

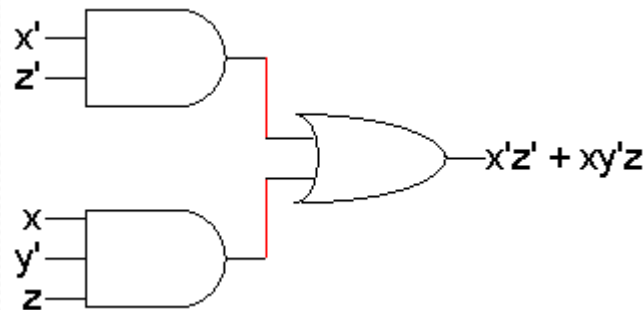
**Gate-level minimization:**  
The map method up to four variable  
don't care conditions  
POS simplification

## Lecture 3

Dronacharya Group of Institutions

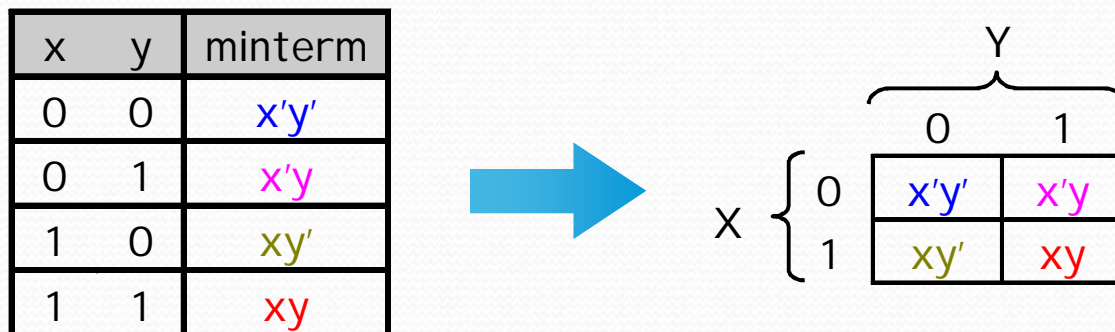
# Karnaugh Maps

- Boolean algebra helps us simplify expressions and circuits
- Karnaugh Map: A graphical technique for simplifying a Boolean expression into either form:
  - minimal sum of products (MSP)
  - minimal product of sums (MPS)
- Goal of the simplification.
  - There are a minimal number of product/sum terms
  - Each term has a minimal number of literals
- Circuit-wise, this leads to a *minimal* two-level implementation

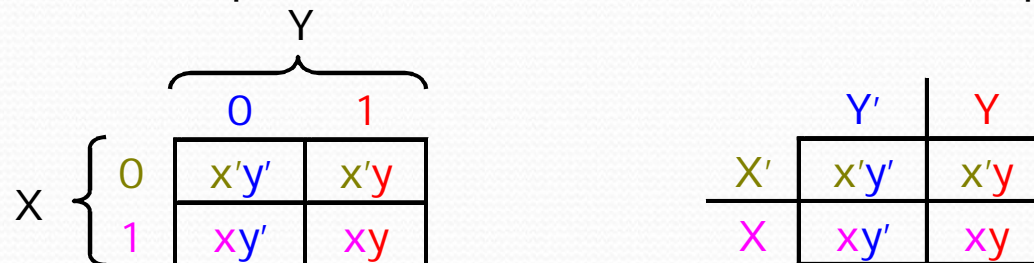


# Re-arranging the Truth Table

- A two-variable function has four possible minterms. We can re-arrange these minterms into a **Karnaugh map**



- Now we can easily see which minterms contain common literals
  - Minterms on the left and right sides contain  $y'$  and  $y$  respectively
  - Minterms in the top and bottom rows contain  $x'$  and  $x$  respectively





# Karnaugh Map Simplifications

- Imagine a two-variable sum of minterms:

$$x'y' + x'y$$

- Both of these minterms appear in the top row of a Karnaugh map, which means that they both contain the literal  $x'$

		Y
	$x'y'$	$x'y$
X	$xy'$	$xy$

- What happens if you simplify this expression using Boolean algebra?

$$\begin{aligned}x'y' + x'y &= x'(y' + y) && [ \text{Distributive} ] \\ &= x' \bullet 1 && [ y + y' = 1 ] \\ &= x' && [ x \bullet 1 = x ]\end{aligned}$$

