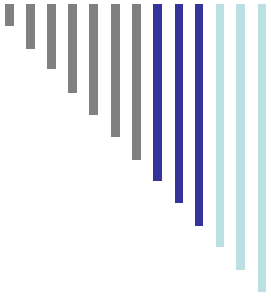




Computer Networks

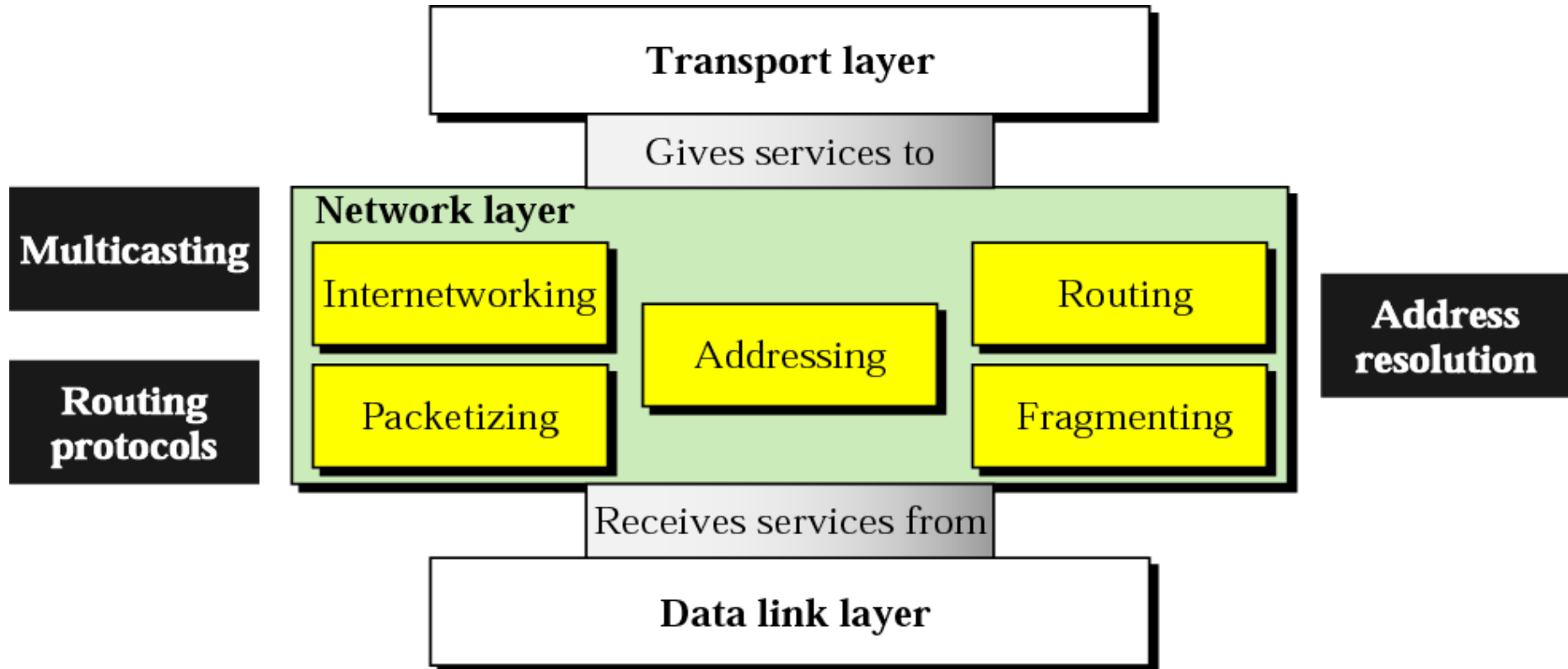
— Unit - 3 —

— Network Layer —

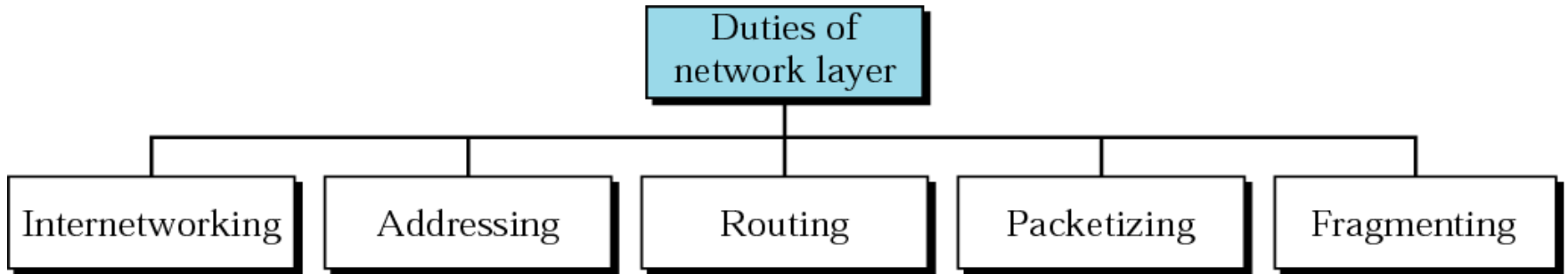


Network Layer

Position of network layer



Network layer duties



*Host-to-Host
Delivery:
Internetworking,
Addressing,
and Routing*

19.1 Internetworks

Need For Network Layer

Internet As A Packet-Switched Network

Internet As A Connectionless Network

