# OPERATING SYSTEMS 

## SECURITY

## SECURITY

## In This Chapter:

-The Security Problem
-Program Threats
-System and Network Threats
-Cryptography as a Security Tool

- User Authentication
-Implementing Security Defenses
-Firewalling to Protect Systems and Networks
-Computer-Security Classifications
-An Example: Windows XP


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## SECURITY ISSUES:

External protection of a system. A classified site goes to extraordinary lengths to keep things physically tight. Among the issues to be considered:

Unauthorized access Mechanism assuring only authorized individuals see classified materials.

Malicious
Accidental
Authentication
modification or destruction introduction of inconsistency. How do we know the user is who she says she is. Can have passwords on domains.

Protection of passwords is difficult. Issues include:

- It's very easy to guess passwords since people use simple and easily remembered words.
- Need exists to change passwords continually.
- Limiting number of tries before locking up.
- How to crack UNIX passwords.

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## Security Issues

| Trojan Horse: | A piece of code that misuses its environment. <br> innocent enough, however when executed, unexpected behavior <br> occurs. <br> Inserting a method of breaching security in a system. For instance, <br> some secret set of inputs to a program might provide special privileges. |
| :--- | :--- |
| Trap Doors: | The program seems <br> Threat monitoring: Look for unusual activity. Once access is gained, how do you identify <br> someone acting in an unusual fashion? |
| Audit Log: | Record time, user, and type of access on all objects. Trace problems <br> back to source. |
| Worms | Use spawning mechanism; standalone programs. <br> In the Internet worm, Robert Morse exploited UNIX networking features <br> (remote access) as well as bugs in finger and sendmail programs. <br> Grappling hook program uploaded main worm program. |
| Viruses | Fragment of code embedded in a legitimate program. Mainly effects <br> personal PC systems. These are often downloaded via e-mail or as <br> active components in web pages. |
| A mechanism that allows only certain traffic between trusted and un- |  |
| frusted systems. Often applied to a way to keep unwanted internet |  |

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## Typical Security Attacks

## ATTACK METHODS:

Attacks on a distributed system include:

- Passive wiretapping. ( unauthorized interception/reading of messages )
- Active wiretapping:

Modification Changing a portion of the message.

Spurious messages Introducing bogus messages with valid addresses and consistency criteria.

Site impersonation Claiming to be some other logical node.

Replay of previous transmission - repeating previous valid messages. (for example, authorization of cash withdrawal.)

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## ATTACK METHODS:

## Typical Security Attacks



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## Typical Security Attacks

## ATTACK METHODS:

- Trojan Horse
- Code segment that misuses its environment
- Exploits mechanisms for allowing programs written by users to be executed by other users
- Spyware, pop-up browser windows, covert channels
- Trap Door
- Specific user identifier or password that circumvents normal security procedures
- Could be included in a compiler
- Logic Bomb
- Program that initiates a security incident under certain circumstances
- Stack and Buffer Overflow
- Exploits a bug in a program (overflow either the stack or memory buffers)


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## Typical Security Attacks

## Example of Buffer Overflow Waiting To Happen:

```
#include <stdio.h>
#define BUFFER SIZE 256
int main(int argc, char *argv[])
{
    char buffer[BUFFER SIZE];
    int other_data;
    if (argc < 2)
        return -1;
    else {
        strcpy (buffer,argv[1]);
        return 0;
    }
}

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\section*{Typical Security Attacks}

\section*{Viruses}
- Code fragment embedded in legitimate program
- Very specific to CPU architecture, operating system, applications
- Usually borne via email or as a macro
- Visual Basic Macro to reformat hard drive
```

