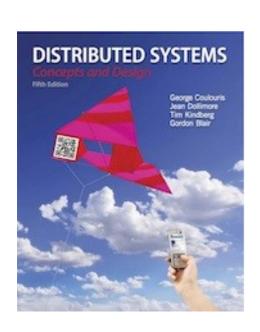
#### Slides for Chapter 3: Networking and Internetworking



# From Coulouris, Dollimore, Kindberg and Blair Distributed Systems: Concepts and Design

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Figure 3.1 Network performance

	Example	Range	Bandwidth (Mbps)	Latency (ms)
Wired:				
LAN	Ethernet	1–2 kms	10-10,000	1-10
WAN	IP routing	worldwide	0.010 – 600	100-500
MAN	ATM	2-50 kms	1-600	10
Internetwork	Internet	worldwide	0.5 - 600	100-500
Wireless:	km			
WPAN	Bluetooth (IEEE 802.15.1)	10-30m	0.5-2	5-20
WLAN	WiFi (IEEE 802.11)	0.15–1.5 km	11 - 108	5-20
WMAN	WiMAX (IEEE 802.16)	5–50 km	1.5 - 20	5-20
WWAN	3G phone	cell: 15	348-14.4	100-500

Figure 3.2 Conceptual layering of protocol software

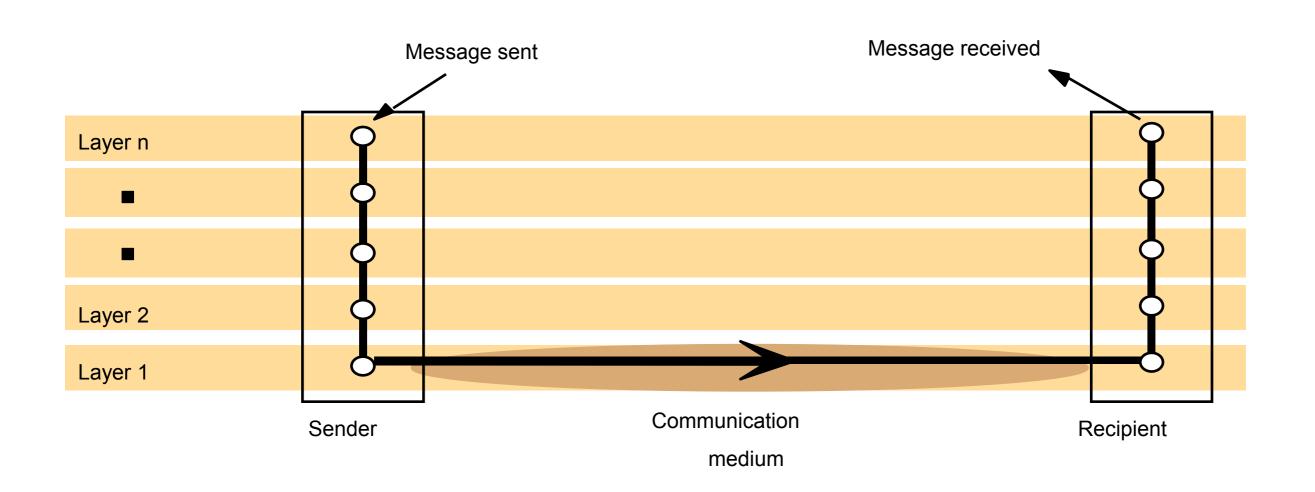


Figure 3.3 Encapsulation as it is applied in layered protocols

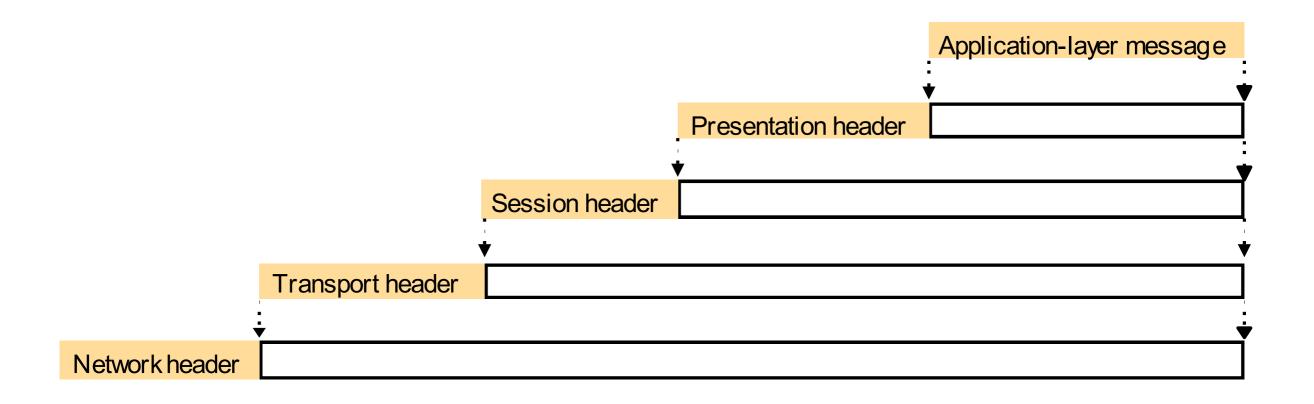
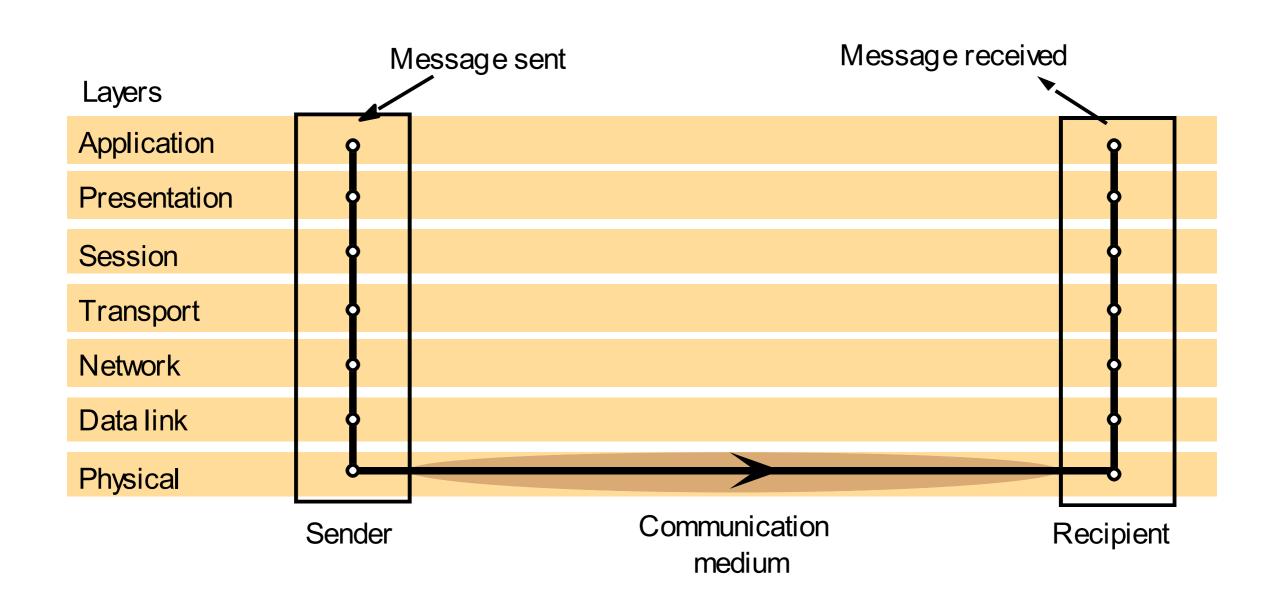