Figure 19.1 Internetwork

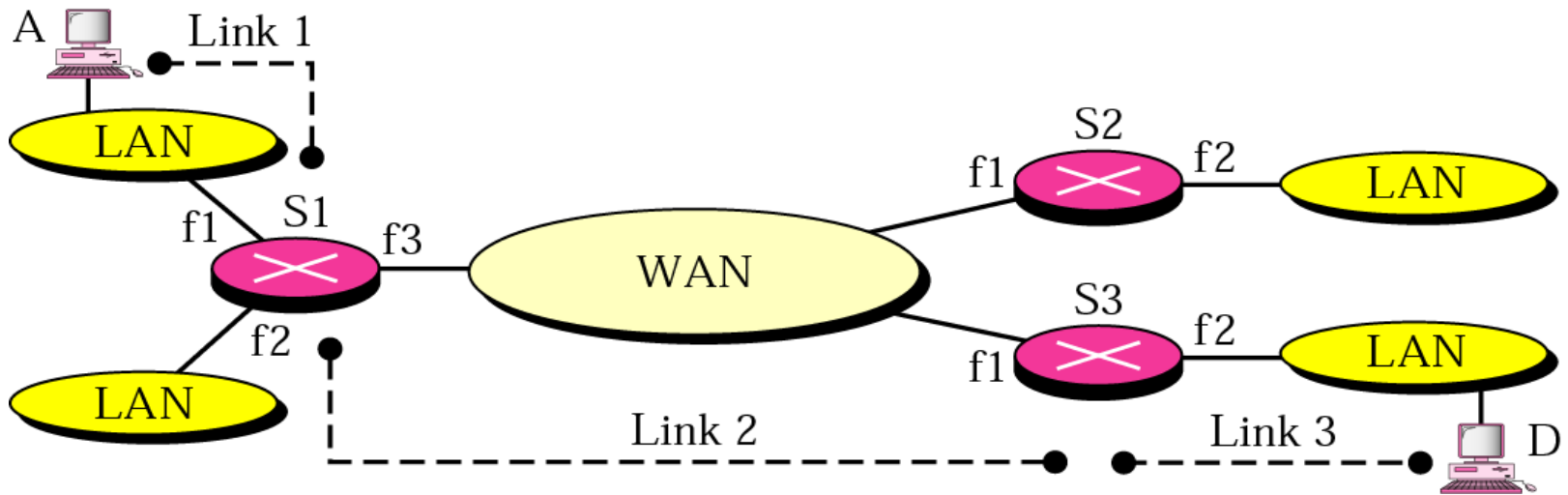


Figure 19.2 Links in an internetwork

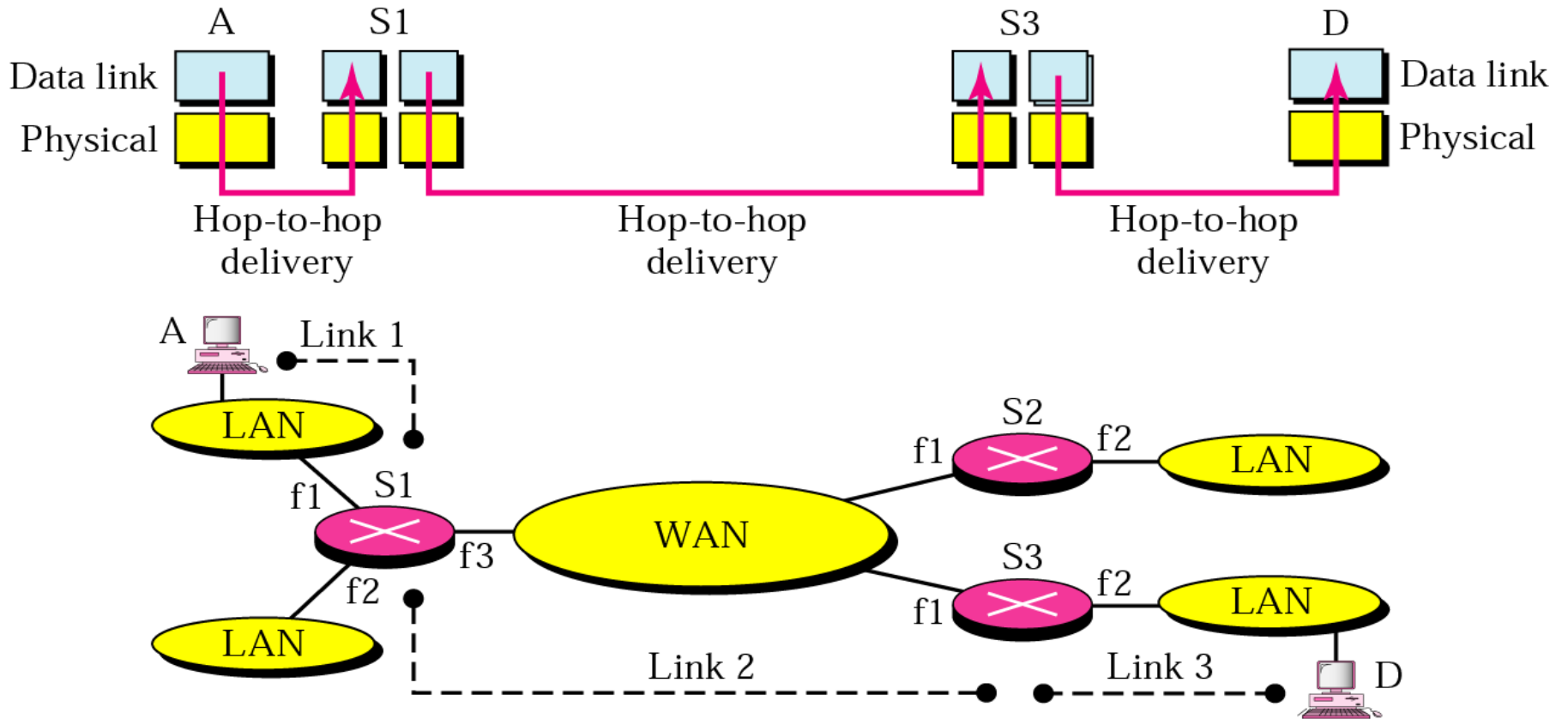


Figure 19.3 Network layer in an internetwork

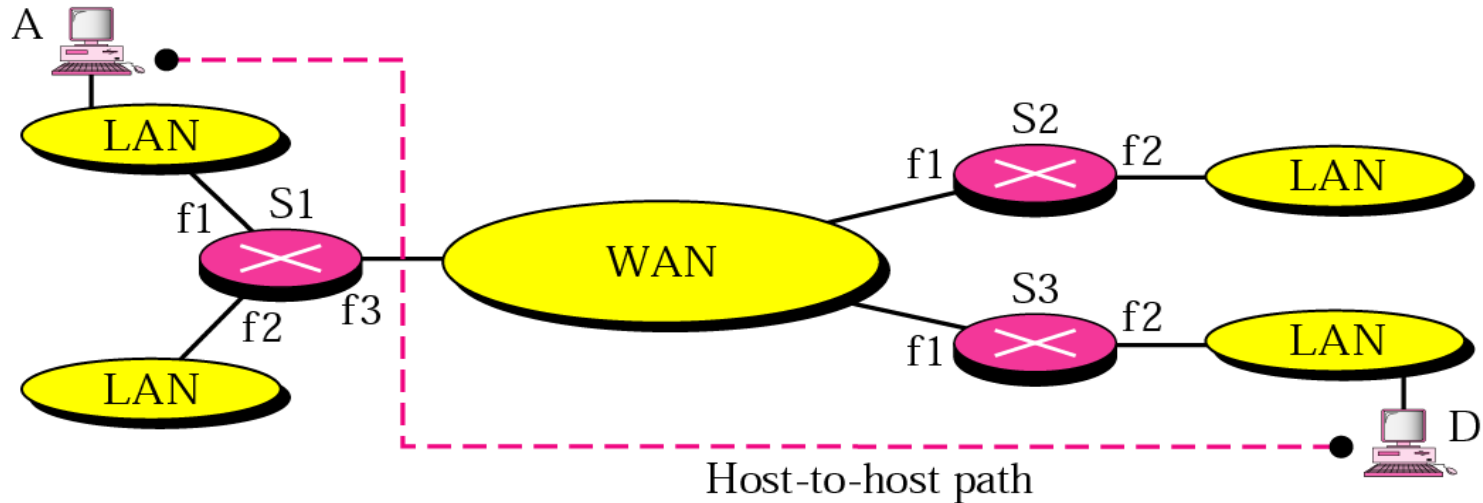


Figure 19.4 Network layer at the source

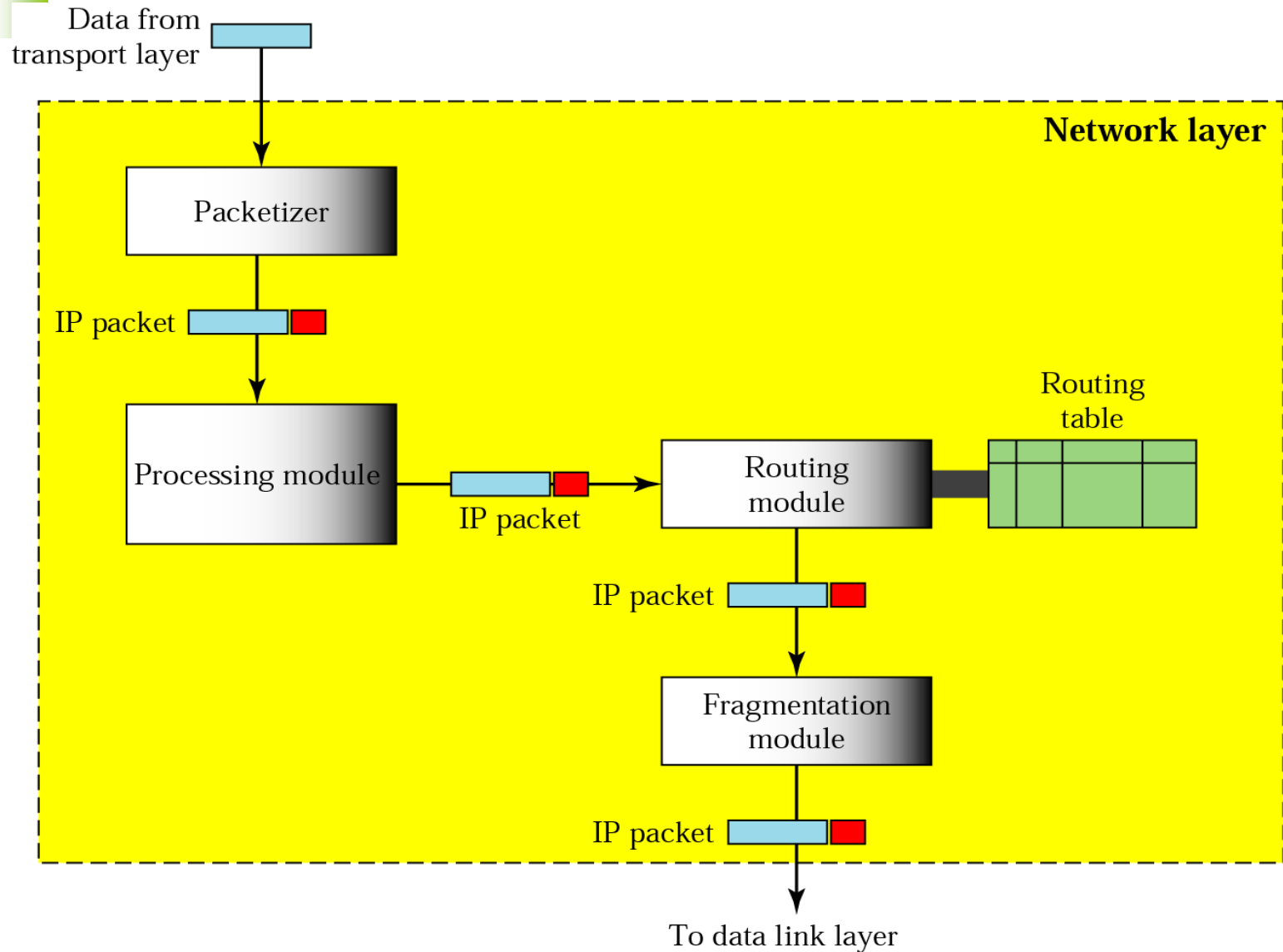


Figure 19.5 Network layer at a router

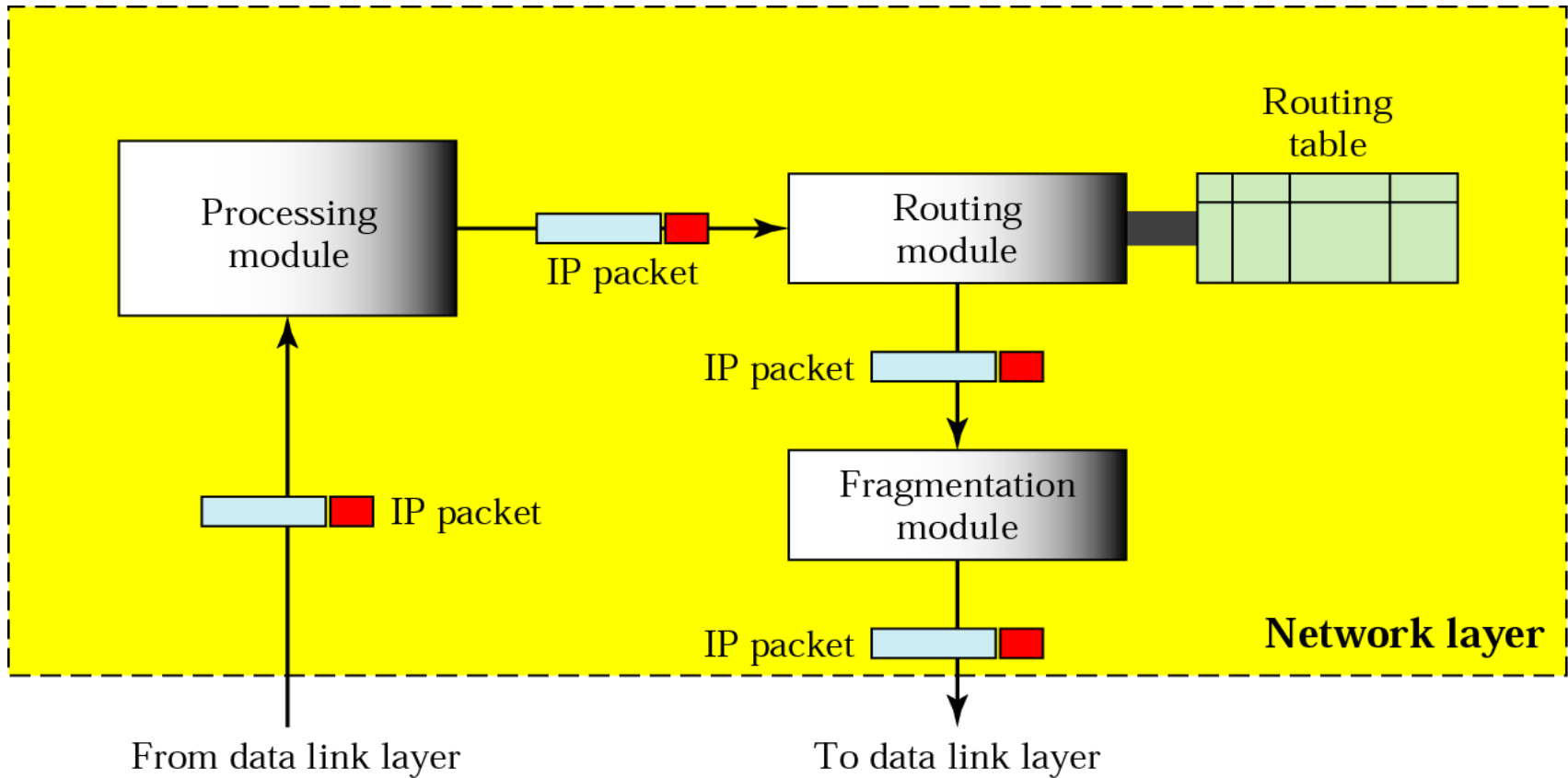
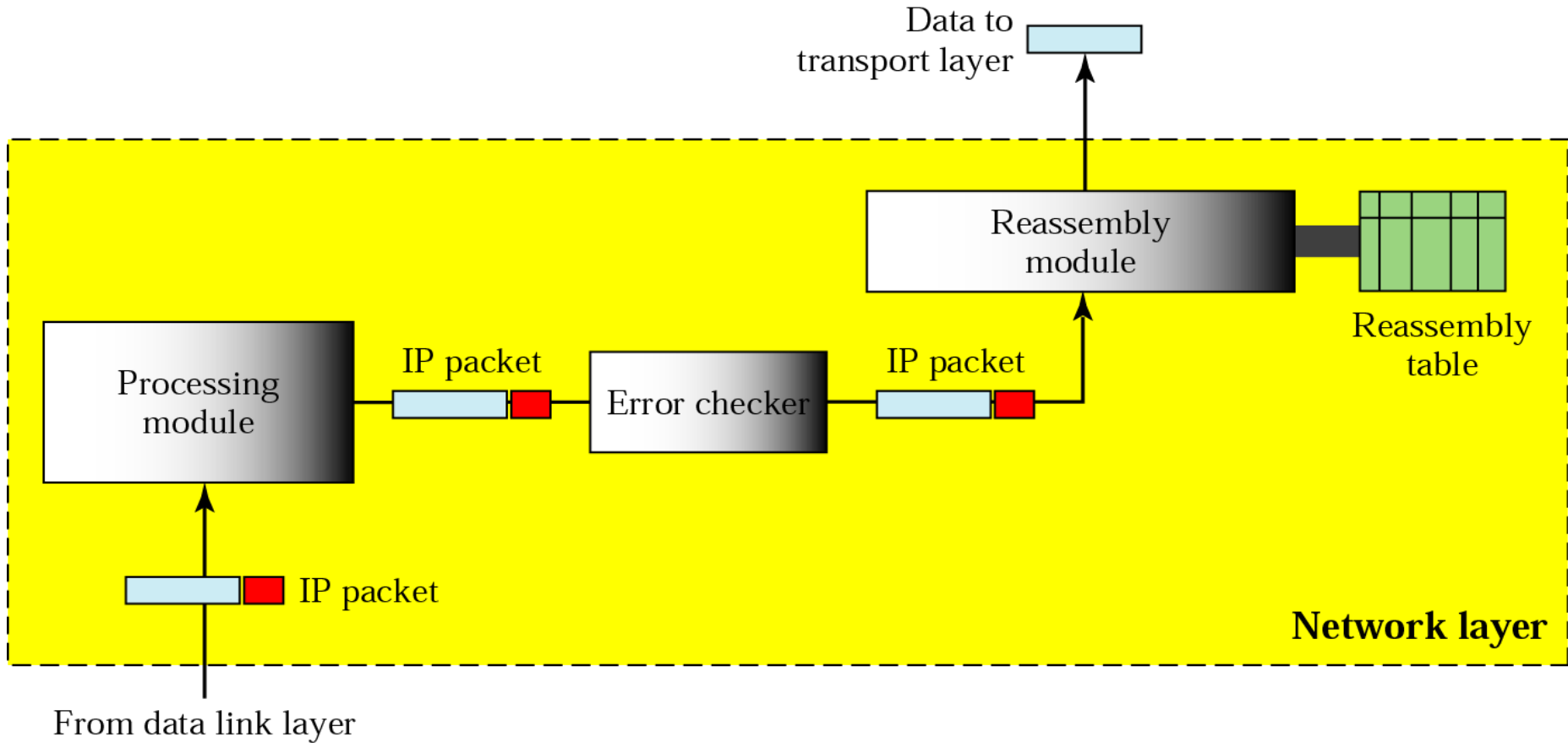