# More Two-Variable Examples

- Another example expression is  $x'y + xy$ 
  - Both minterms appear in the right side, where  $y$  is uncomplemented
  - Thus, we can reduce  $x'y + xy$  to just  $y$

		Y
X	$x'y'$	$x'y$
	$xy'$	$xy$

- How about  $x'y' + x'y + xy$ ?
  - We have  $x'y' + x'y$  in the top row, corresponding to  $x'$
  - There's also  $x'y + xy$  in the right side, corresponding to  $y$
  - This whole expression can be reduced to  $x' + y$

		Y
X	$x'y'$	$x'y$
	$xy'$	$xy$



# A Three-Variable Karnaugh Map

- For a three-variable expression with inputs  $x, y, z$ , the arrangement of minterms is more tricky:

		YZ			
		00	01	11	10
X	0	$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
	1	$xy'z'$	$xy'z$	$xyz$	$xyz'$

		YZ			
		00	01	11	10
X	0	$m_0$	$m_1$	$m_3$	$m_2$
	1	$m_4$	$m_5$	$m_7$	$m_6$

- Another way to label the K-map (use whichever you like):

		Y			
		$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
X		$xy'z'$	$xy'z$	$xyz$	$xyz'$
	Z				

		Y			
		$m_0$	$m_1$	$m_3$	$m_2$
X		$m_4$	$m_5$	$m_7$	$m_6$
	Z				

# Why the funny ordering?

- With this ordering, any group of 2, 4 or 8 adjacent squares on the map contains common literals that can be factored out

			Y	
	$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
X	$xy'z'$	$xy'z$	$xyz$	$xyz'$
		Z		

$$\begin{aligned}
 & \textcolor{red}{x'y'z} + \textcolor{red}{x'yz} \\
 &= x'z(y' + y) \\
 &= x'z \bullet 1 \\
 &= x'z
 \end{aligned}$$

- "Adjacency" includes wrapping around the left and right sides:

			Y	
	$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
X	$xy'z'$	$xy'z$	$xyz$	$xyz'$
		Z		

$$\begin{aligned}
 & \textcolor{blue}{x'y'z'} + \textcolor{blue}{xy'z'} + \textcolor{blue}{x'yz'} + \textcolor{blue}{xyz'} \\
 &= z'(x'y' + xy' + x'y + xy) \\
 &= z'(y'(x' + x) + y(x' + x)) \\
 &= z'(y' + y) \\
 &= z'
 \end{aligned}$$

- We'll use this property of adjacent squares to do our simplifications.



# K-maps From Truth Tables

- We can fill in the K-map directly from a truth table
  - The output in row  $i$  of the table goes into square  $m_i$  of the K-map
  - Remember that the rightmost columns of the K-map are "switched"

x	y	z	$f(x,y,z)$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0

			Y	
	$m_0$	$m_1$	$m_3$	$m_2$
X	$m_4$	$m_5$	$m_7$	$m_6$
		Z		

			Y	
	0	1	0	0
X	0	1	1	1
		Z		

1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1



# Reading the MSP from the K-map

- You can find the minimal SoP expression
  - Each rectangle corresponds to one product term
  - The product is determined by finding the common literals in that rectangle

				Y
	0	1	0	0
X	0	1	1	1
		Z		

				Y
	$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
X	$xy'z'$	$xy'z$	$xyz$	$xyz'$
		Z		

$y'z$                        $xy$

$$F(x,y,z) = y'z + xy$$

# Grouping the Minterms Together

- The most difficult step is grouping together all the 1s in the K-map
  - Make **rectangles** around groups of one, two, four or eight 1s
  - All of the 1s in the map should be included in at least one rectangle
  - Do *not* include any of the 0s
  - Each group corresponds to one product term

			Y	
	0	1	0	0
X	0	1	1	1
		Z		



# For the Simplest Result

- Make as few rectangles as possible, to minimize the number of products in the final expression.
- Make each rectangle as large as possible, to minimize the number of literals in each term.
- Rectangles can be overlapped, if that makes them larger.