Figure 3.4 Protocol layers in the ISO Open Systems Interconnection (OSI) model



## Figure 3.5 OSI protocol summary

Layer	Description	Examples
Application	Protocols that are designed to meet the communication requirements of specific applications, often defining the interface to a service.	HTTP, FTP , SMTP, CORBA IIOP
Presentation	Protocols at this level transmit data in a network representation that is independent of the representations used in individual computers, which may differ. Encryption is also performed in this layer, if required.	Secure Sockets (SSL),CORBA Data Rep.
Session	At this level reliability and adaptation are performed, such as detection of failures and automatic recovery.	
Transport	This is the lowest level at which messages (rather than packets) are handled. Messages are addressed to communication ports attached to processes, Protocols in this layer may be connection-oriented or connectionless.	TCP, UDP
Network	Transfers data packets between computers in a specific network. In a WAN or an internetwork this involves the generation of a route passing through routers. In a single LAN no routing is required.	IP, ATM virtual circuits
Data link	Responsible for transmission of packets between nodes that are directly connected by a physical link. In a WAN transmission is between pairs of routers or between routers and hosts. In a LAN it is between any pair of hosts.	Ethernet MAC, ATM cell transfer, PPP
Physical	The circuits and hardware that drive the network. It transmits sequences of binary data by analogue signalling, using amplitude or frequency modulation of electrical signals (on cable circuits), light signals (on fibre optic circuits) or other electromagnetic signals (on radio and microwave circuits).	Ethernet base- band signalling, ISDN