Su.b AutoOpen()
Dim OFS
Set OFS =
CreateObject(''Scripting.FileSystemObject'')
vS = Shell(''c:command.com /k format c:'r'vbHide)
End Sub

```

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\section*{Typical Security Attacks}

A Boot Sector Virus


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\section*{Typical Security Attacks}

\section*{System And Network Threats}
- Worms - use spawn mechanism; standalone program
- Internet worm
- Exploited UNIX networking features (remote access) and bugs in finger and sendmail programs. (See next slide)
- Grappling hook program uploaded main worm program
- Port scanning
- Automated attempt to connect to a range of ports on one or a range of IP addresses
- Denial of Service
- Overload the targeted computer preventing it from doing any useful work
- Distributed denial-of-service (DDOS) come from multiple sites at once

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\section*{Typical Security Attacks}

Morris Internet Worm


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\section*{Cryptography}

\section*{DEFINITIONS:}

Encryption:
\[
C=E(M, K e)
\]

E = Encyphering Algorithm
M = Message - plain text
\(\mathrm{Ke}=\) Encryption key
C = Cyphered text

\section*{Decryption:}
\(M=D(C, K d)\)

D = Decyphering Algorithm
Kd = Decryption key


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\section*{Cryptography}

\section*{DEFINITIONS:}

\section*{Cryptosystems are either Conventional or Public Key}
- Conventional is symmetric; \(\mathrm{Ke}=\mathrm{Kd}\), so the key must be kept secret. Algorithms are simple to describe, but complex in the number of operations.
- Public key is asymmetric; Ke != Kd , so Ke can be made public. Kd is secret and can't easily be derived from Ke .

Security against attack is either:
- Unconditionally secure - Ke can't be determined regardless of available computational power.
- Computationally secure: - calculation of Kd is economically unfeasible (it would overwhelm all available computing facilities.)

The only known unconditionally secure system in common use!
- Involves a random key that has the same length as the plain text to be encrypted.
- The key is used once and then discarded. The key is exclusively OR'd with the message to produce the cypher.
- Given the key and the cypher, the receiver uses the same method to reproduce the message.

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\section*{Data Encryption Standard}

\section*{DATA ENCRYPTION STANDARD ( DES ):}
- The official National Institute of Standards and Technology (NIST), (formerly the National Bureau of Standards) encryption for use by Federal agencies.
- The source of security is the non-linear many-to-one function applied to a block of data. This function uses transposition and substitution. The algorithm is public, but the key ( 56 bits) is secret.
- Computational power today can crack a 56 bit code.
- In common use today is Triple DES in which 3 different keys are used, making the effective key length 168 bits.

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\section*{Public Key Cryptosystems}

The general principle is this:
1. Any RECEIVER A uses an algorithm to calculate an encryption key KEa and a decryption key KDa.
2. Then the receiver PUBLICIZES KEa to anyone who cares to hear. But the receiver keeps secret the decryption key KDa.
3. User B sends a message to A by first encrypting that message using the publicized key for that receiver A, KEa.
4. Since only \(\mathbf{A}\) knows how to decrypt the message, it's secure.


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\section*{Public Key Cryptosystems}

To be effective, a system must satisfy the following rules:
a) Given plaintext and ciphertext, the problem of determining the keys is computationally complex.
b) It is easy to generate matched pairs of keys \(\mathrm{Ke}, \mathrm{Kd}\) that satisfy the property
\[
\mathrm{D}(\mathrm{E}(\mathrm{M}, \mathrm{Ke}), \mathrm{Kd})=\mathrm{M} .
\]

This implies some sort of trapdoor, such that Ke and Kd can be calculated from first principles, but one can't be derived from the other.
c) The encryption and decryption functions E and D are efficient and easy to use.
d) Given Ke , the problem of determining Kd is computationally complex.

What is computationally difficult? Problems that can't easily be calculated in a finite time.
Examples include: factoring the product of two very large prime numbers; the knapsack problem.
These problems are NP complete - solution times are exponential in the size of the sample.

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\section*{Public Key Cryptosystems}

To be effective, a system must satisfy the following rules:
e) For almost all messages it must be computationally unfeasible to find ciphertext key pairs that will produce the message.
(In other words, an attacker is forced to discover the true (M,Ke) pair that was used to create the ciphertext C .)
f) Decryption is the inverse of encryption.
\[
E(D(M, K d), K e)=D(E(M, K e), K d)
\]

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\section*{Public Key Cryptosystems}

\section*{AN EXAMPLE:}
1. Two large prime numbers \(p\) and \(q\) are selected using some efficient test for primality. These numbers are
\[
\text { Let } p=3, q=11
\] secret:
\[
n=3 * 11=33 .
\]
2. The product \(n=p\) * \(q\) is computed.
3. The number \(K d>\max (p, q)\) is picked at random from the set of integers that are relatively prime to and less than \(L(n)=(p-1)(q-1)\).
\(L(n)=(p-1)(q-1)=20\). Choose Kd > 11 and prime to 20. Choose Kd = 13.
4. The integer \(\mathrm{Ke}, \quad 0<\mathrm{Ke}<\mathrm{L}(\mathrm{n})\) is computed from \(L(n)\) and \(K d\) such that \(K e * K d=1(\bmod L(n))\).
\[
\begin{aligned}
& 0<K e<20 \\
& \mathrm{Ke}=17 . \quad(\text { since } 17 * 13=221=1(\bmod 20))
\end{aligned}
\]

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\section*{Public Key Cryptosystems}

\section*{AN EXAMPLE:}

Separate the text to be encoded into chunks with values \(0-(n-1)\).
```