Figure 19.6 Network layer at the destination



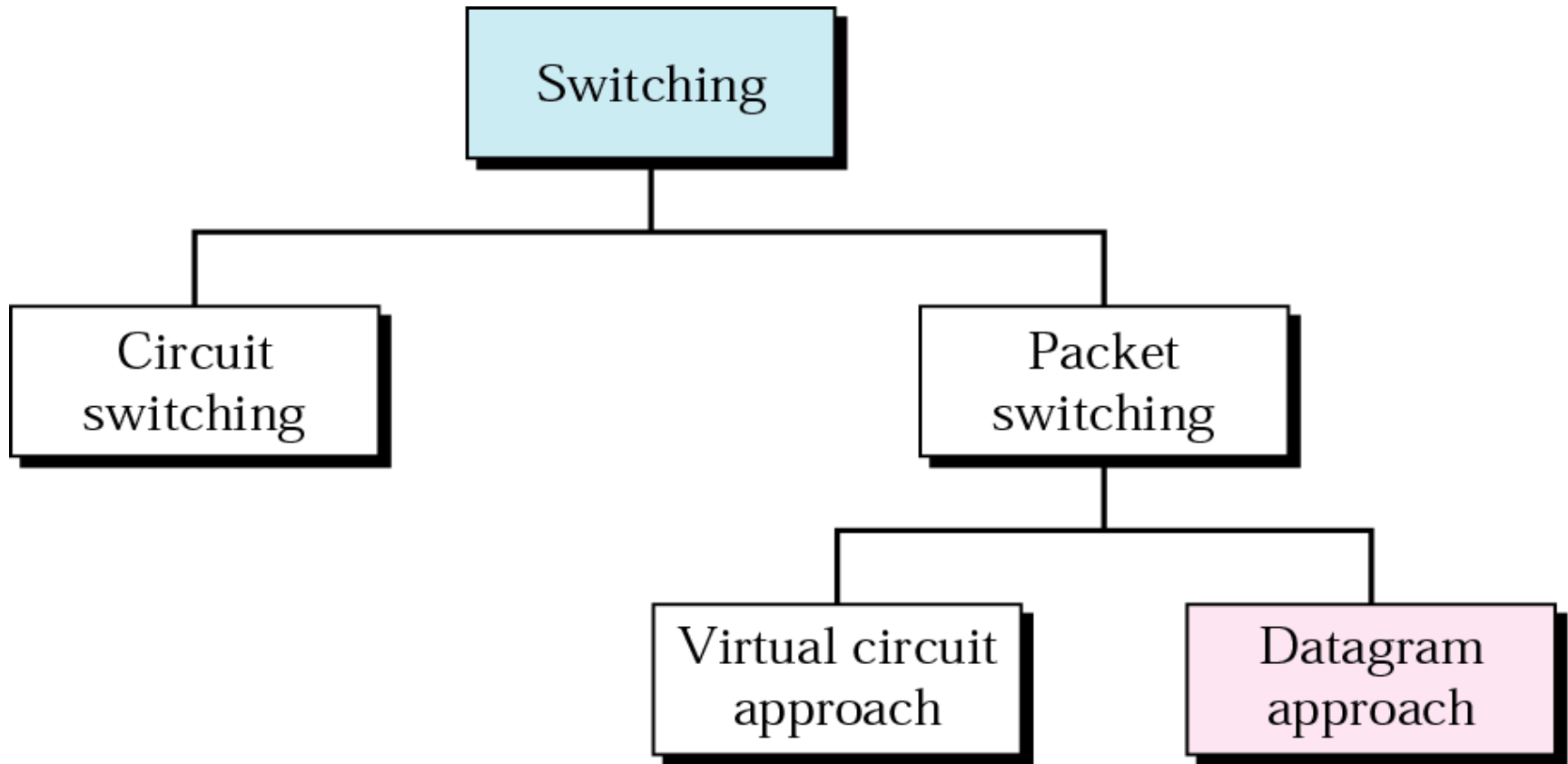
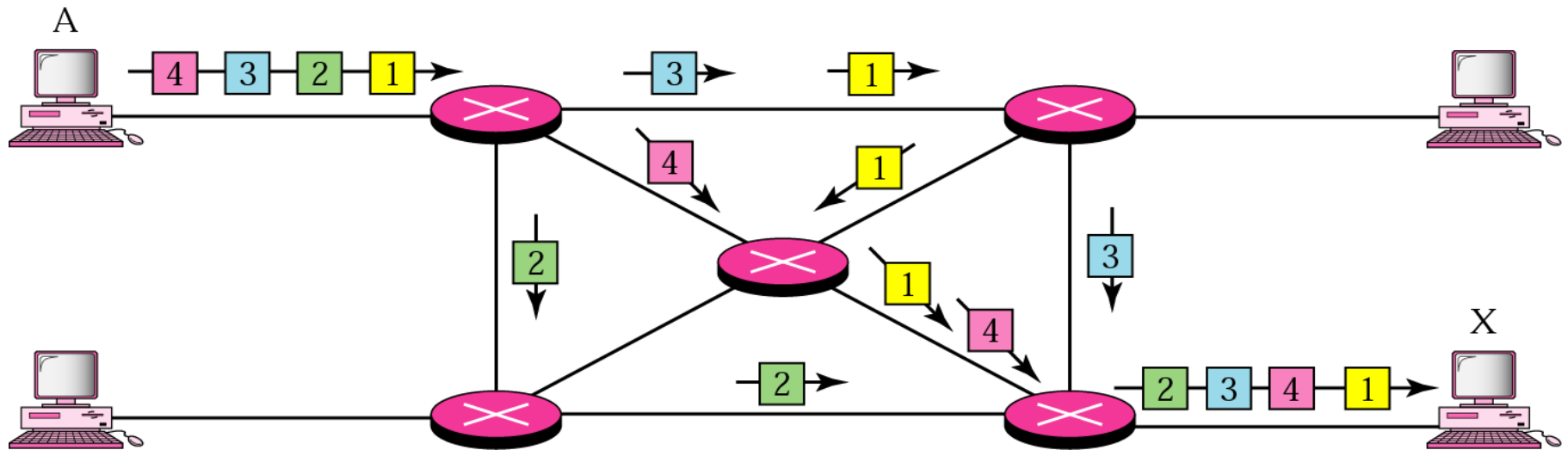
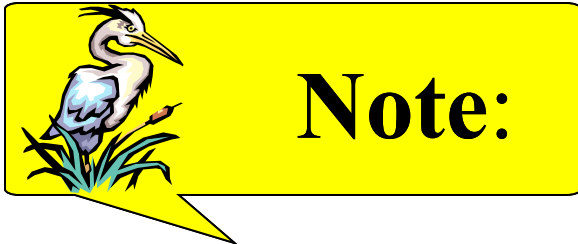


Figure 19.8 Datagram approach





Switching at the network layer in the Internet is done using the datagram approach to packet switching.



*Communication at the network layer
in the Internet is connectionless.*

19.2 Addressing

Internet Address

Classful Addressing

Subnetting

Supernetting

Classless Addressing

Dynamic Address Configuration

Network Address Translation

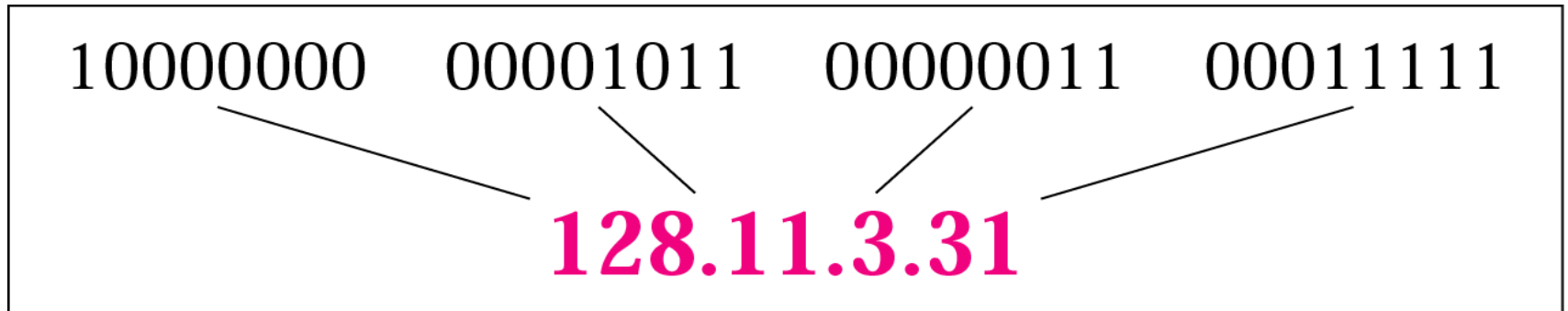


An IP address is a 32-bit address.



*The IP addresses are unique
and universal.*

Figure 19.9 Dotted-decimal notation





The binary, decimal, and hexadecimal number systems are reviewed in Appendix B.

Example 1

Change the following IP addresses from binary notation to dotted-decimal notation.

- a. 10000001 00001011 00001011 11101111
- b. 11111001 10011011 11111011 00001111

Solution

We replace each group of 8 bits with its equivalent decimal number (see Appendix B) and add dots for separation:

- a. **129.11.11.239**
- b. **249.155.251.15**

Example 2

Change the following IP addresses from dotted-decimal notation to binary notation.

- a. 111.56.45.78
- b. 75.45.34.78

Solution

We replace each decimal number with its binary equivalent (see Appendix B):

- a. **01101111 00111000 00101101 01001110**
- b. **01001011 00101101 00100010 01001110**

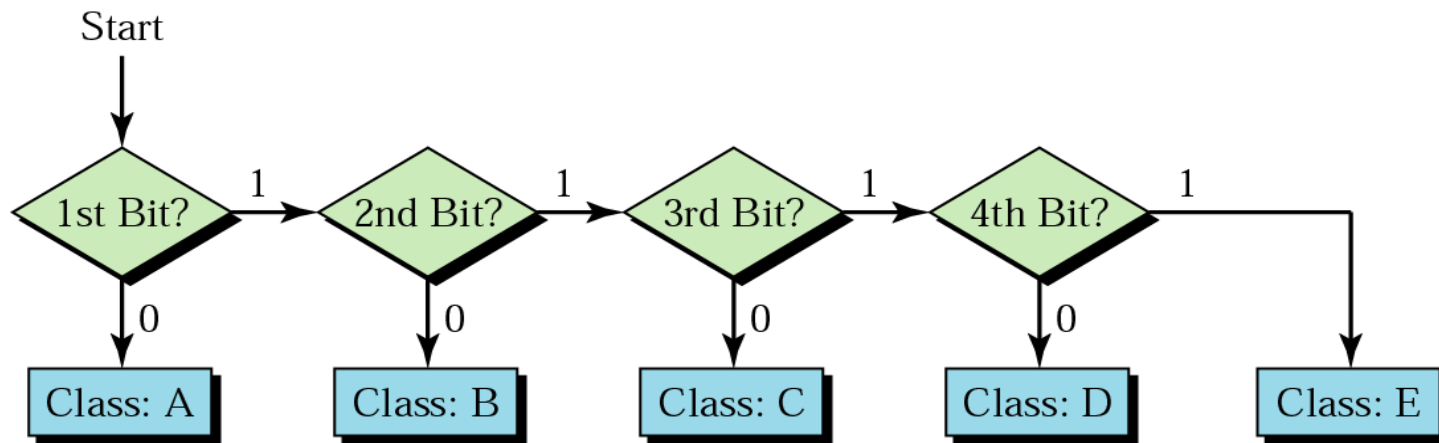


In classful addressing, the address space is divided into five classes: A, B, C, D, and E.

Figure 19.10 Finding the class in binary notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

Figure 19.11 Finding the address class



Example 3

Find the class of each address:

- a. **00000001 00001011 00001011 11101111**
- b. **11110011 10011011 11111011 00001111**

Solution

See the procedure in **Figure 19.11**.

- a. **The first bit is 0; this is a class A address.**
- b. **The first 4 bits are 1s; this is a class E address.**

Figure 19.12 Finding the class in decimal notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0 to 127			
Class B	128 to 191			
Class C	192 to 223			
Class D	224 to 239			
Class E	240 to 255			

Example 4

Find the class of each address:

- a. **227.12.14.87**
- b. **252.5.15.111**
- c. **134.11.78.56**

Solution

- a. **The first byte is 227 (between 224 and 239); the class is D.**
- b. **The first byte is 252 (between 240 and 255); the class is E.**
- c. **The first byte is 134 (between 128 and 191); the class is B.**

Figure 19.13 Netid and hostid

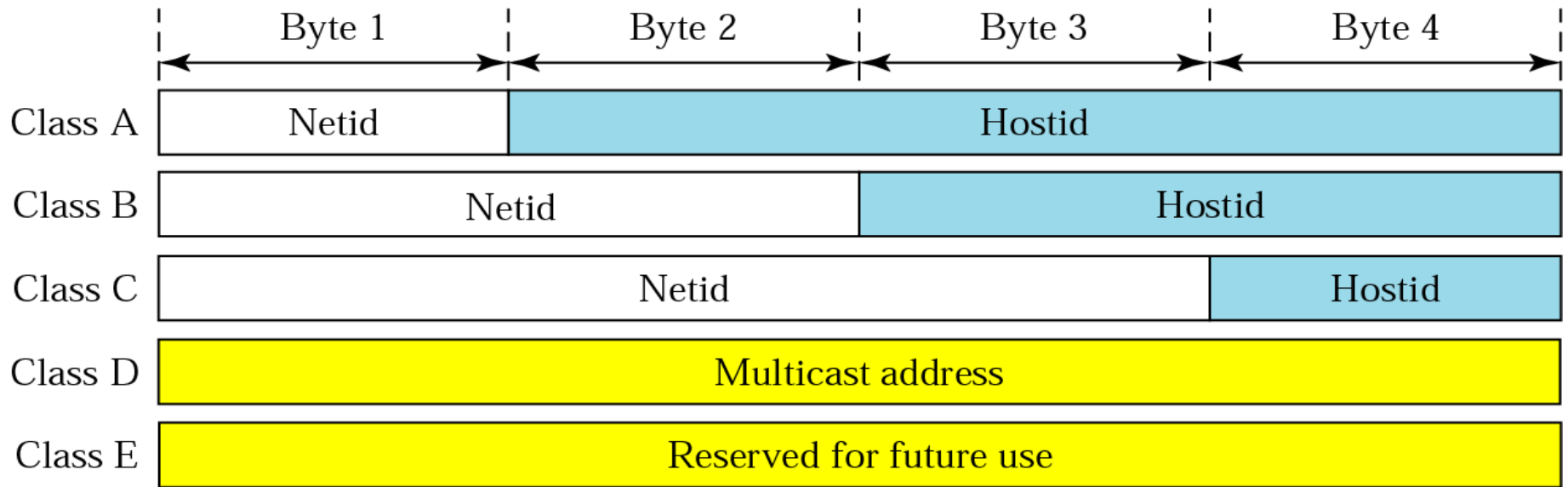
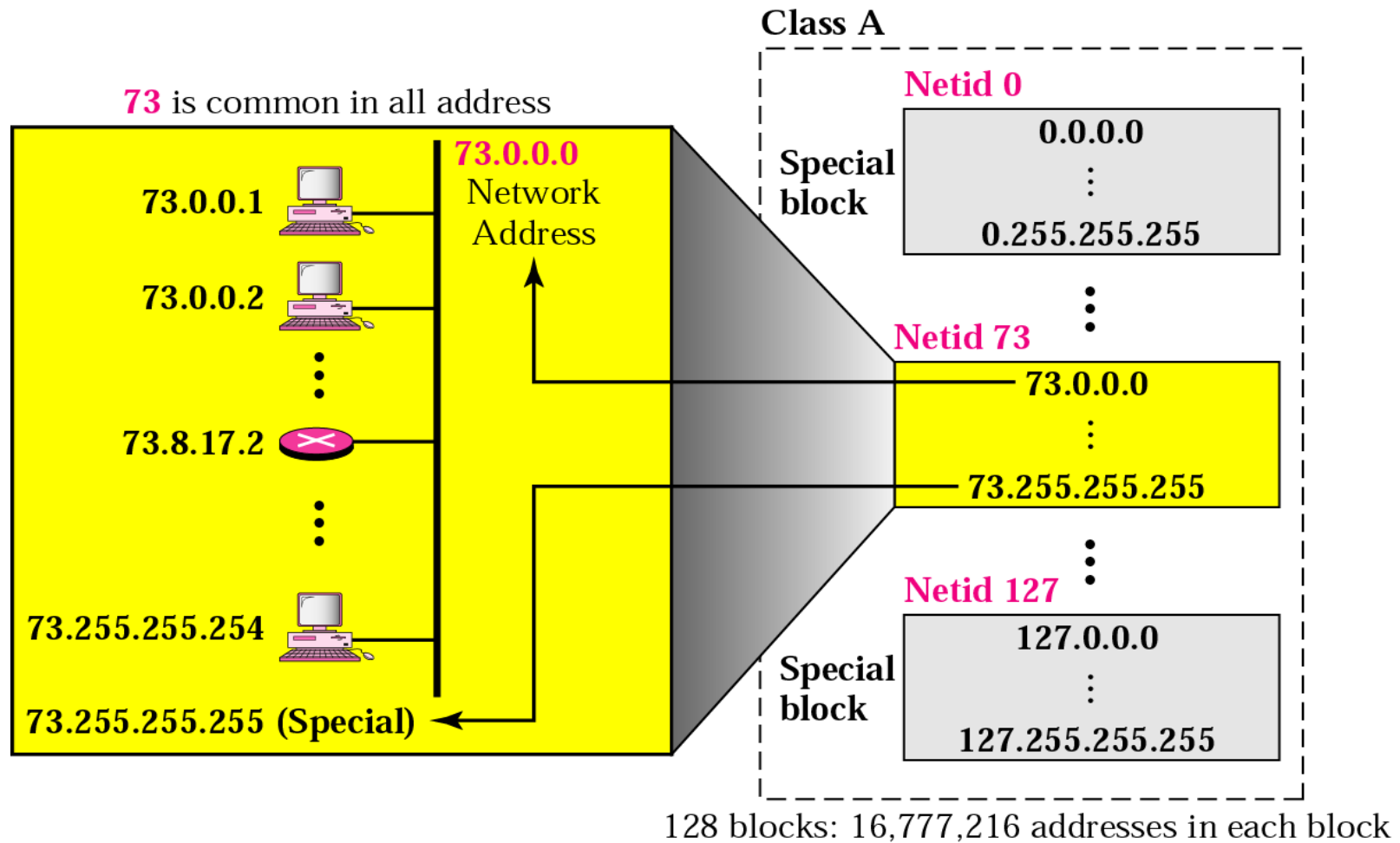