Figure 3.6 Internetwork layers

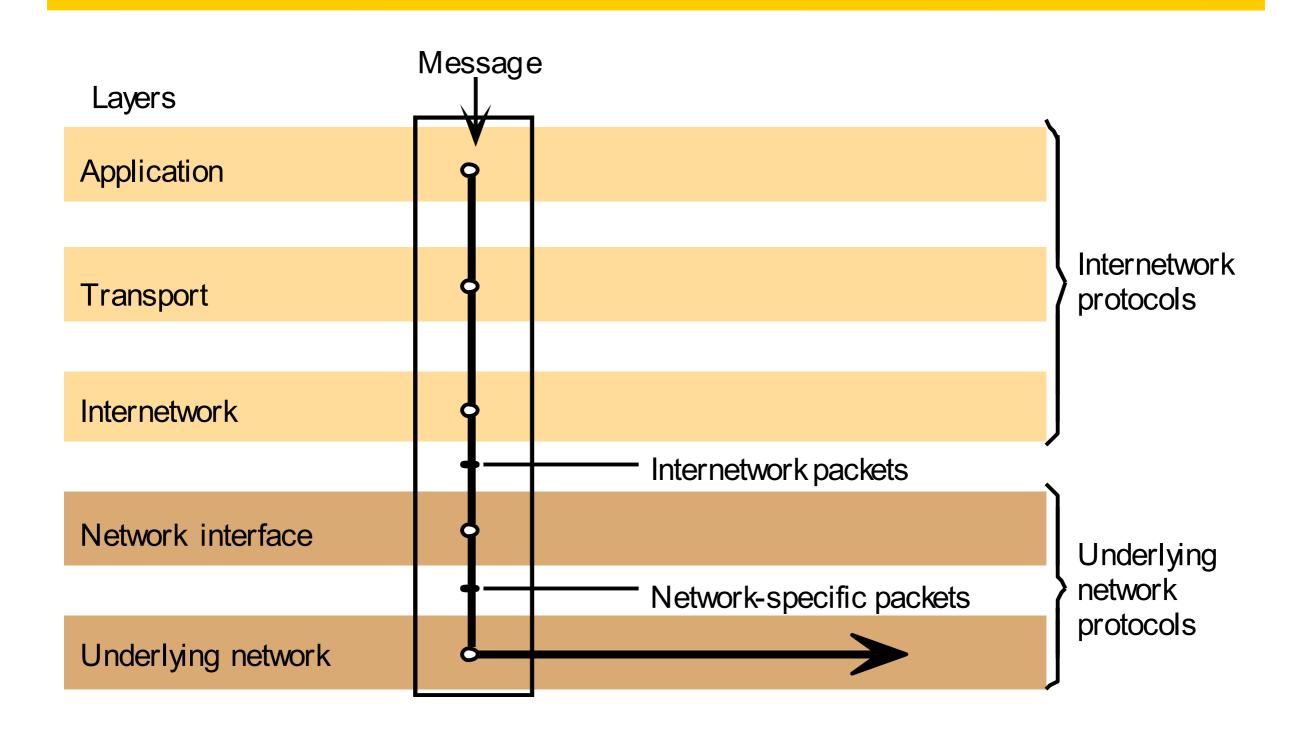


Figure 3.7 Routing in a wide area network

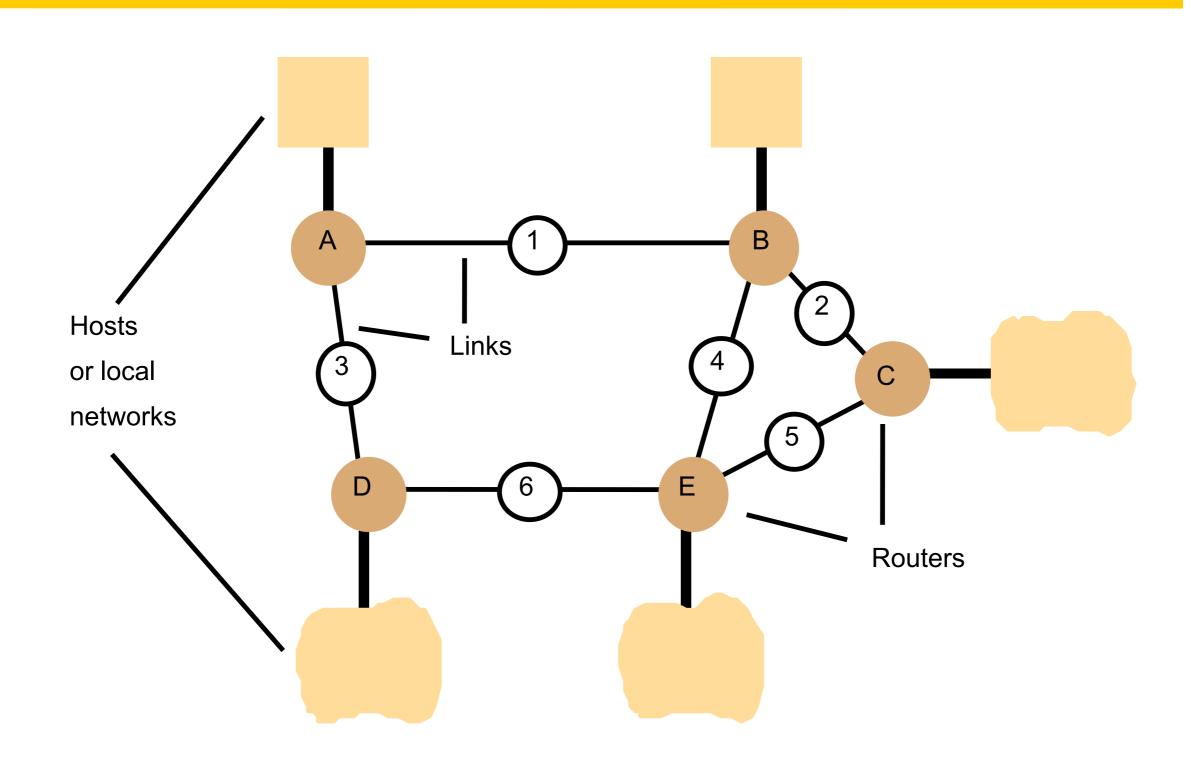


Figure 3.8 Routing tables for the network in Figure 3.7

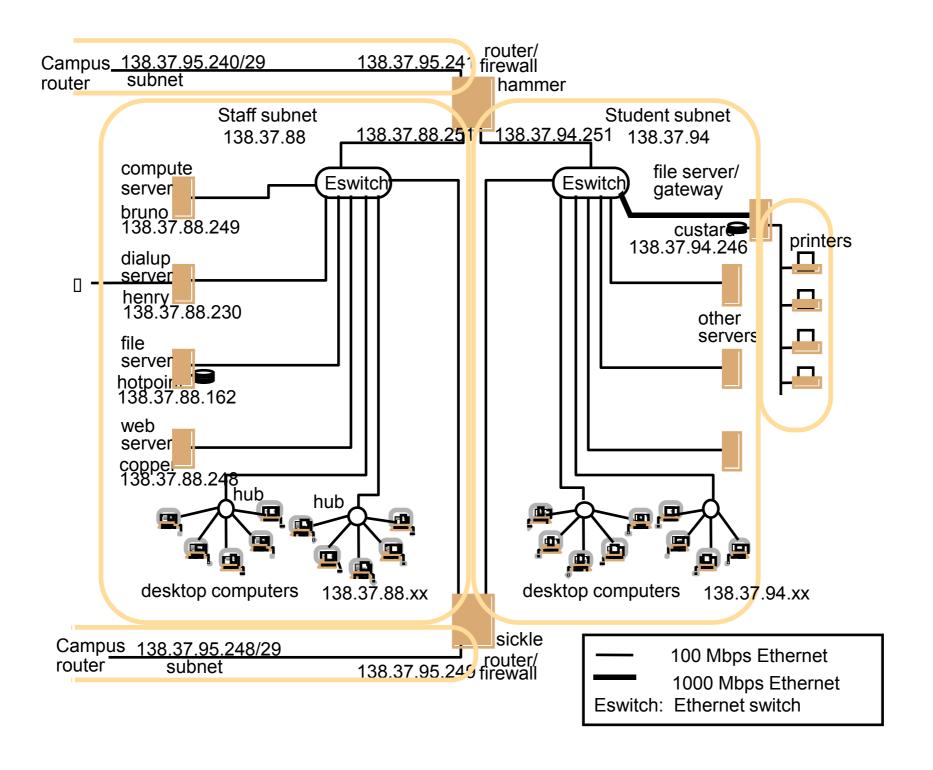
Routings from A		Routings from B			Routings from C				
То	Link	Cost		То	Link	Cost	То	Link	Cost
A	local	0		A	1	1	A	2	2
В	1	1		В	local	0	В	2	1
$\mathbf{C}$	1	2		C	2	1	$\mathbf{C}$	local	0
D	3	1		D	1	2	D	5	2
E	1	2		E	4	1	E	5	1

