# Chapter 14: Protection

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- Goals of Protection
- Principles of Protection
- Domain of Protection
- Access Matrix
- Implementation of Access Matrix
- Access Control
- Revocation of Access Rights
- Capability-Based Systems
- Language-Based Protection

### Objectives

- Discuss the goals and principles of protection in a modern computer system
- Explain how protection domains combined with an access matrix are used to specify the resources a process may access
- Examine capability and language-based protection systems

#### Goals of Protection

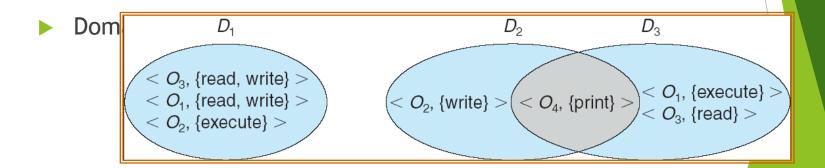
- Operating system consists of a collection of objects, hardware or software
- ► Each object has a unique name and can be accessed through a well-defined set of operations.
- Protection problem ensure that each object is accessed correctly and only by those processes that are allowed to do so.

## Principles of Protection

- Guiding principle principle of least privilege
  - Programs, users and systems should be given just enough privileges to perform their tasks

#### **Domain Structure**

Access-right = <object-name, rights-set> where rights-set is a subset of all valid operations that can be performed on the object.

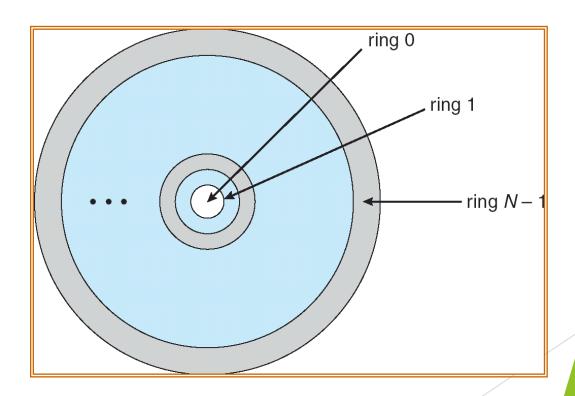


# Domain Implementation (UNIX)

- System consists of 2 domains:
  - User
  - Supervisor
- UNIX
  - Domain = user-id
  - Domain switch accomplished via file system.
    - ▶ Each file has associated with it a domain bit (setuid bit).
    - ▶ When file is executed and setuid = on, then user-id is set to owner of the file being executed. When execution completes user-id is reset.

# Domain Implementation (MULTICS)

- Let  $D_i$  and  $D_j$  be any two domain rings.
- $\blacktriangleright \quad \mathsf{lf} \ j < I \Rightarrow \mathsf{D}_i \ \subseteq \mathsf{D}_j$



#### **Access Matrix**

- View protection as a matrix (access matrix)
- Rows represent domains
- Columns represent objects
- Access(i, j) is the set of operations that a process executing in Domain; can invoke on Object;

### **Access Matrix**

object domain	F <sub>1</sub>	F <sub>2</sub>	<i>F</i> <sub>3</sub>	printer
$D_1$	read		read	
$D_2$				print
$D_3$		read	execute	
$D_4$	read write		read write	

#### Use of Access Matrix

- If a process in Domain  $D_i$  tries to do "op" on object  $O_j$ , then "op" must be in the access matrix.
- Can be expanded to dynamic protection.
  - Operations to add, delete access rights.
  - Special access rights:
    - $\triangleright$  owner of  $O_i$
    - $\triangleright$  copy op from  $O_i$  to  $O_j$
    - $\triangleright$  control  $D_i$  can modify  $D_i$  access rights
    - transfer switch from domain  $D_i$  to  $D_j$

### Use of Access Matrix (Cont.)

- Access matrix design separates mechanism from policy.
  - Mechanism
    - Operating system provides access-matrix + rules.
    - If ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced.
  - Policy
    - User dictates policy.
    - ▶ Who can access what object and in what mode.

# Implementation of Access Matrix

► Each column = Access-control list for one object Defines who can perform what operation.

```
Domain 1 = Read, Write
Domain 2 = Read
Domain 3 = Read
```

:

Each Row = Capability List (like a key) Fore each domain, what operations allowed on what objects.