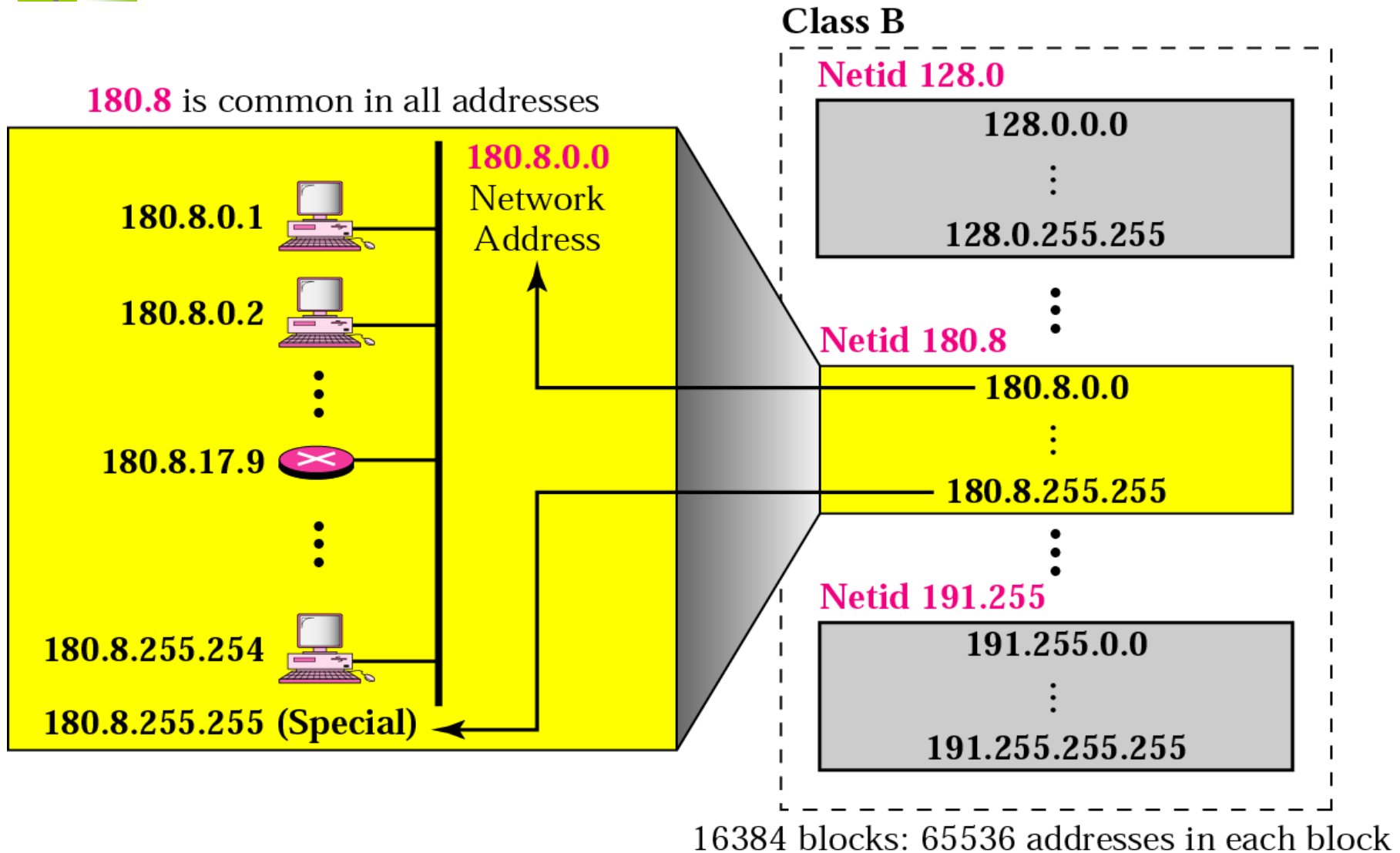
Figure 19.14 Blocks in class A





*Millions of class A addresses are
wasted.*

Figure 19.15 Blocks in class B





Many class B addresses are wasted.



The number of addresses in class C is smaller than the needs of most organizations.

Figure 19.16 Blocks in class C

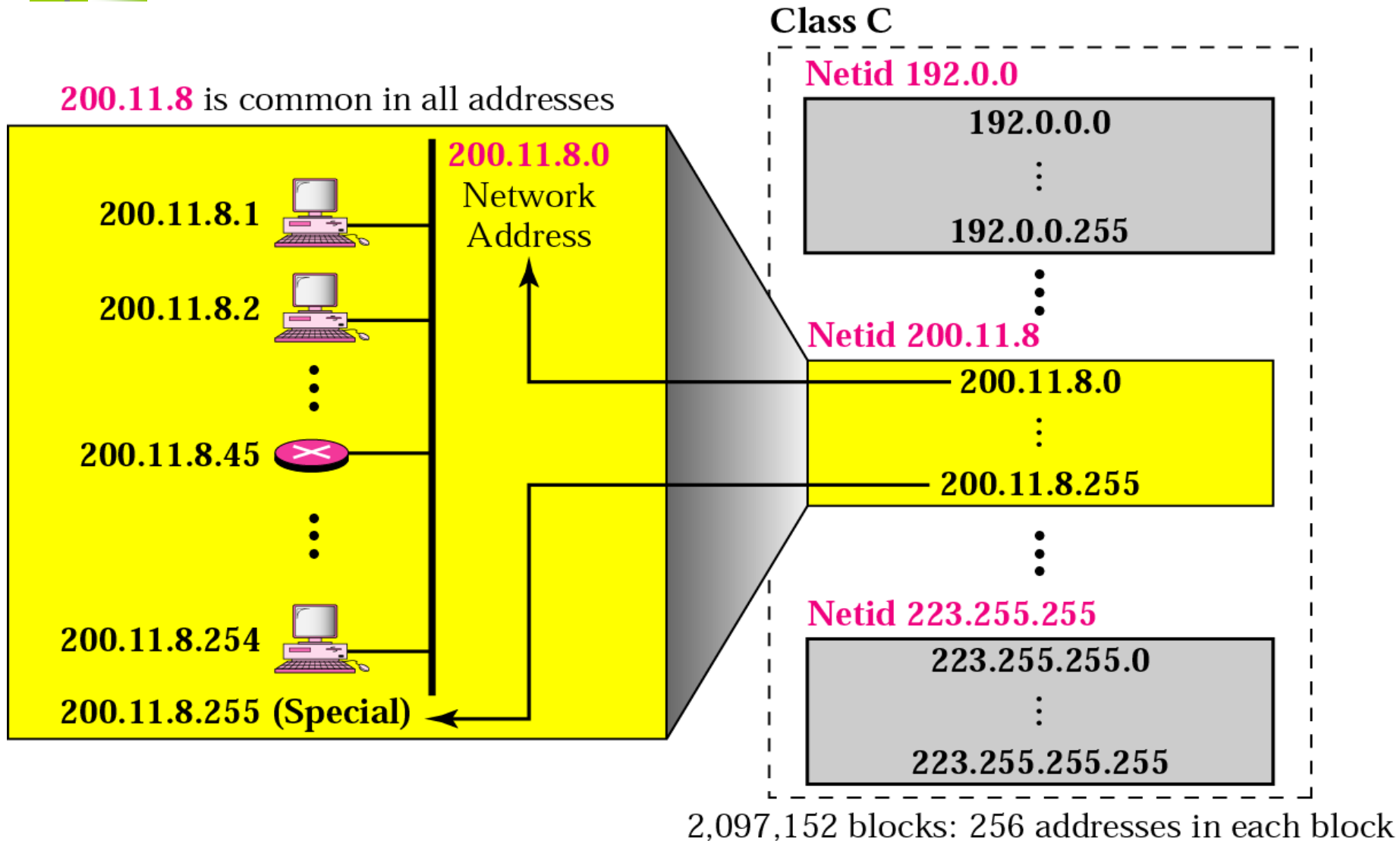
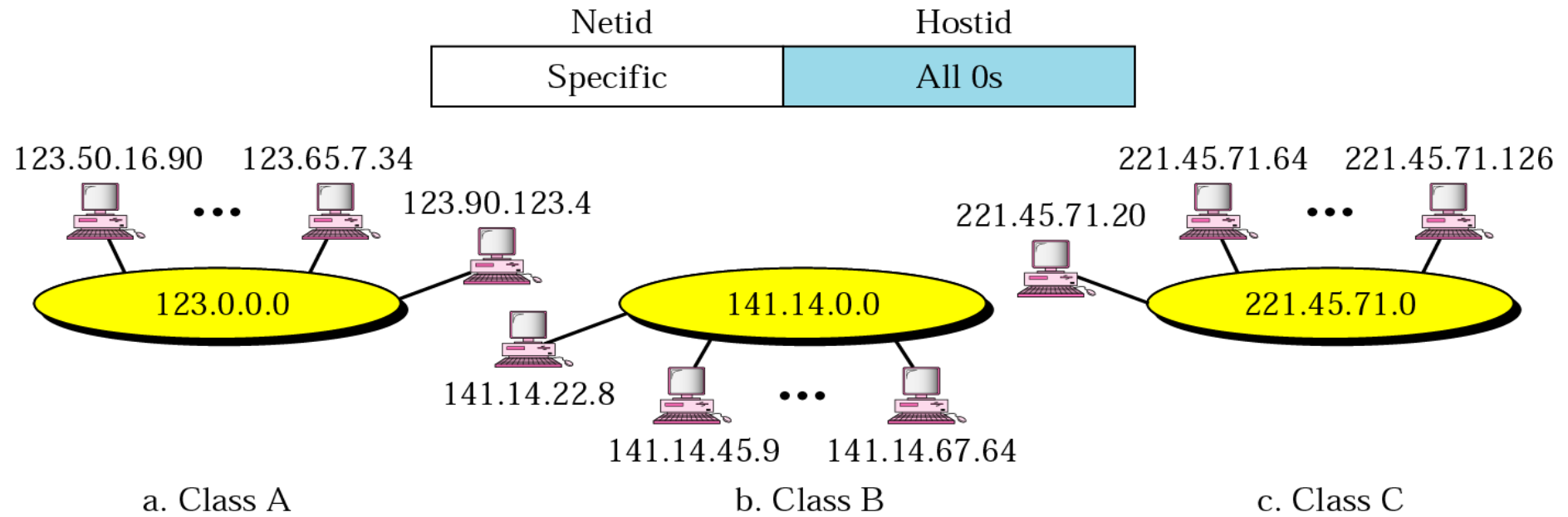


Figure 19.17 Network address





In classful addressing, the network address is the one that is assigned to the organization.

Example 5

Given the address 23.56.7.91, find the network address.

Solution

The class is A. Only the first byte defines the netid. We can find the network address by replacing the hostid bytes (56.7.91) with 0s. Therefore, the network address is 23.0.0.0.

Example 6

Given the address 132.6.17.85, find the network address.

Solution

The class is B. The first 2 bytes defines the netid. We can find the network address by replacing the hostid bytes (17.85) with 0s. Therefore, the network address is 132.6.0.0.

Example 7

Given the network address 17.0.0.0, find the class.