Rout	ings from D		· -	Rout	ings from E	
То	Link	Cost		То	Link	Cost
A	3	1	-	A	4	2
В	3	2		В	4	1
C	6	2		C	5	1
D	local	0		D	6	1
E	6	1		E	local	0

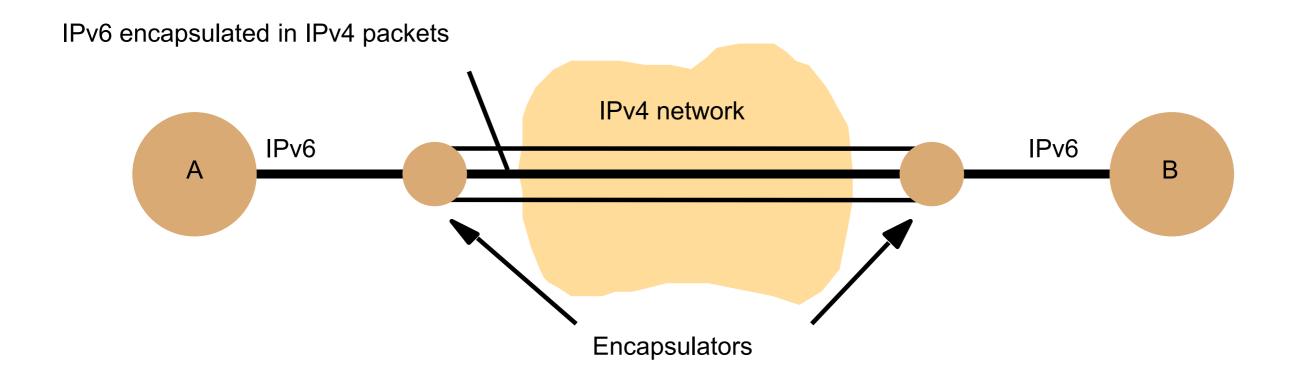
#### Figure 3.9 Pseudo-code for RIP routing algorithm

```
Send: Each t seconds or when Tl changes, send Tl on each non-faulty outgoing link.
Receive: Whenever a routing table Tr is received on link n:
for all rows Rr in Tr {
if (Rr.link \mid n) {
Rr.cost = Rr.cost + 1;
Rr.link = n;
if (Rr.destination is not in Tl) add Rr to Tl;
// add new destination to Tl
else for all rows Rl in Tl {
if (Rr.destination = Rl.destination and
       (Rr.cost < Rl.cost \text{ or } Rl.link = n)) Rl = Rr;
// Rr.cost < Rl.cost : remote node has better route
// Rl.link = n: remote node is more authoritative
```

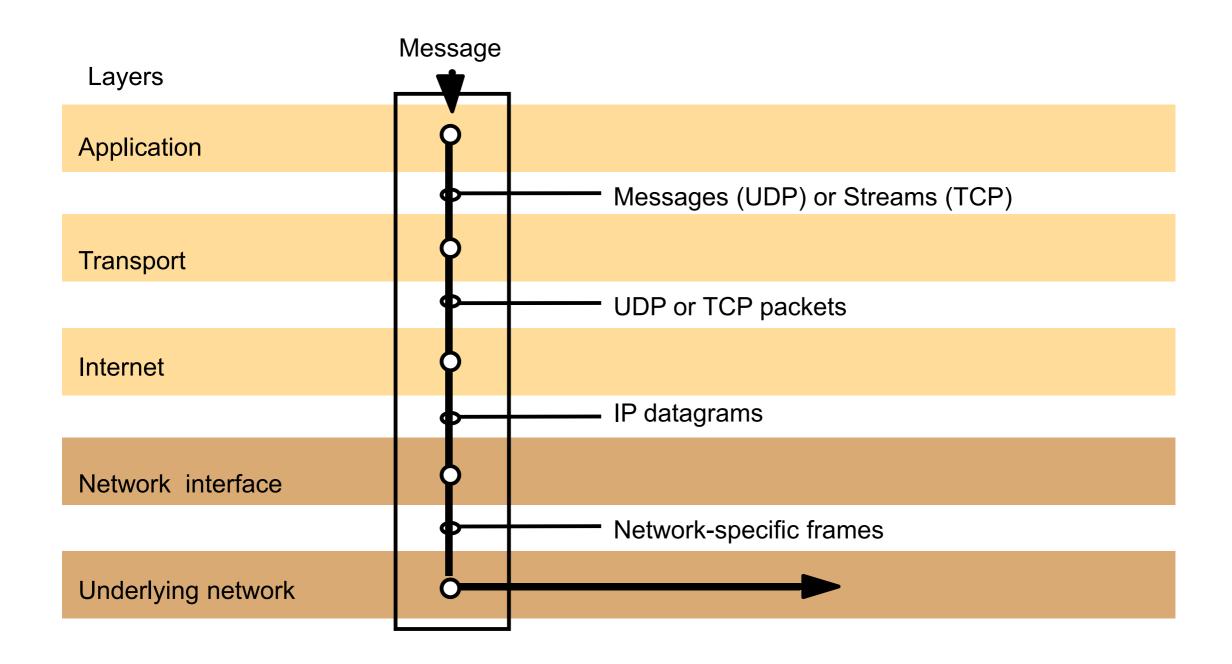
Figure 3.10 Simplified view of part of a university campus network