# K-map Simplification of SoP Expressions

- Let's consider simplifying  $f(x,y,z) = xy + y'z + xz$
- You should convert the expression into a sum of minterms form,
  - The easiest way to do this is to make a truth table for the function, and then read off the minterms
  - You can either write out the literals or use the minterm shorthand
- Here is the truth table and sum of minterms for our example:

x	y	z	$f(x,y,z)$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$\begin{aligned}f(x,y,z) &= x'y'z + xy'z + xyz' + \\ &\quad xyz \\ &= m_1 + m_5 + m_6 + \\ &\quad m_7\end{aligned}$$

# Unsimplifying Expressions

- You can also convert the expression to a sum of minterms with Boolean algebra
  - Apply the distributive law in reverse to add in missing variables.
  - Very few people actually do this, but it's occasionally useful.

$$\begin{aligned}xy + y'z + xz &= (xy \bullet 1) + (y'z \bullet 1) + (xz \bullet 1) \\&= (xy \bullet (z' + z)) + (y'z \bullet (x' + x)) + (xz \bullet (y' + y)) \\&= (xyz' + xyz) + (x'y'z + xy'z) + (xy'z + xyz) \\&= xyz' + xyz + x'y'z + xy'z \\&= m_1 + m_5 + m_6 + m_7\end{aligned}$$

- In both cases, we're actually "unsimplifying" our example expression
  - The resulting expression is larger than the original one!
  - But having all the individual minterms makes it easy to combine them together with the K-map



# Making the Example K-map

- In our example, we can write  $f(x,y,z)$  in two equivalent ways

$$f(x,y,z) = x'y'z + xy'z + xyz' + xyz$$

$$f(x,y,z) = m_1 + m_5 + m_6 + m_7$$

			Y	
	$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
X	$xy'z'$	$xy'z$	$xyz$	$xyz'$
		Z		

			Y	
	$m_0$	$m_1$	$m_3$	$m_2$
X	$m_4$	$m_5$	$m_7$	$m_6$
		Z		

- In either case, the resulting K-map is shown below

			Y	
	0	1	0	0
X	0	1	1	1
		Z		



# Practice K-map 1

- Simplify the sum of minterms  $m_1 + m_3 + m_5 + m_6$

			Y
X			
	Z		

			Y	
	m <sub>0</sub>	m <sub>1</sub>	m <sub>3</sub>	m <sub>2</sub>
X	m <sub>4</sub>	m <sub>5</sub>	m <sub>7</sub>	m <sub>6</sub>
		Z		

# Solutions for Practice K-map 1

- Here is the filled in K-map, with all groups shown
  - The magenta and green groups overlap, which makes each of them as large as possible
  - Minterm  $m_6$  is in a group all by its lonesome

