

# SOLUTION OF SYSTEM OF LINEAR EQUATIONS



#### Solution of linear system of equations

- Numerical solution of differential equations (Finite Difference Method)
- Numerical solution of integral equations (Finite Element Method, Method of Moments)

$$\begin{vmatrix}
a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1 \\
a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2 \\
\vdots & \vdots & \ddots & \vdots \\
a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = b_n
\end{vmatrix} \Rightarrow \begin{bmatrix}
a_{11} & a_{12} & \dots & a_{1n} \\
a_{21} & a_{22} & \dots & a_{2n} \\
\vdots & \vdots & \ddots & \vdots \\
a_{n1} & a_{n2} & \dots & a_{nn}
\end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$



#### Consistency (Solvability)

The linear system of equations Ax=b has a solution, or said to be consistent IFF Rank{A}=Rank{A|b}

A system is inconsistent when Rank{A}<Rank{A|b}</li>

Rank{A} is the maximum number of linearly independent columns or rows of A. Rank can be found by using ERO (Elementary Row Oparations) or ECO (Elementary column operations).

ERO⇒# of rows with at least one nonzero entry ECO⇒# of columns with at least one nonzero entry



#### Elementary row operations

- The following operations applied to the augmented matrix [A|b], yield an equivalent linear system
  - Interchanges: The order of two rows can be changed
  - Scaling: Multiplying a row by a nonzero constant
  - Replacement: The row can be replaced by the sum of that row and a nonzero multiple of any other row.



## An inconsistent example

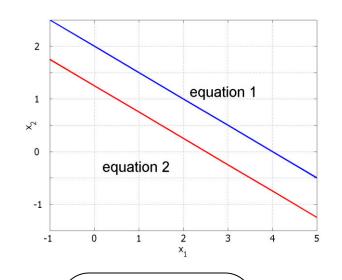
$$\begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 4 \\ 5 \end{bmatrix}$$

ERO: Multiply the first row with -2 and add to the second row

$$\begin{bmatrix} 1 & 2 \\ 0 & 0 \end{bmatrix}$$

$$Rank{A}=1$$

$$\begin{bmatrix} 1 & 2 & 4 \\ 0 & 0 & -3 \end{bmatrix} \quad \text{Rank}\{A|b\}=2$$



Then this system of equations is not solvable



# Uniqueness of solutions

The system has a unique solution IFF

$$Rank{A}=Rank{A|b}=n$$

n is the order of the system

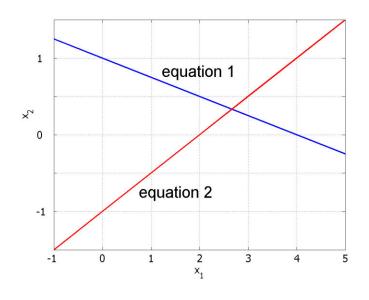
Such systems are called full-rank systems



#### Full-rank systems

If Rank{A}=n
 Det{A} ≠ 0 ⇒ A is nonsingular so invertible
 Unique solution

$$\begin{bmatrix} 1 & 2 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 4 \\ 2 \end{bmatrix}$$



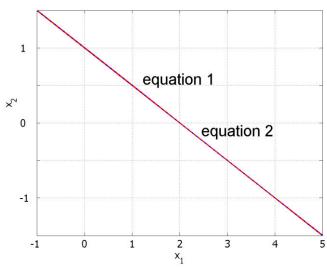


#### Rank deficient matrices

■ If Rank{A}=m<n Det{A} = 0 ⇒ A is singular so not invertible infinite number of solutions (n-m free variables) under-determined system

$$\begin{bmatrix} 1 & 2 \\ 2 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 4 \\ 8 \end{bmatrix}$$

Rank{A}=Rank{A|b}=1 Consistent so solvable





# Ill-conditioned system of equations

 A small deviation in the entries of A matrix, causes a large deviation in the solution.