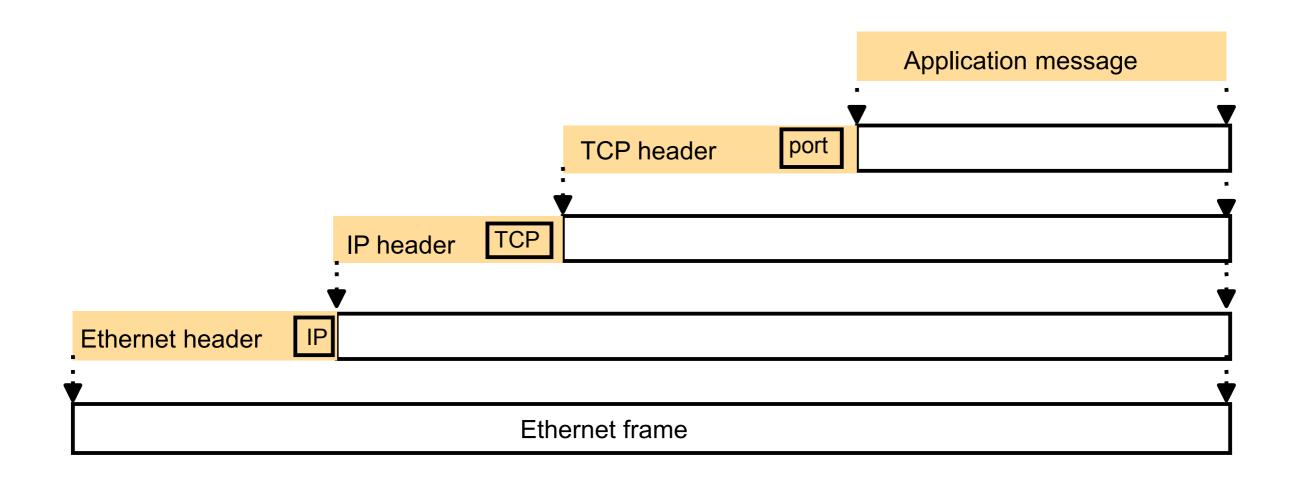
### Figure 3.11 Tunnelling for IPv6 migration