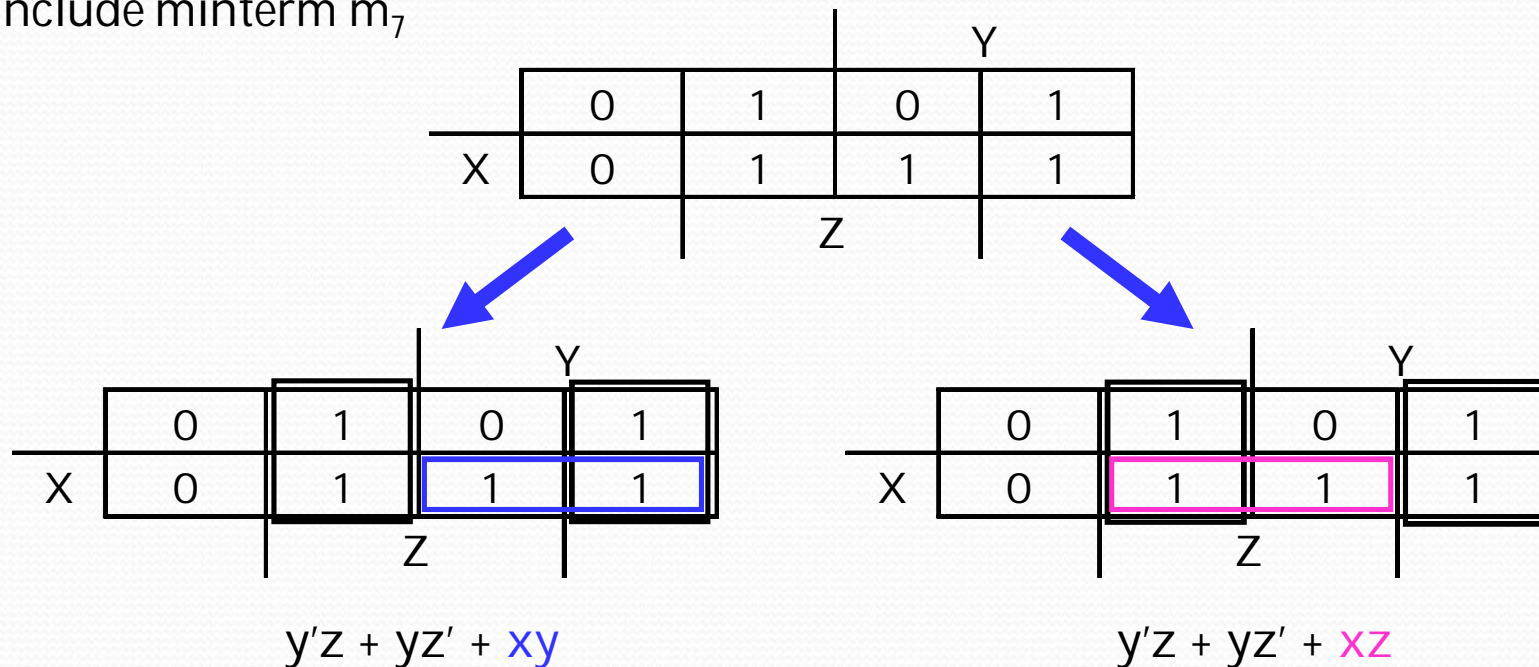
				Y
	0	1	1	0
X	0	1	0	1
		Z		

- The final MSP here is  $x'z + y'z + xyz'$



# K-maps can be tricky!

- There may not necessarily be a *unique* MSP. The K-map below yields two valid and equivalent MSPs, because there are two possible ways to include minterm  $m_7$

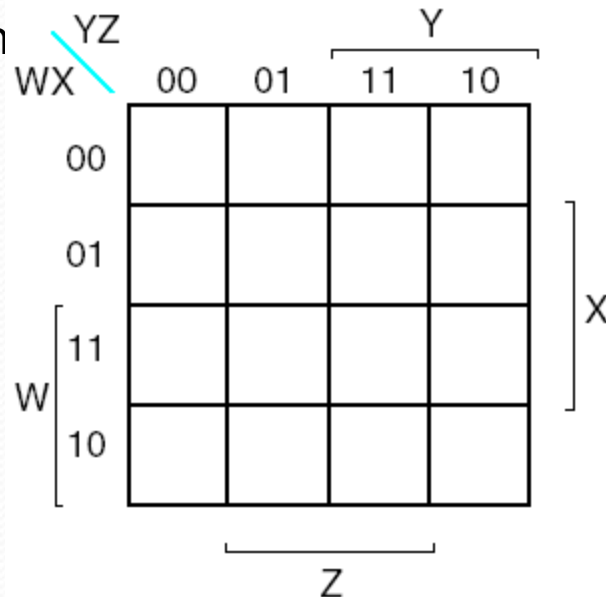


- Remember that overlapping groups is possible, as shown above



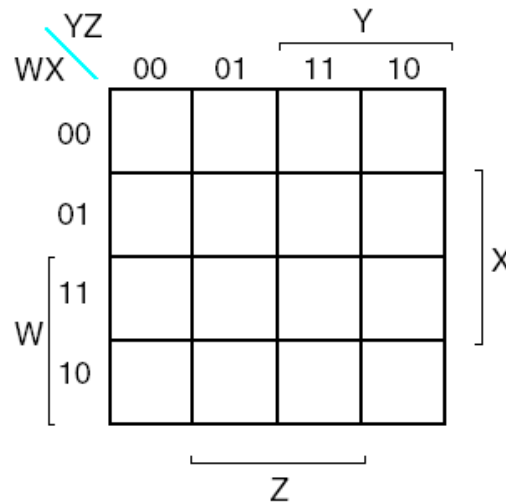
# Four-variable K-maps – $f(W,X,Y,Z)$

- We can do four-variable expressions too!
  - The minterms in the third and fourth columns, *and* in the third and fourth rows, are switched around.
  - Again, this ensures that adjacent cells share common literals