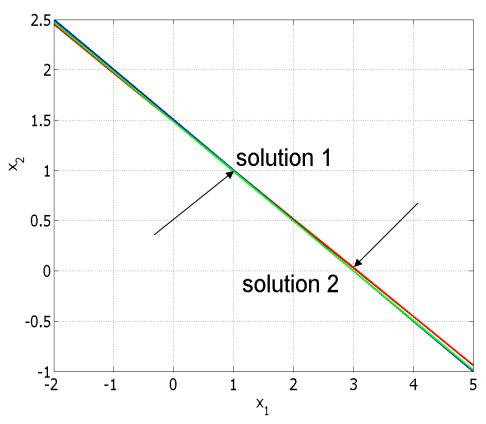
$$\begin{bmatrix} 1 & 2 \\ 0.48 & 0.99 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1.47 \end{bmatrix} \Rightarrow \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 2 \\ 0.49 & 0.99 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1.47 \end{bmatrix} \Rightarrow \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 0 \end{bmatrix}$$



#### Ill-conditioned continued.....

A linear system of equations is said to be "ill-conditioned" if the coefficient matrix tends to be singular





## Types of linear system of equations

- Coefficient matrix A is square and real
- The RHS vector b is nonzero and real
- Consistent system, solvable
- Full-rank system, unique solution
- Well-conditioned system



#### Solution Techniques

- Direct solution methods
  - Finds a solution in a finite number of operations by transforming the system into an <u>equivalent system</u> that is 'easier' to solve.
  - Diagonal, upper or lower triangular systems are easier to solve
  - Number of operations is a function of system size n.
- Iterative solution methods
  - Computes succesive approximations of the solution vector for a given A and b, starting from an initial point x<sub>0</sub>.
  - Total number of operations is uncertain, may not converge.



#### **Direct solution Methods**

- **Gaussian Elimination** 
  - By using ERO, matrix A is transformed into an upper triangular matrix (all elements below diagonal 0)
  - Back substitution is used to solve the uppertriangular system

$$\begin{bmatrix} a_{11} & \cdots & a_{1i} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots & & \vdots \\ a_{i1} & \cdots & a_{ii} & \cdots & a_{in} \\ \vdots & & \vdots & \ddots & \vdots \\ a_{n1} & \cdots & a_{ni} & \cdots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_i \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_i \\ \vdots \\ b_n \end{bmatrix} \Rightarrow \begin{bmatrix} a_{11} & \cdots & a_{1i} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots & & \vdots \\ 0 & \cdots & \widetilde{a}_{ii} & \cdots & \widetilde{a}_{in} \\ \vdots & \ddots & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & \cdots & \widetilde{a}_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ \widetilde{b}_i \\ \vdots \\ \widetilde{b}_n \end{bmatrix}$$

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#### First step of elimination

#### **Pivotal element**

$$\begin{bmatrix} a_{11}^{(1)} & a_{12}^{(1)} & a_{13}^{(1)} & \cdots & a_{1n}^{(1)} \\ a_{21}^{(1)} & a_{22}^{(1)} & a_{23}^{(1)} & \cdots & a_{2n}^{(1)} \\ a_{31}^{(1)} & a_{32}^{(1)} & a_{33}^{(1)} & \cdots & a_{3n}^{(1)} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{n1}^{(1)} & a_{n2}^{(1)} & a_{n3}^{(1)} & \cdots & a_{nn}^{(1)} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_2 \\ x_3 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1^{(1)} \\ b_2^{(1)} \\ b_2^{(1)} \\ b_3^{(1)} \\ \vdots \\ b_n^{(1)} \end{bmatrix}$$

$$m_{2,1} = a_{21}^{(1)} / a_{11}^{(1)} \begin{bmatrix} a_{11}^{(1)} & a_{12}^{(1)} & a_{13}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & a_{23}^{(2)} & \cdots & a_{2n}^{(2)} \\ m_{3,1} = a_{31}^{(1)} / a_{11}^{(1)} & 0 & a_{32}^{(2)} & a_{33}^{(2)} & \cdots & a_{3n}^{(2)} \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ m_{n,1} = a_{n1}^{(1)} / a_{11}^{(1)} & 0 & a_{n2}^{(2)} & a_{n3}^{(2)} & \cdots & a_{nn}^{(2)} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1^{(1)} \\ b_2^{(2)} \\ b_2^{(2)} \\ b_3^{(2)} \\ \vdots \\ b_n^{(2)} \end{bmatrix}$$