### Figure 3.12 TCP/IP layers



#### Figure 3.13 Encapsulation in a message transmitted via TCP over an Ethernet



#### Figure 3.14 The programmer's conceptual view of a TCP/IP Internet

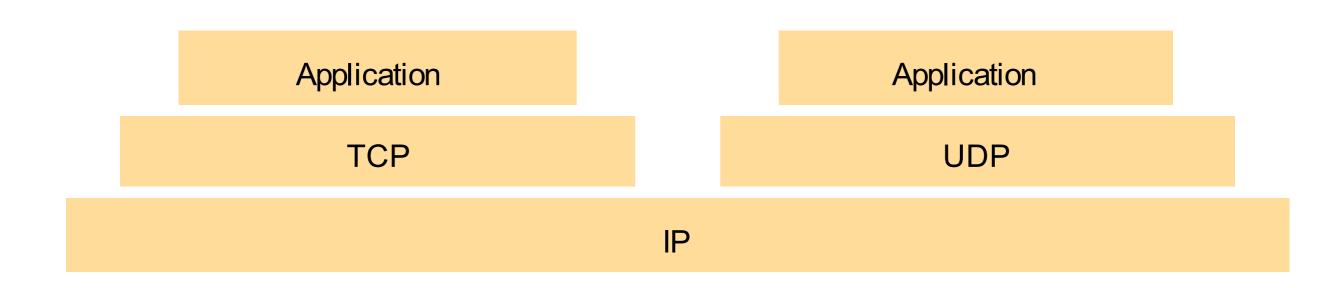


Figure 3.15 Internet address structure, showing field sizes in bits

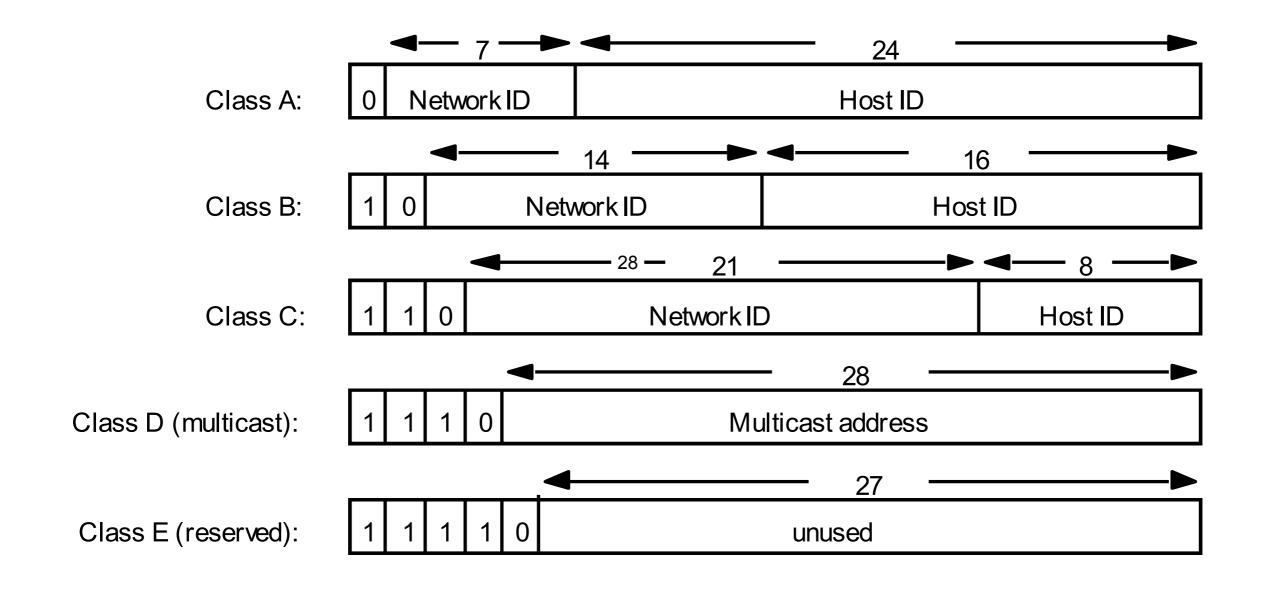
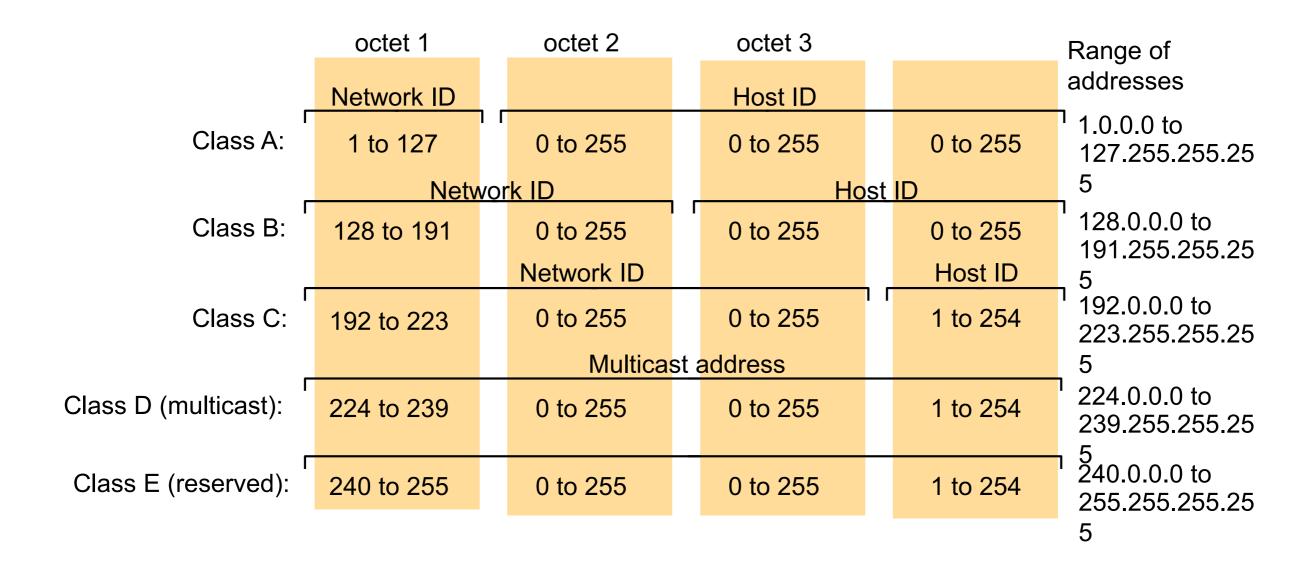


Figure 3.16
Decimal representation of Internet addresses



#### Figure 3.17 IP packet layout

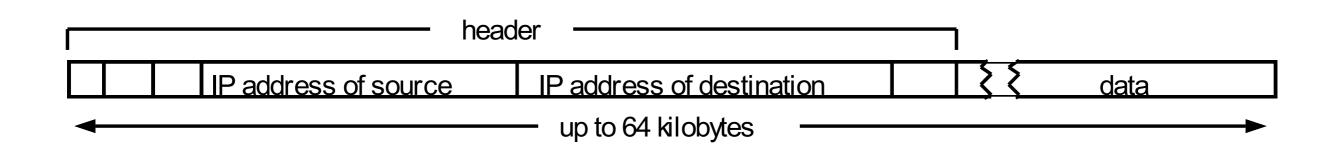
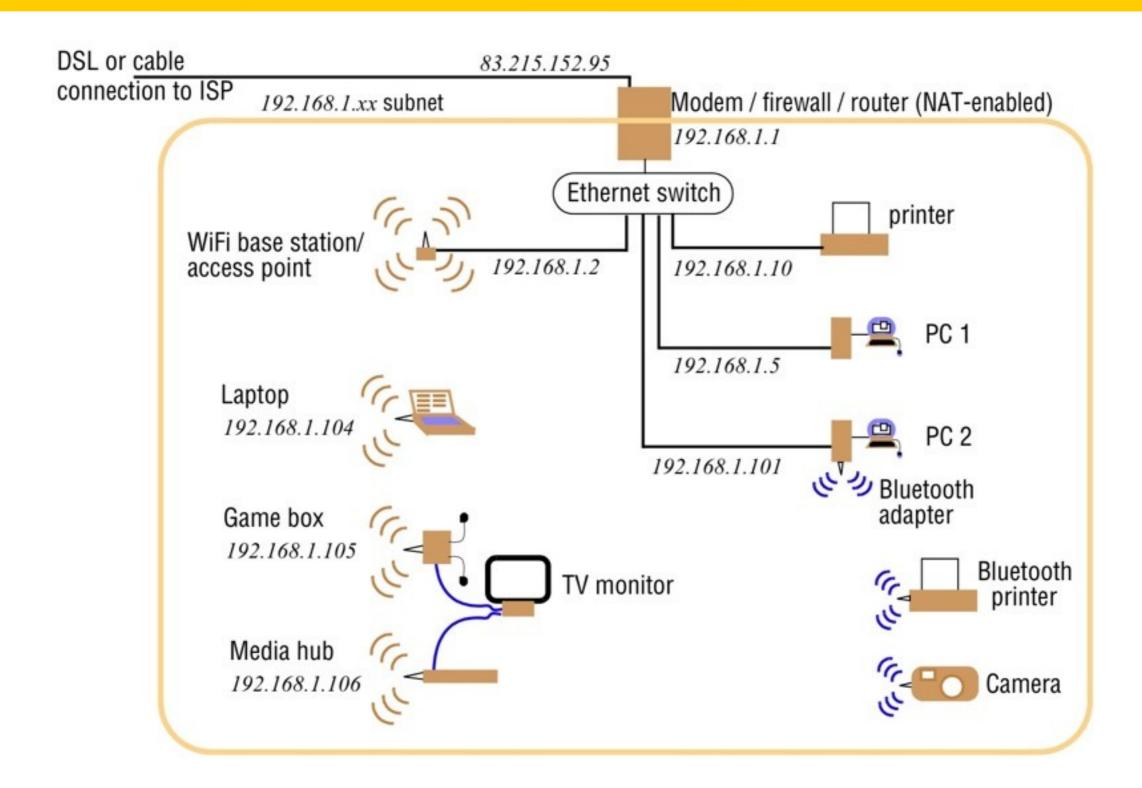