- Grouping minterms is similar to the three-variable case, but:
  - You can have rectangular groups of 1, 2, 4, 8 or 16 minterms
  - You can wrap around *all four* sides

# Four-variable K-maps



		Y		
		w'x'y'z'	w'x'y'z	w'x'yz
		w'xy'z'	w'xy'z	w'xyz
W	wxy'z'	wxy'z	wxyz	wxyz'
	wx'y'z'	wx'y'z	wx'yz	wx'yz'
		Z		X

		Y		
		m <sub>0</sub>	m <sub>1</sub>	m <sub>3</sub>
		m <sub>4</sub>	m <sub>5</sub>	m <sub>7</sub>
W	m <sub>12</sub>	m <sub>13</sub>	m <sub>15</sub>	m <sub>14</sub>
	m <sub>8</sub>	m <sub>9</sub>	m <sub>11</sub>	m <sub>10</sub>
		Z		X

# Example: Simplify $m_0 + m_2 + m_5 + m_8 + m_{10} + m_{13}$

- The expression is already a sum of minterms, so here's the K-map:

--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--

		Y			
		$m_0$	$m_1$	$m_3$	$m_2$
		$m_4$	$m_5$	$m_7$	$m_6$
W		$m_{12}$	$m_{13}$	$m_{15}$	$m_{14}$
		$m_8$	$m_9$	$m_{11}$	$m_{10}$
		Z			

- We can make the following groups, resulting in the MSP  $x'z' + xy'z$

		Y			
		1	0	0	1
		0	1	0	0
W		0	1	0	0
		1	0	0	1
		Z			

				Y	
		w'x'y'z'	w'x'y'z	w'x'yz	w'x'yz'
		w'xy'z'	w'xy'z	w'xyz	w'xyz'
W		wxy'z'	wxy'z	wxyz	wxyz'
		wx'y'z'	wx'y'z	wx'yz	wx'yz'
		Z			



# PoS Optimization

- Maxterms are grouped to find minimal PoS expression

		yz			
		00	01	11	10
x	0	$x + y + z$	$x + y + z'$	$x + y' + z'$	$x + y' + z$
	1	$x' + y + z$	$x' + y + z'$	$x' + y' + z'$	$x' + y' + z$

# PoS Optimization

•  $F(W,X,Y,Z) = \prod M(0,1,2,4,5)$

		00	01	yz	11		10
x	0	$x + y + z$	$x + y + z'$		$x + y' + z'$		$x + y' + z$
	1	$x' + y + z$	$x' + y + z'$		$x' + y' + z'$		$x' + y' + z$

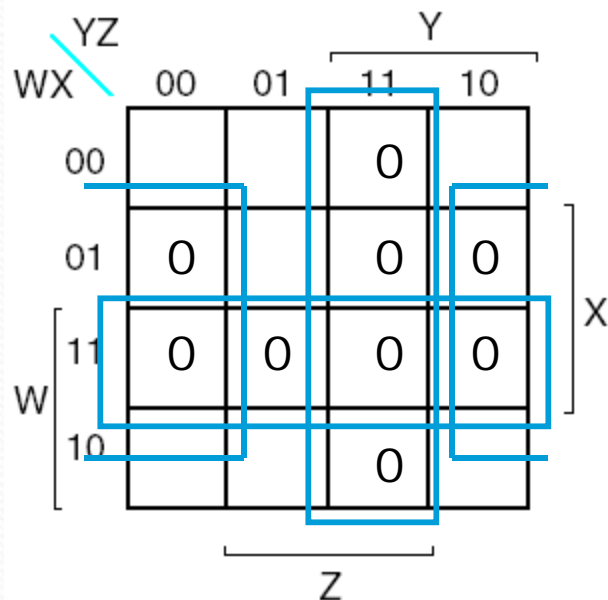
		00	01	yz	11		10
x	0	0	0		1		0
	1	0	0		1		1