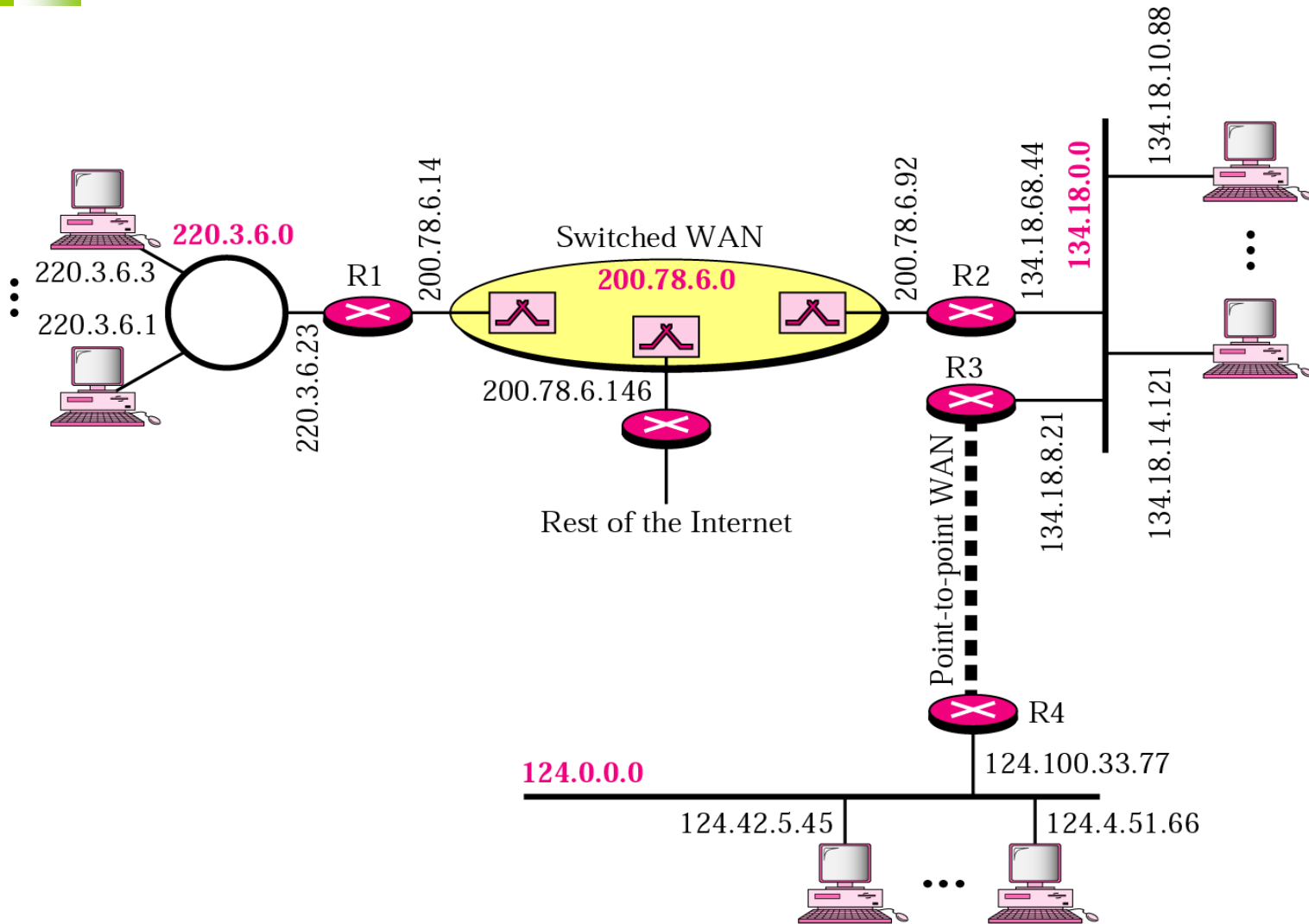
Solution

The class is A because the netid is only 1 byte.



A network address is different from a netid. A network address has both netid and hostid, with 0s for the hostid.

Figure 19.18 Sample internet





IP addresses are designed with two levels of hierarchy.

Figure 19.19 A network with two levels of hierarchy

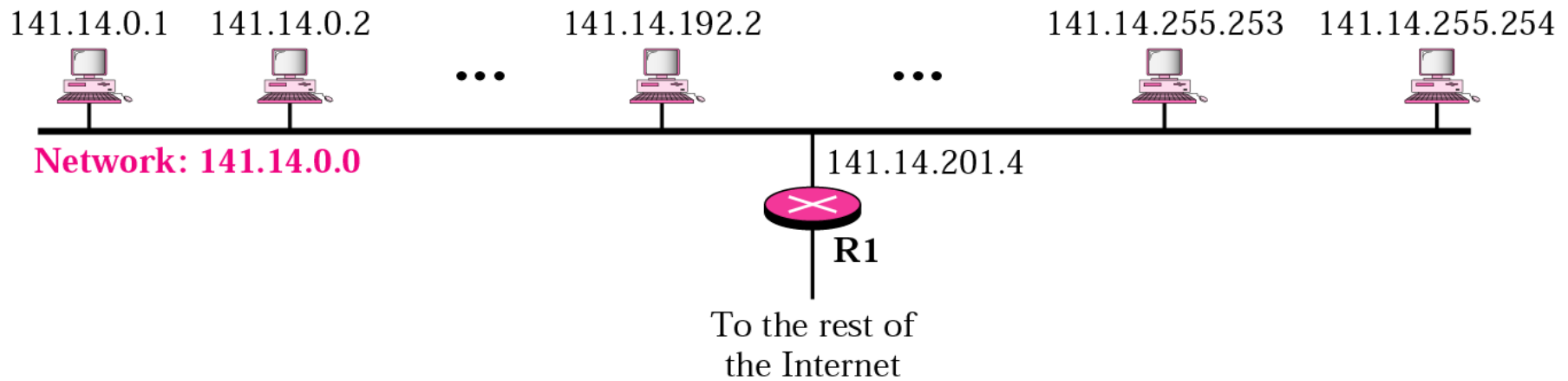


Figure 19.20 A network with three levels of hierarchy (subnetted)

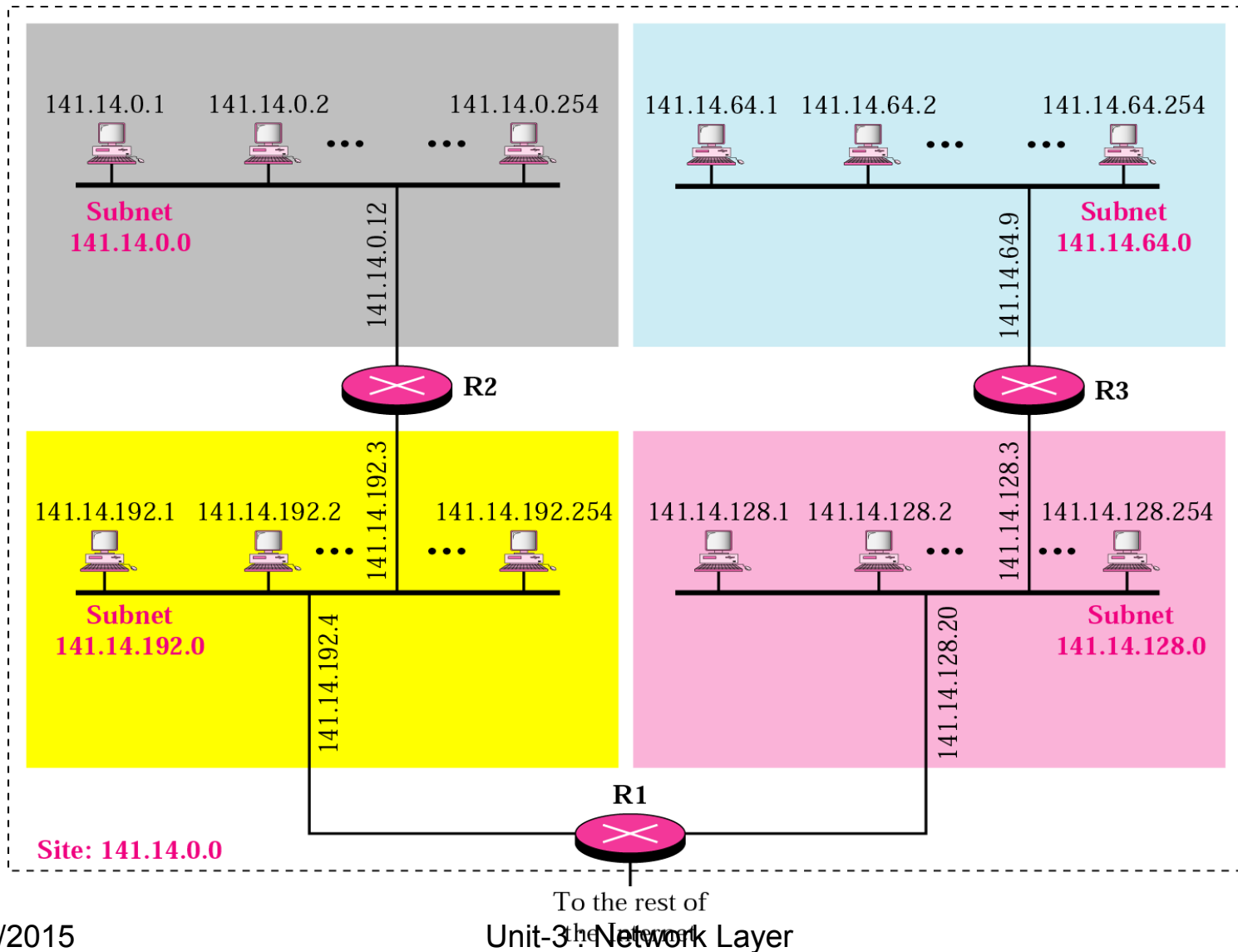
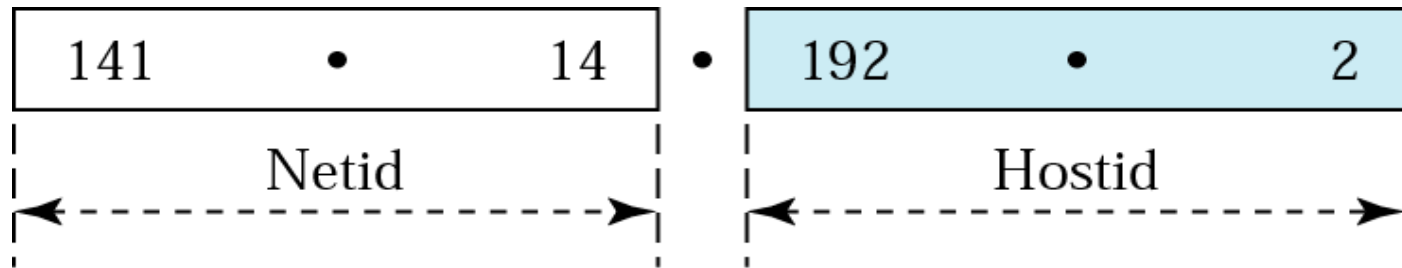
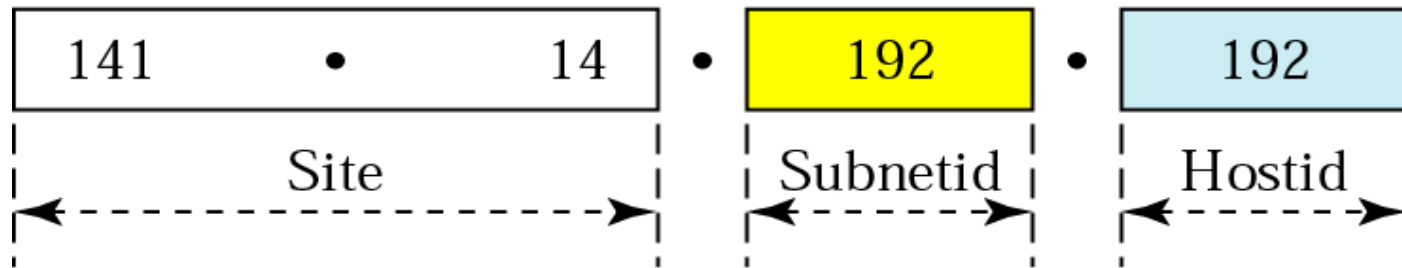


Figure 19.21 Addresses in a network with and without subnetting



a. Without subnetting



b. With subnetting

Figure 19.22 Hierarchy concept in a telephone number



Table 19.1 Default masks

Class	<i>In Binary</i>	<i>In Dotted-Decimal</i>	<i>Using Slash</i>
A	11111111 00000000 00000000 00000000	255.0.0.0	/8
B	11111111 11111111 00000000 00000000	255.255.0.0	/16
C	11111111 11111111 11111111 00000000	255.255.255.0	/24



The network address can be found by applying the default mask to any address in the block (including itself). It retains the netid of the block and sets the hostid to 0s.

Example 8

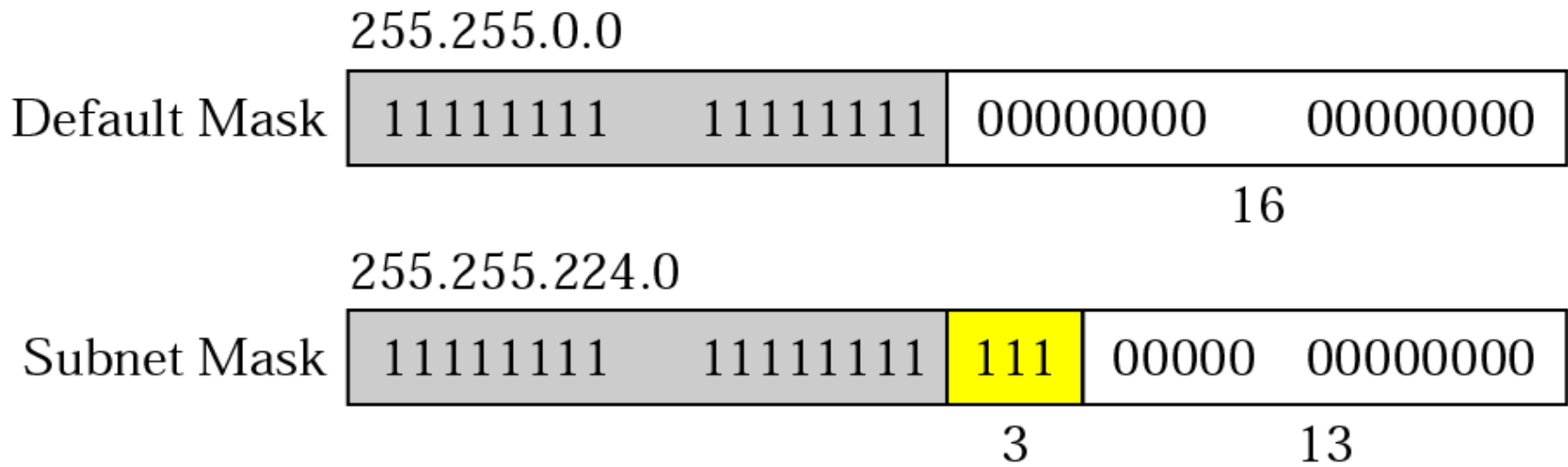
A router outside the organization receives a packet with destination address 190.240.7.91. Show how it finds the network address to route the packet.

Solution

The router follows three steps:

1. The router looks at the first byte of the address to find the class. It is class B.
2. The default mask for class B is 255.255.0.0. The router ANDs this mask with the address to get 190.240.0.0.
3. The router looks in its routing table to find out how to route the packet to this destination. Later, we will see what happens if this destination does not exist.

Figure 19.23 Subnet mask



Example 9

A router inside the organization receives the same packet with destination address 190.240.33.91. Show how it finds the subnetwork address to route the packet.