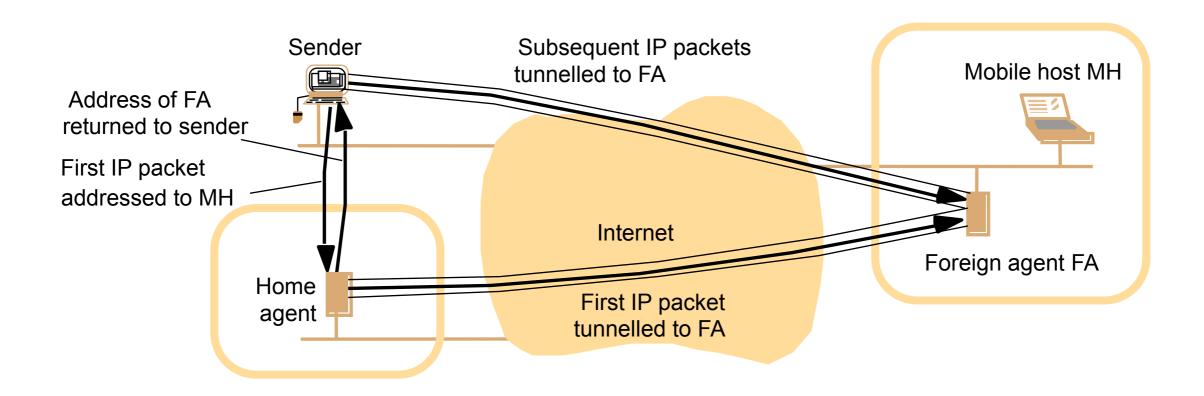
Figure 3.18 A typical NAT-based home network



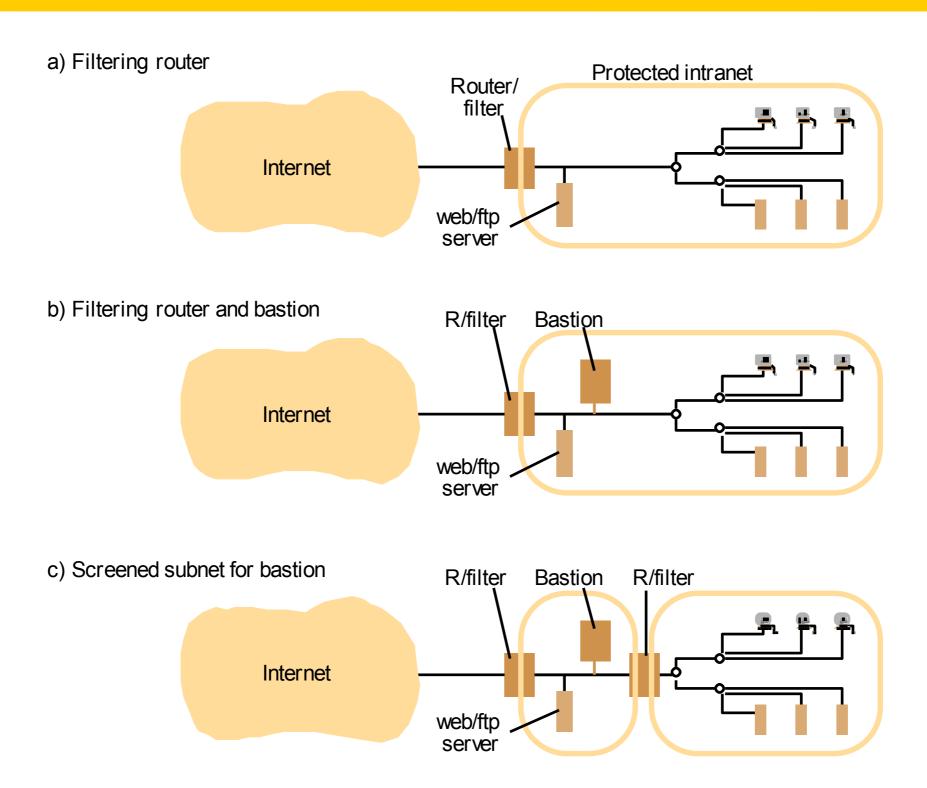
## Figure 3.19 IPv6 header layout

Version (4	Traffic class (8	Flow label (2	Flow label (20					
bits)Payload lengerts)16		Next <sup>hitea</sup> der (8	Hop limit (8					
bits)		bits)	bits)					
	Source address bits)							
	addited.							
Destination addrése8 bits)								