$$F(W,X,Y,Z) = Y \cdot (X + Z)$$

# PoS Optimization from SoP

$$F(W,X,Y,Z) = \sum m(0,1,2,5,8,9,10)$$

$$= \prod M(3,4,6,7,11,12,13,14,15)$$



$$F(W,X,Y,Z) = (W' + X')(Y' + Z')(X' + Z)$$

Or,

$$F(W,X,Y,Z) = X'Y' + X'Z' + W'Y'Z$$

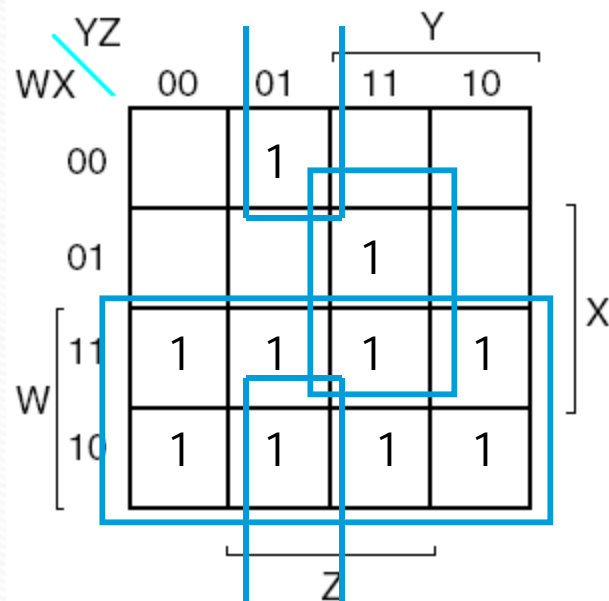
Which one is the minimal one?



# SoP Optimization from PoS

$$F(W,X,Y,Z) = \prod M(0,2,3,4,5,6)$$

$$= \sum m(1,7,8,9,10,11,12,13,14,15)$$



$$F(W,X,Y,Z) = W + XYZ + X'Y'Z$$

# I don't care!

- You don't always need all  $2^n$  input combinations in an n-variable function
  - If you can guarantee that certain input combinations never occur
  - If some outputs aren't used in the rest of the circuit
- We mark don't-care outputs in truth tables and K-maps with Xs.

x	y	z	f(x,y,z)
0	0	0	0
0	0	1	1
0	1	0	X
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	X
1	1	1	1

- Within a K-map, each X can be considered as either 0 or 1. You should pick the interpretation that allows for the most simplification.

# Practice K-map

- Find a MSP for

$$f(w,x,y,z) = \sum m(0,2,4,5,8,14,15), d(w,x,y,z) = \sum m(7,10,13)$$

This notation means that input combinations  $wxyz = 0111, 1010$  and  $1101$  (corresponding to minterms  $m_7, m_{10}$  and  $m_{13}$ ) are unused.

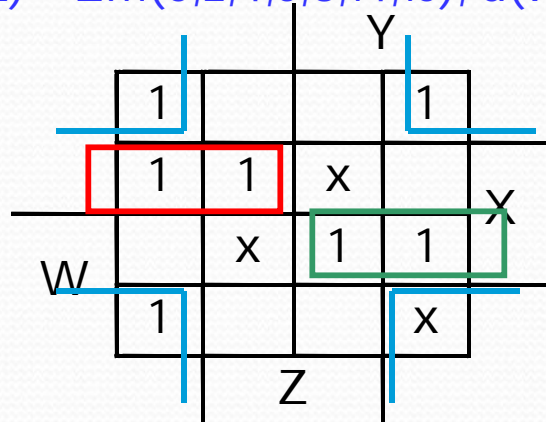
		Y		X
		1	0	
W	1	0	0	1
	1	1	x	0
	0	x	1	1
	1	0	0	x
		Z		



# Solutions for Practice K-map

- Find a MSP for:

$$f(w,x,y,z) = \sum m(0,2,4,5,8,14,15), d(w,x,y,z) = \sum m(7,10,13)$$



$$f(w,x,y,z) = x'z' + w'xy' + wxy$$

# K-map Summary

- K-maps are an alternative to algebra for simplifying expressions
  - The result is a MSP/MPS, which leads to a minimal two-level circuit
  - It's easy to handle don't-care conditions
  - K-maps are really only good for manual simplification of small expressions...
- Things to keep in mind:
  - Remember the correct order of minterms/maxterms on the K-map
  - When grouping, you can wrap around all sides of the K-map, and your groups can overlap
  - Make as few rectangles as possible, but make each of them as large as possible. This leads to fewer, but simpler, product terms
  - There may be more than one valid solution