Solution

The router follows three steps:

- 1. The router must know the mask. We assume it is /19, as shown in Figure 19.23.**
- 2. The router applies the mask to the address, 190.240.33.91. The subnet address is 190.240.32.0.**
- 3. The router looks in its routing table to find how to route the packet to this destination. Later, we will see what happens if this destination does not exist.**

Figure 19.24 DHCP transition diagram

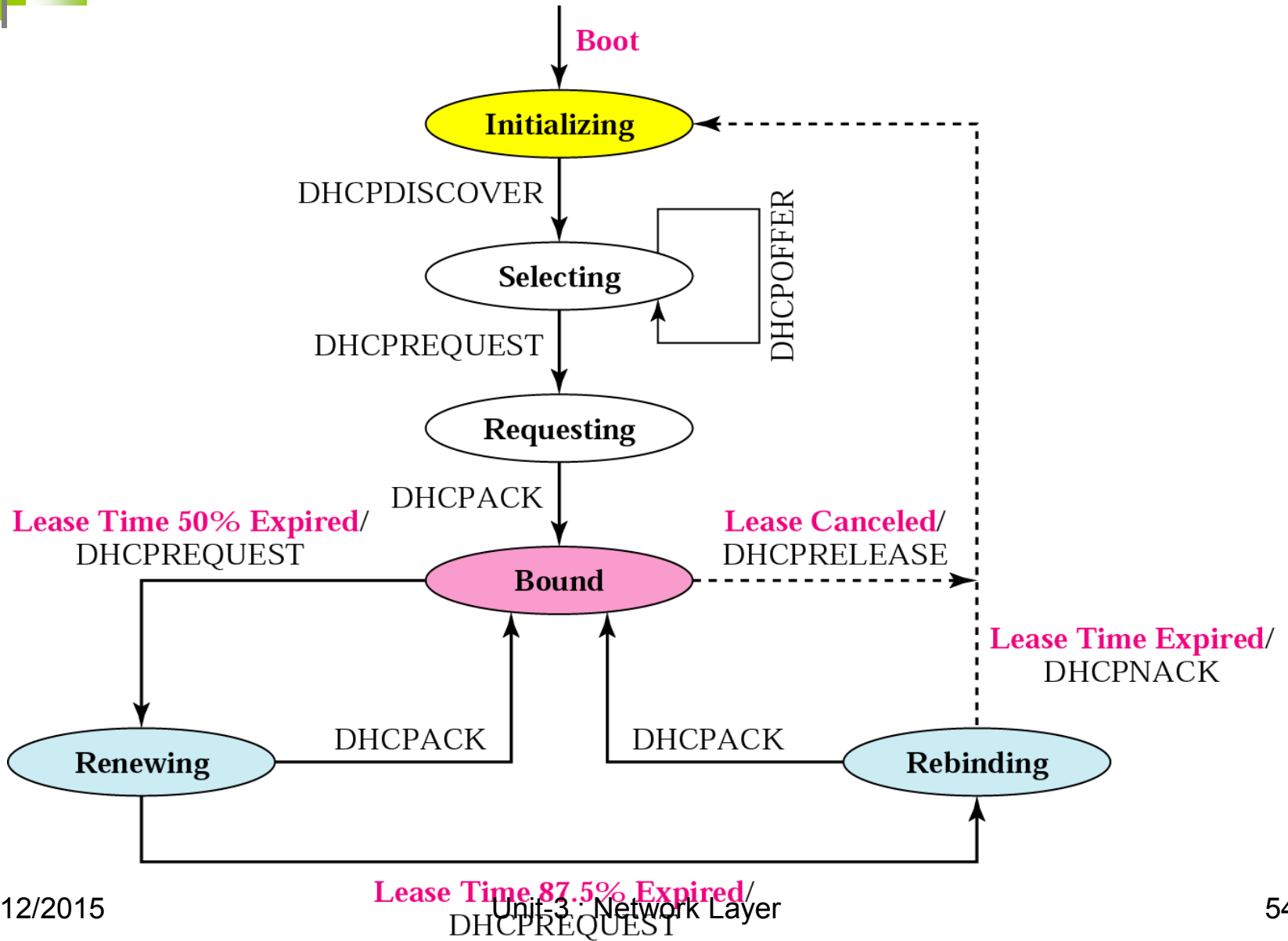


Table 19.2 Default masks

<i>Range</i>		<i>Total</i>
10.0.0.0	to 10.255.255.255	2^{24}
172.16.0.0	to 172.31.255.255	2^{20}
192.168.0.0	to 192.168.255.255	2^{16}

Figure 19.25 NAT

Site using private addresses

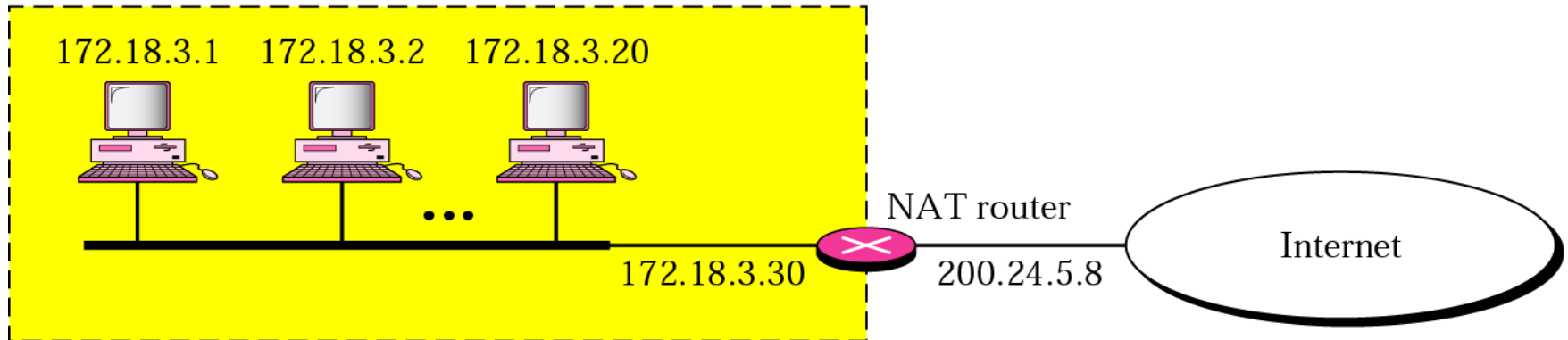


Figure 19.26 Address translation

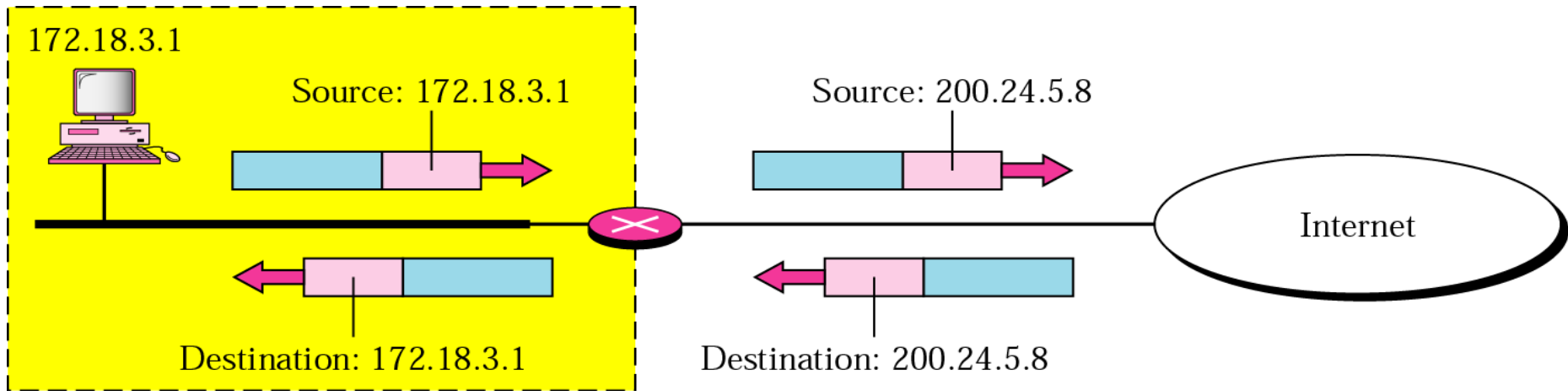


Figure 19.27 Translation

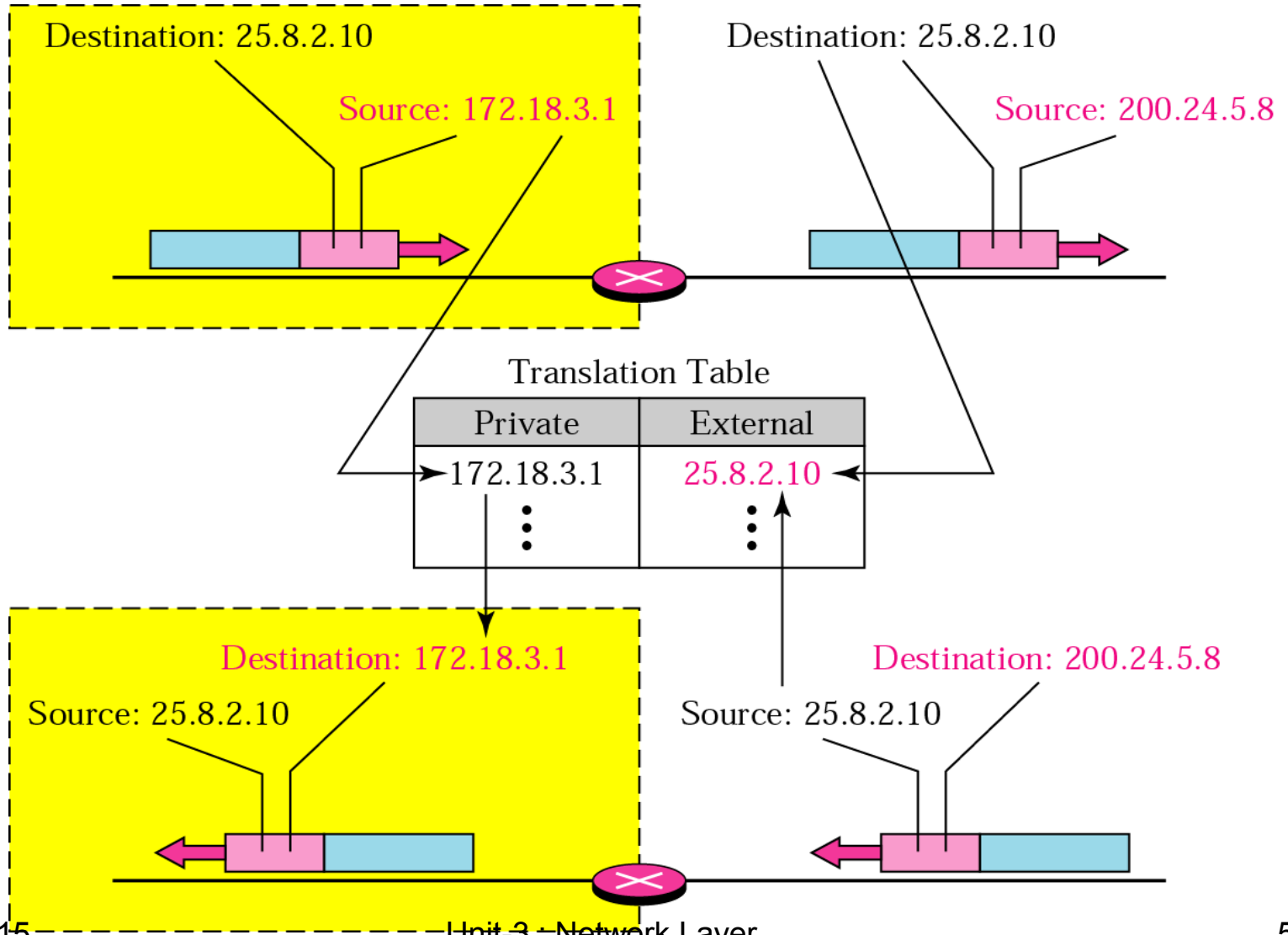


Table 19.3 Five-column translation table

<i>Private Address</i>	<i>Private Port</i>	<i>External Address</i>	<i>External Port</i>	<i>Transport Protocol</i>
172.18.3.1	1400	25.8.3.2	80	TCP
172.18.3.2	1401	25.8.3.2	80	TCP
...