#### Figure 3.20 The MobileIP routing mechanism



#### Figure 3.21 Firewall configurations



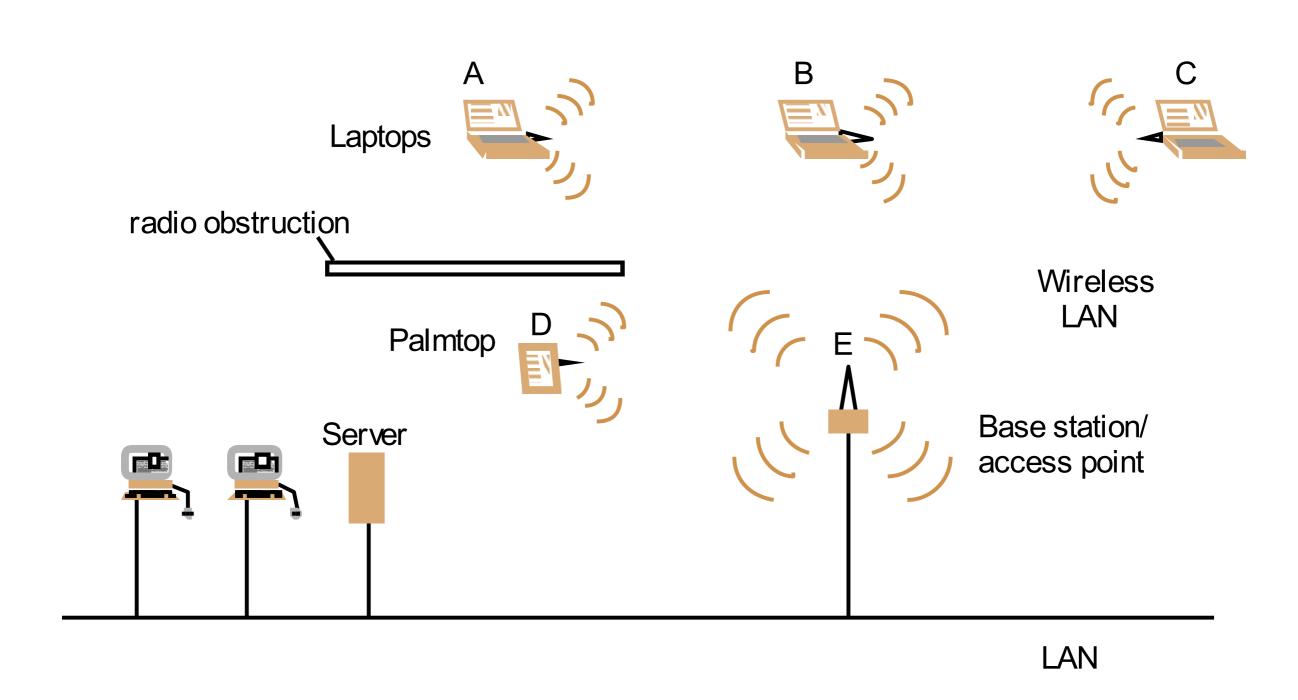
#### Figure 3.22 IEEE 802 network standards

IEEE No.	Name	Title	Reference
802.3	Ethernet	CSMA/CD Networks (Ethernet)	[IEEE 1985a]
802.4		Token Bus Networks	[IEEE 1985b]
802.5		Token Ring Networks	[IEEE 1985c]
802.6		Metropolitan Area Networks	[IEEE 1994]
802.11	WiFi	Wireless Local Area Networks	[IEEE 1999]
802.15.1	Bluetooth	Wireless Personal Area Networks	[IEEE 2002]
802.15.4	ZigBee	Wireless Sensor Networks	[IEEE 2003]
802.16	WiMAX	Wireless Metropolitan Area Networks	[IEEE 2004a]

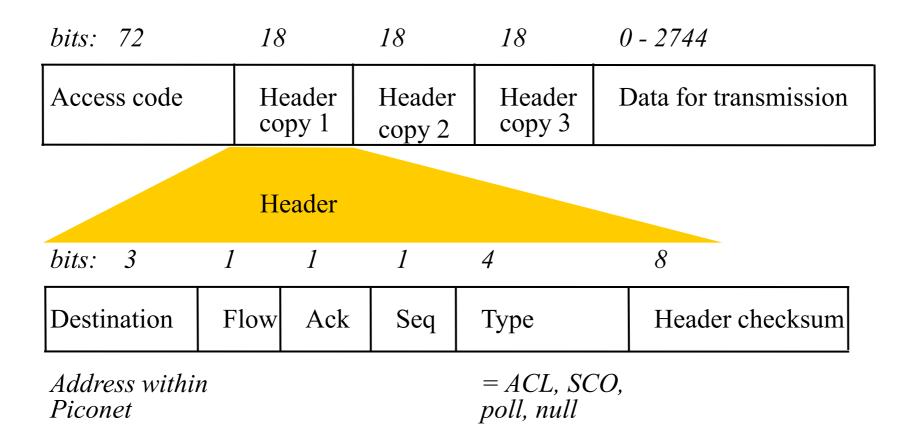
Figure 3.23 Ethernet ranges and speeds

	10Base5	10BaseT	100BaseT	1000BaseT		
Data rate	10 Mbps	10 Mbps	100 Mbps	1000 Mbps		
Max. segment lengths	Max. segment lengths:					
Twisted wire (UTP)	100 m	100 m	100 m	25 m		
Coaxial cable (STP)	500 m	500 m	500 m	25 m		
Multi-mode fibre	2000 m	2000 m	500 m	500 m		
Mono-mode fibre	25000 m	25000 m	20000 m	2000 m		

Figure 3.24
Wireless LAN configuration



#### Figure 3.25 Bluetooth frame structure



SCO packets (e.g. for voice data) have a 240-bit payload containing 80 bits of data triplicated, filling exactly one timeslot.