#### Second step of elimination

Pivotal element 
$$\begin{bmatrix} a_{11}^{(1)} & a_{12}^{(1)} & a_{13}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & a_{23}^{(2)} & \cdots & a_{2n}^{(2)} \\ 0 & a_{32}^{(2)} & a_{33}^{(2)} & \cdots & a_{3n}^{(2)} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & a_{n2}^{(2)} & a_{n3}^{(2)} & \cdots & a_{nn}^{(2)} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1^{(1)} \\ b_2^{(2)} \\ b_3^{(2)} \\ \vdots \\ b_n^{(2)} \end{bmatrix}$$

$$m_{3,2} = a_{32}^{(2)} / a_{22}^{(2)} \begin{bmatrix} a_{11}^{(1)} & a_{12}^{(1)} & a_{13}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & a_{23}^{(2)} & \cdots & a_{2n}^{(2)} \\ 0 & 0 & a_{33}^{(3)} & \cdots & a_{3n}^{(3)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ m_{n,2} = a_{n2}^{(2)} / a_{22}^{(2)} \begin{bmatrix} 0 & 0 & a_{n3}^{(3)} & \cdots & a_{nn}^{(3)} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1^{(1)} \\ b_2^{(2)} \\ b_2^{(3)} \\ \vdots \\ b_n^{(3)} \end{bmatrix}$$



# Gaussion elimination algorithm

$$m_{r,p} = a_{rp}^{(p)} / a_{pp}^{(p)}$$
 $a_{rp}^{(p)} = 0$ 
 $b_r^{(p+1)} = b_r^{(p)} - m_{r,p} \times b_p^{(p)}$ 

For c=p+1 to n

$$a_{rc}^{(p+1)} = a_{rc}^{(p)} - m_{r,p} \times a_{pc}^{(p)}$$

# Back substitution algorithm

$$\begin{bmatrix} a_{11}^{(1)} & a_{12}^{(1)} & a_{13}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & a_{23}^{(2)} & \cdots & a_{2n}^{(2)} \\ 0 & 0 & a_{33}^{(3)} & \cdots & a_{3n}^{(3)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & a_{n-1n-1}^{(n)} & a_{n-1n}^{(n)} \\ 0 & 0 & 0 & 0 & a_{nn}^{(n)} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_{n-1} \\ x_n \end{bmatrix} = \begin{bmatrix} b_1^{(1)} \\ b_2^{(2)} \\ b_2^{(3)} \\ \vdots \\ b_{n-1}^{(n-1)} \\ b_n^{(n)} \end{bmatrix}$$

$$x_{n} = \frac{b_{n}^{(n)}}{a_{nn}^{(n)}} \qquad x_{n-1} = \frac{1}{a_{n-1}^{(n-1)}} \left[ b_{n-1}^{(n-1)} - a_{n-1}^{n-1} x_{n} \right]$$

$$x_{i} = \frac{1}{a_{ii}^{(i)}} \left[ b_{i}^{(i)} - \sum_{k=i+1}^{n} a_{ik}^{(i)} x_{k} \right] \quad i = n-1, n-2, \dots, 1$$



#### Operation count

- Number of arithmetic operations required by the algorithm to complete its task.
- Generally only multiplications and divisions are counted
- Elimination process  $\left(\frac{n^3}{3} + \frac{n^2}{2} \frac{5n}{6}\right)$
- Back substitution  $\frac{n^2 + n}{2}$
- Total  $\frac{n^3}{3} + n^2 \frac{n}{3}$

Dominates
Not efficient for
different RHS vectors