In our example, we'll use < space = 0, A = 1, B = 2, C = 3, D = 4, E = 5 >.
Then " B A D <sp> B E E " --> "21 0400 25 05"
21^17(mod 33) = 21. 21^13 (mod 33) = 21.
04^17(mod 33) = 16. 16^13 (mod 33) = 04.
00^17(mod 33) = 00. 00^13 (mod 33) = 00.
25^17(mod 33) = 31. 31^13 (mod 33) = 25.
05^17(mod 33) = 14. 14^ 13 (mod 33) = 05.

```

This whole operation works because, though \(n\) and \(K e\) are known, \(p\) and \(q\) are not public. Thus Kd is hard to guess.
[Note: recently a 100 digit number was successfully factored into two prime numbers.]

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\section*{Public Key Cryptosystems}

\section*{AUTHENTICATION AND DIGITAL SIGNATURES:}

\section*{Sender Authentication:}

In a public key system, how does the receiver know who sent a message (since the receiver's encryption key is public)?

Suppose A sends message M to B:
a) \(\mathbf{A}\) DECRYPTS \(\mathbf{M}\) using \(\mathbf{A} \mathbf{s} \mathbf{K d}(\mathbf{A})\).
b) A attaches its identification to the message.
c) A ENCRYPTS the entire message using \(\mathbf{B}\) 's encryption, \(\mathbf{K e}(\mathbf{B})\)
\[
C=E((A, D(M, K d(A))), \operatorname{Ke}(B))
\]
d) B decrypts using its private key \(\operatorname{Kd}(\mathbf{A})\) to produce the pair \(\mathbf{A}, \mathrm{D}(\mathbf{M}, \operatorname{Kd}(\mathbf{A}))\).
e) Since the proclaimed sender is \(\mathbf{A}, \quad \mathbf{B}\) knows to use the public encryption key \(\operatorname{Ke(A)}\).

\section*{Capture/Replay}

In this case, a third party could capture / replay a message.
The solution is to use a rapidly changing value such as time or a sequence number as part of the message.

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\section*{Man-in-the-middle Attack on Asymmetric Cryptography}

Here are the attack steps for this scenario:
1. Sender wishes to send a message to Receiver.
2. \(S\) asks \(R\) for its encryption key.
3. When R returns key, that key is intercepted by the attacker who substitutes her key.
4. Sender encrypts message using this bogus key and returns it.
5. Since the attacker is the owner of this bogus key, the attacker can read the message.

Public Key Cryptosystems


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\section*{Example - SSL}
- Insertion of cryptography at one layer of the ISO network model (the transport layer)
- SSL - Secure Socket Layer (also called TLS)
- Cryptographic protocol that limits two computers to only exchange messages with each other
- Very complicated, with many variations
- Used between web servers and browsers for secure communication (credit card numbers)
- The server is verified with a certificate assuring client is talking to correct server
- Asymmetric cryptography used to establish a secure session key (symmetric encryption) for bulk of communication during session
- Communication between each computer uses symmetric key cryptography

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\section*{Example - Windows XP}
- Security is based on user accounts
- Each user has unique security ID
- Login to ID creates security access token
- Includes security ID for user, for user's groups, and special privileges
- Every process gets copy of token
- System checks token to determine if access allowed or denied
- Uses a subject model to ensure access security. A subject tracks and manages permissions for each program that a user runs
- Each object in Windows XP has a security attribute defined by a security descriptor
- For example, a file has a security descriptor that indicates the access permissions for all users
U.S. Department of Defense outlines four divisions of computer security: A, B, C, and D.
- D - Minimal security.
- C - Provides discretionary protection through auditing. Divided into C1 and C2. C1 identifies cooperating users with the same level of protection. C2 allows user-level access control.
- B - All the properties of \(\mathbf{C}\), however each object may have unique sensitivity labels. Divided into B1, B2, and B3.
- A - Uses formal design and verification techniques to ensure security.

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\section*{Wrap Up}

In this chapter we've looked at how to secure information that may be placed in hazardous public forums.

Data on the net is an excellent example here.```