19.3 Routing

Routing Techniques

Static Versus Dynamic Routing

Routing Table for Classful Addressing

Routing Table for Classless Addressing

Figure 19.28 Next-hop routing

Routing table for host A

Destination	Route
Host B	R1, R2, Host B

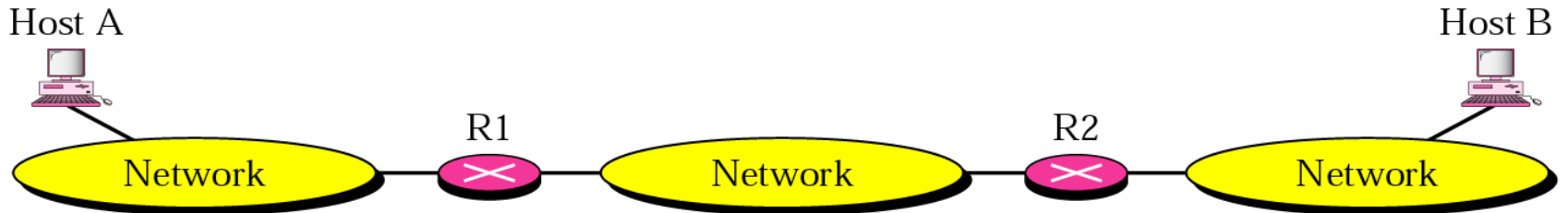
Routing table for R1

Destination	Route
Host B	R2, Host B

Routing table for R2

Destination	Route
Host B	Host B

a. Routing tables based on route



Routing table for host A

Destination	Next Hop
Host B	R1

Routing table for R1

Destination	Next Hop
Host B	R2

Routing table for R2

Destination	Next Hop
Host B	—

b. Routing tables based on next hop

Figure 19.29 Network-specific routing

Routing table for host S based on host-specific routing

Destination	Next Hop
A	R1
B	R1
C	R1
D	R1

Routing table for host S based on network-specific routing

Destination	Next Hop
N2	R1

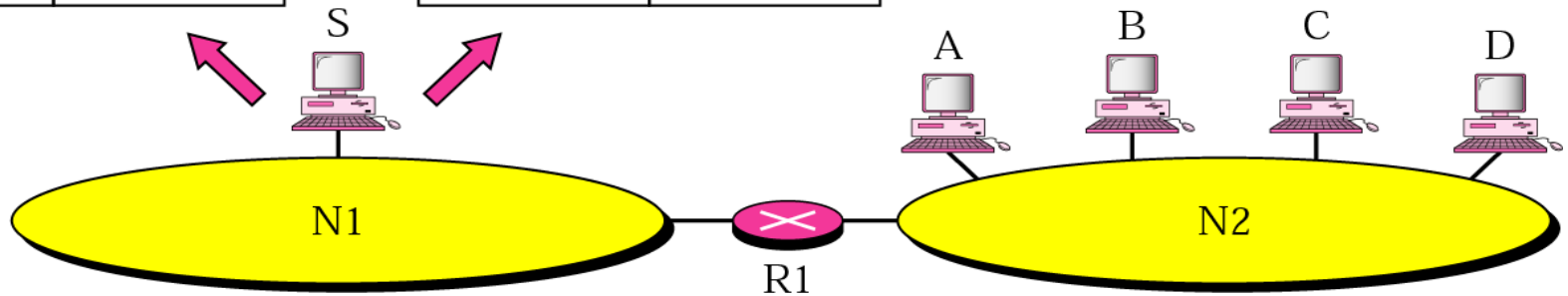


Figure 19.30 Host-specific routing

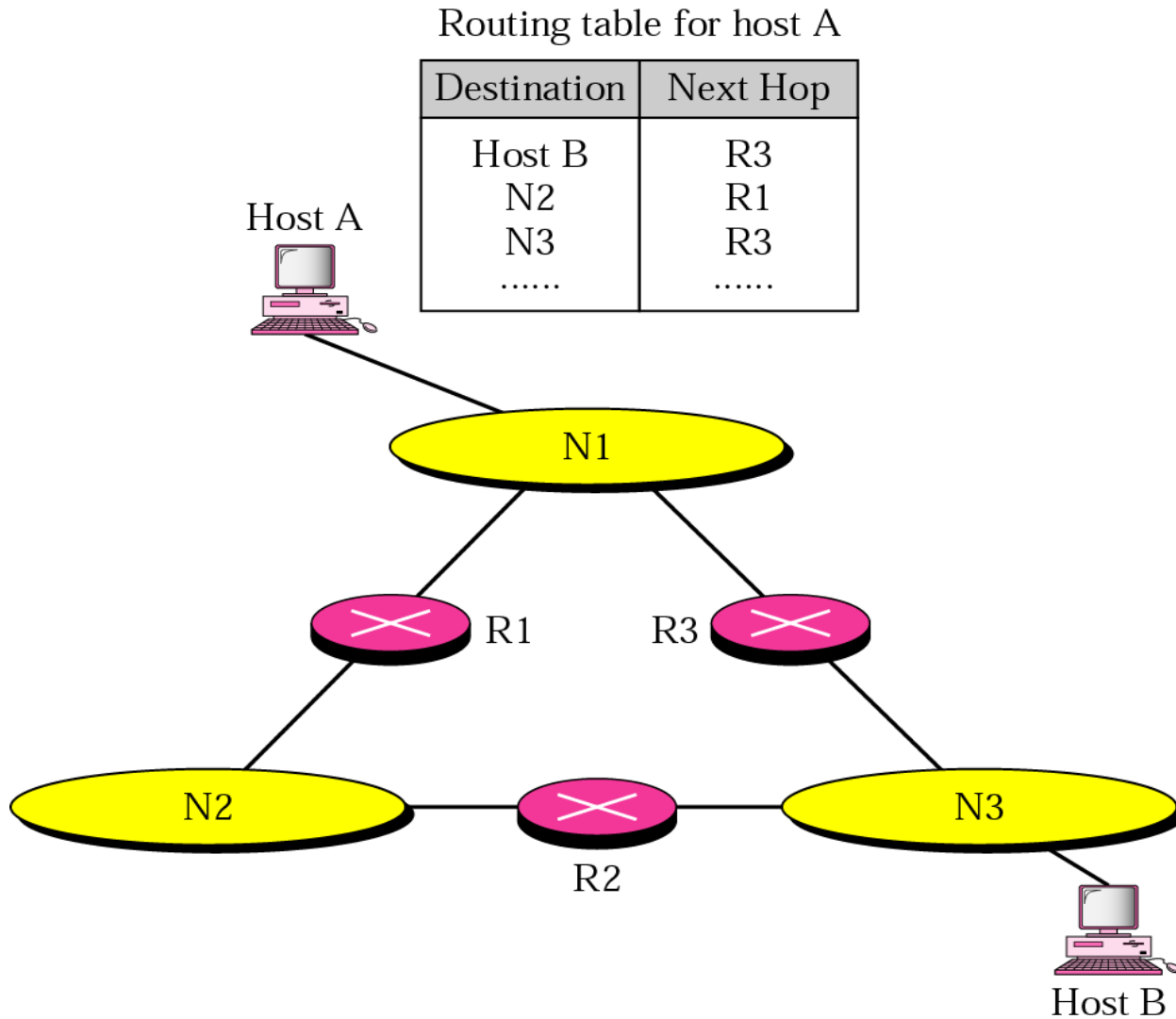


Figure 19.31 Default routing

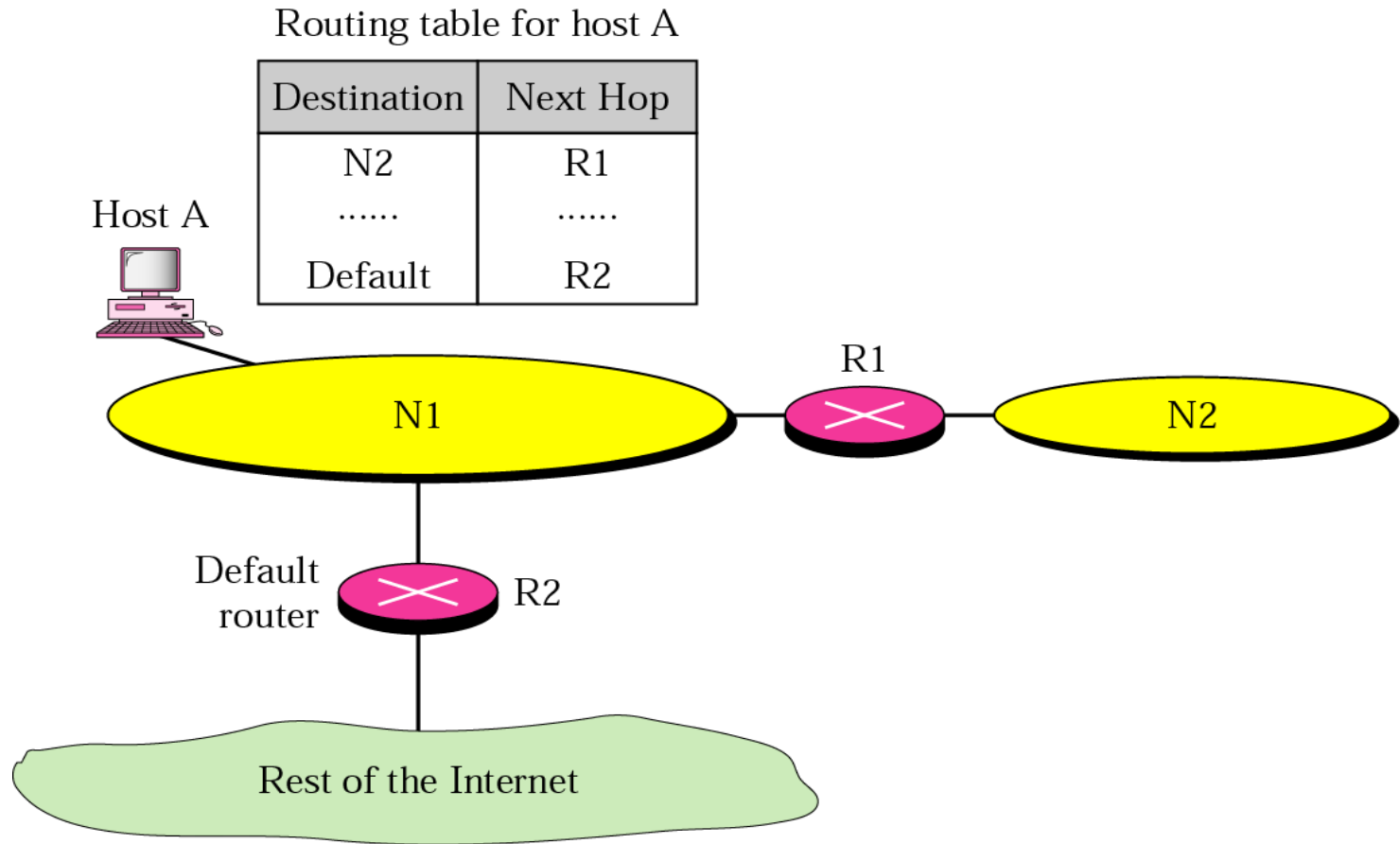


Figure 19.32 Classful addressing routing table

	Mask	Destination address	Next-hop address	Interface
	/8	14.0.0.0	118.45.23.8	m1
Host-specific →	/32	192.16.7.1	202.45.9.3	m0
	/24	193.14.5.0	84.78.4.12	m2
Default →	/0	/0	145.11.10.6	m0

Example 10

Using the table in Figure 19.32, the router receives a packet for destination 192.16.7.1. For each row, the mask is applied to the destination address until a match with the destination address is found. In this example, the router sends the packet through interface m0 (host specific).

Example 11

Using the table in Figure 19.32, the router receives a packet for destination 193.14.5.22. For each row, the mask is applied to the destination address until a match with the next-hop address is found. In this example, the router sends the packet through interface m2 (network specific).

Example 12

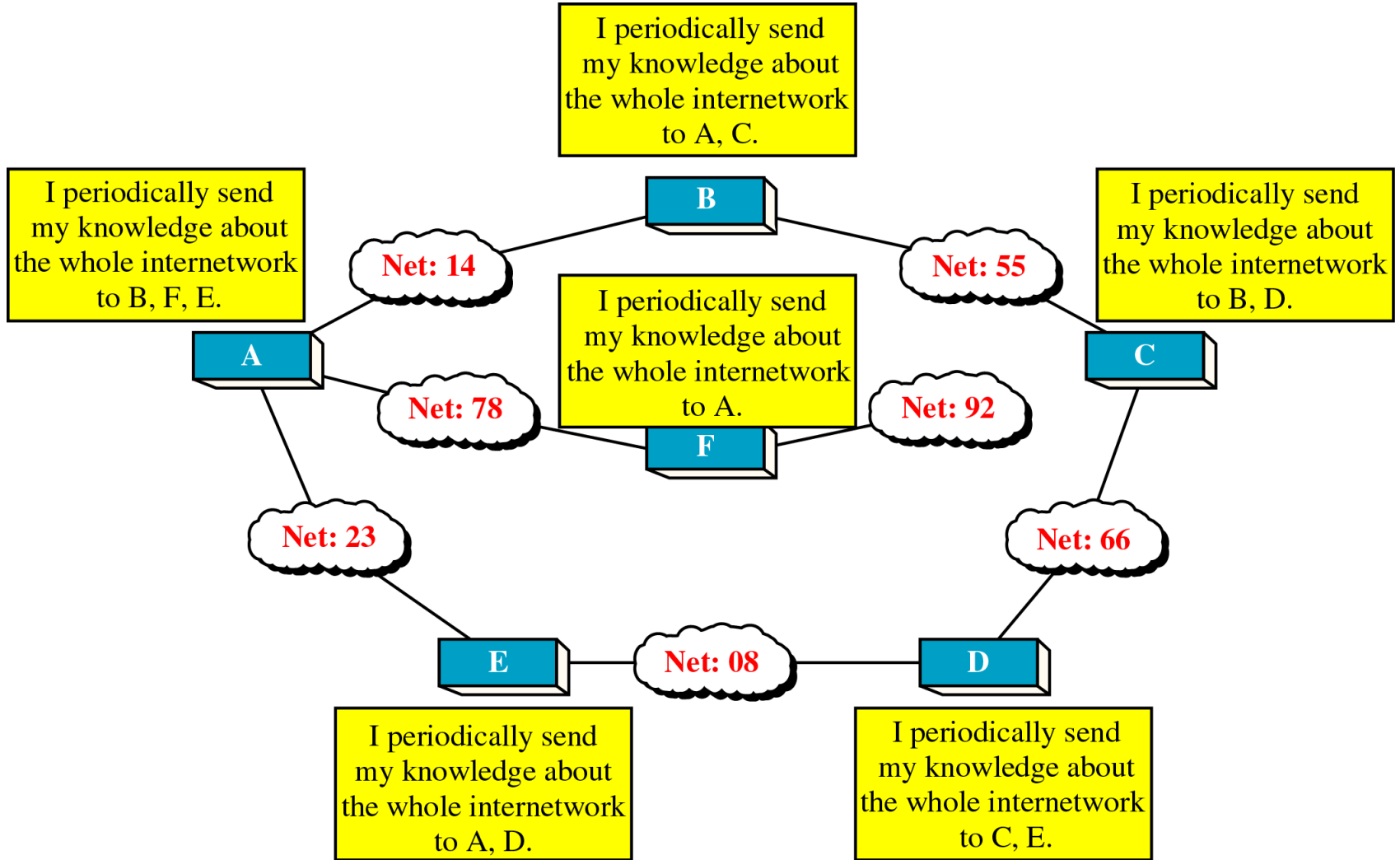
Using the table in Figure 19.32, the router receives a packet for destination 200.34.12.34. For each row, the mask is applied to the destination address, but no match is found. In this example, the router sends the packet through the default interface m0.