Object 1 - Read

Object 4 - Read, Write, Execute

Object 5 - Read, Write, Delete, Copy

#### Access Matrix of Figure A With Domains as Objects

object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>	laser printer	<i>D</i> <sub>1</sub>	$D_2$	$D_3$	$D_4$
$D_1$	read		read			switch		
$D_2$				print			switch	switch
$D_3$		read	execute					
$D_4$	read write		read write		switch			

Figure B

# Access Matrix with Copy Rights

object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>			
$D_1$	execute		write*			
$D_2$	execute	read*	execute			
$D_3$	execute					
(a)						
object domain	F <sub>1</sub>	$F_2$	$F_3$			
$D_1$	execute		write*			
$D_2$	execute	read*	execute			
$D_3$	execute	read				
(b)						

# Access Matrix With *Owner*Rights

object domain	F <sub>1</sub>	F <sub>2</sub>	F <sub>3</sub>			
$D_1$	owner execute		write			
$D_2$		read* owner	read* owner write			
$D_3$	execute					
(a)						
object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>			
<i>D</i> <sub>1</sub>	owner execute		write			
$D_2$		owner read* write*	read* owner write			
$D_3$		write	write			
(b)						

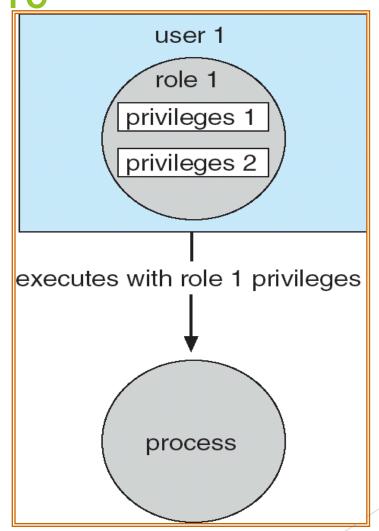
# Modified Access Matrix of Figure B

object domain	F <sub>1</sub>	$F_2$	F <sub>3</sub>	laser printer	<i>D</i> <sub>1</sub>	$D_2$	$D_3$	$D_4$
$D_1$	read		read			switch		
$D_2$				print			switch	switch control
$D_3$		read	execute					
$D_4$	write		write		switch			

#### **Access Control**

- Protection can be applied to non-file resources
- Solaris 10 provides role-based access control to implement least privilege
  - Privilege is right to execute system call or use an option within a system call
  - Can be assigned to processes
  - Users assigned roles granting access to privileges and programs

Role-based Access Control in Solaris 10



# Revocation of Access Rights

- Access List Delete access rights from access list.
  - Simple
  - Immediate
- Capability List Scheme required to locate capability in the system before capability can be revoked.
  - Reacquisition
  - Back-pointers
  - Indirection
  - Keys

# Capability-Based Systems

#### Hydra

- ► Fixed set of access rights known to and interpreted by the system.
- Interpretation of user-defined rights performed solely by user's program; system provides access protection for use of these rights.

#### Cambridge CAP System

- Data capability provides standard read, write, execute of individual storage segments associated with object.
- Software capability -interpretation left to the subsystem, through its protected procedures.

### Language-Based Protection

- Specification of protection in a programming language allows the high-level description of policies for the allocation and use of resources.
- Language implementation can provide software for protection enforcement when automatic hardwaresupported checking is unavailable.
- Interpret protection specifications to generate calls on whatever protection system is provided by the hardware and the operating system.

#### Protection in Java 2

- Protection is handled by the Java Virtual Machine (JVM)
- A class is assigned a protection domain when it is loaded by the JVM.
- ► The protection domain indicates what operations the class can (and cannot) perform.
- If a library method is invoked that performs a privileged operation, the stack is inspected to ensure the operation can be performed by the library.

# **Stack Inspection**

protection domain:	untrusted applet	URL loader	networking
socket permission:	none	*.lucent.com:80, connect	any
class:	gui: get(url); open(addr);	get(URL u): doPrivileged { open('proxy.lucent.com:80'); } <request from="" proxy="" u=""></request>	open(Addr a):  checkPermission (a, connect); connect (a);