Routing Algorithms

- 1.Distance Vector Routing
- 2.Link State Routing

Figure 21-18

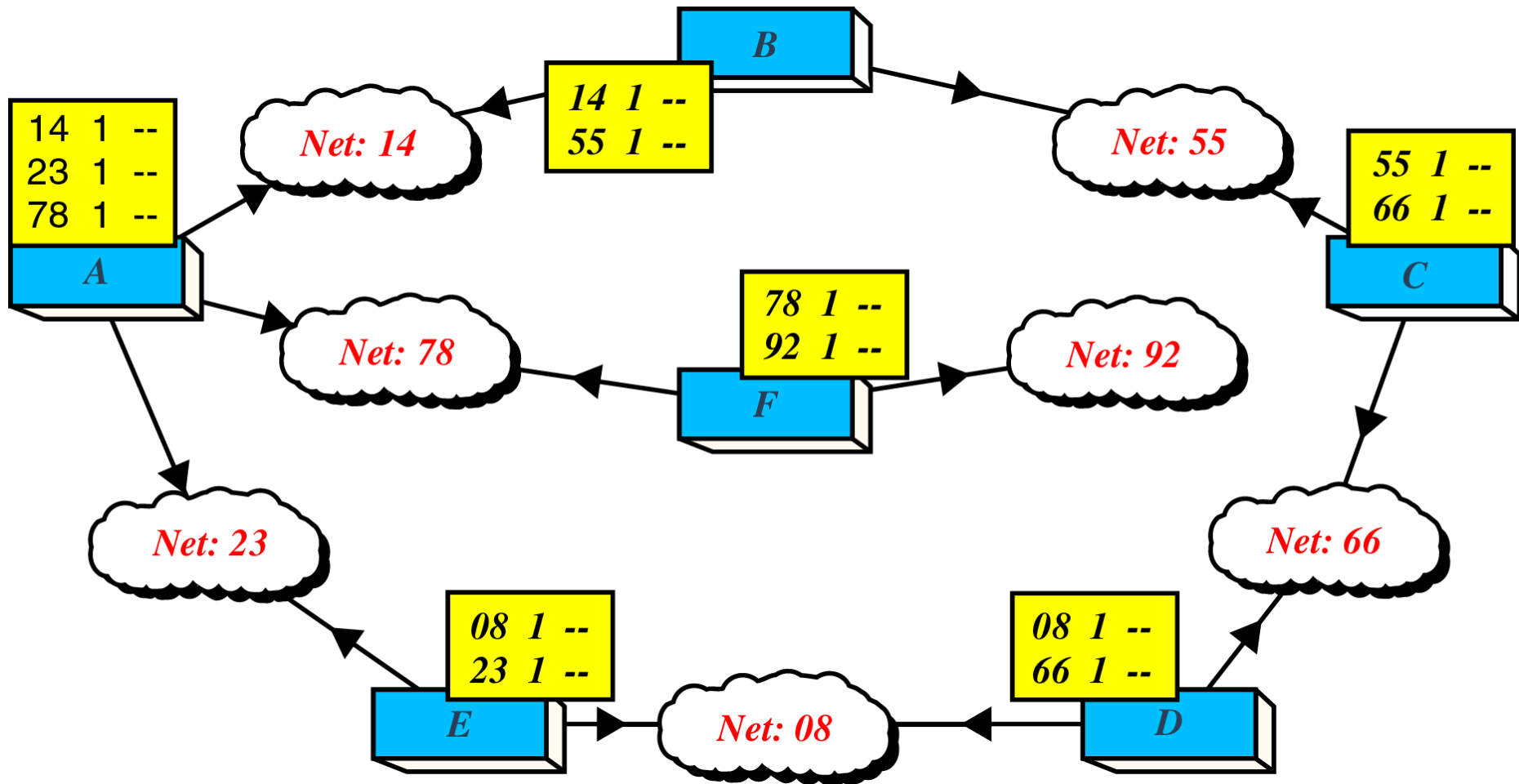
The Concept of Distance Vector Routing



Distance Vector Routing Table

Network ID	Cost	Next Hop
• • • • • • • • • •	• • • • • • • •	• • • • • • • • • •
• • • • • • • • • •	• • • • • • • •	• • • • • • • • • •
• • • • • • • • • •	• • • • • • • •	• • • • • • • • • •
• • • • • • • • • •	• • • • • • • •	• • • • • • • • • •
• • • • • • • • • •	• • • • • • • •	• • • • • • • • • •

Routing Table Distribution



Updating Routing Table for Router A

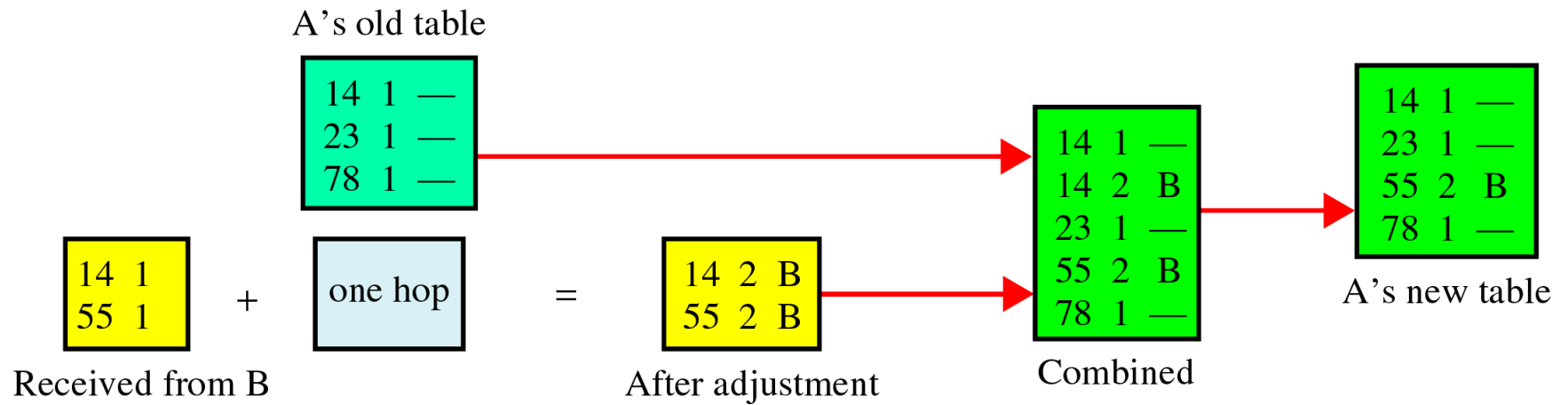
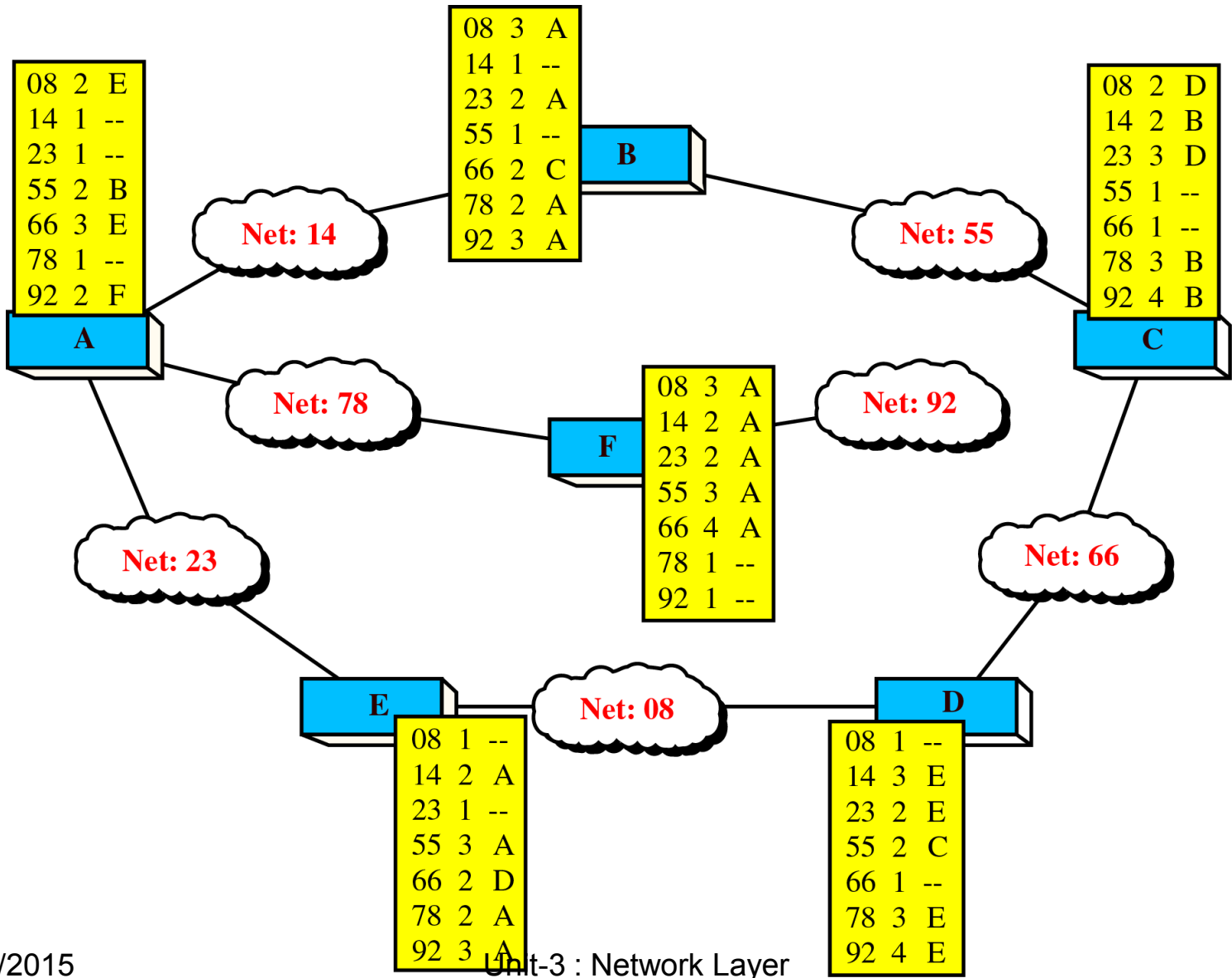
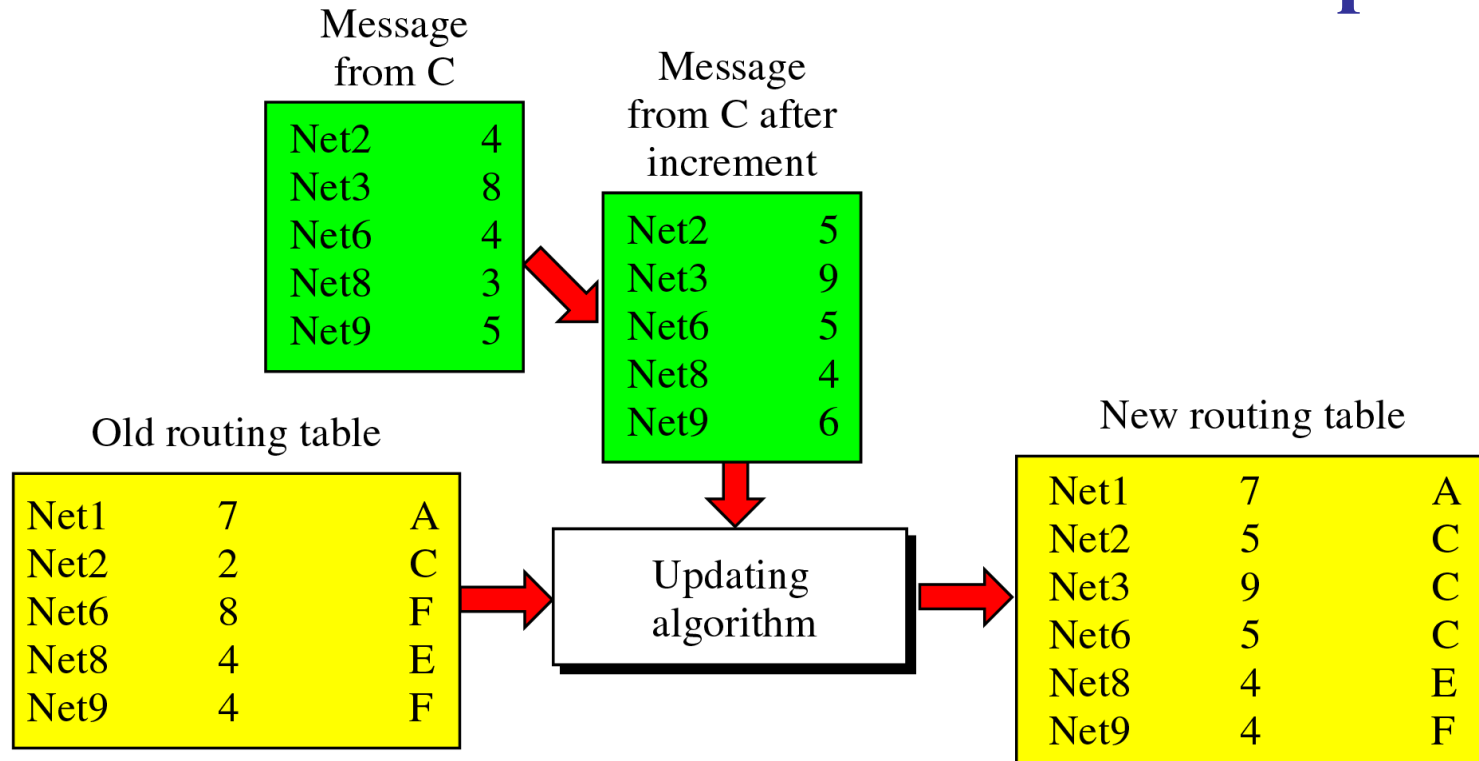


Figure 21-22

Final Routing Tables





Rules

Net2: Replace (**Rule 2.a**)

Net3: Add (**Rule 1**)

Net6: Replace (**Rule 2.b.i**)

Net8: No change (**Rule 2.b.ii**)

Net9: No change (**Rule 2.b.ii**)

Note that there is no news about Net1 in the advertised message, so none of the rules apply to this entry.

Figure 21-24

Concept of Link State Routing

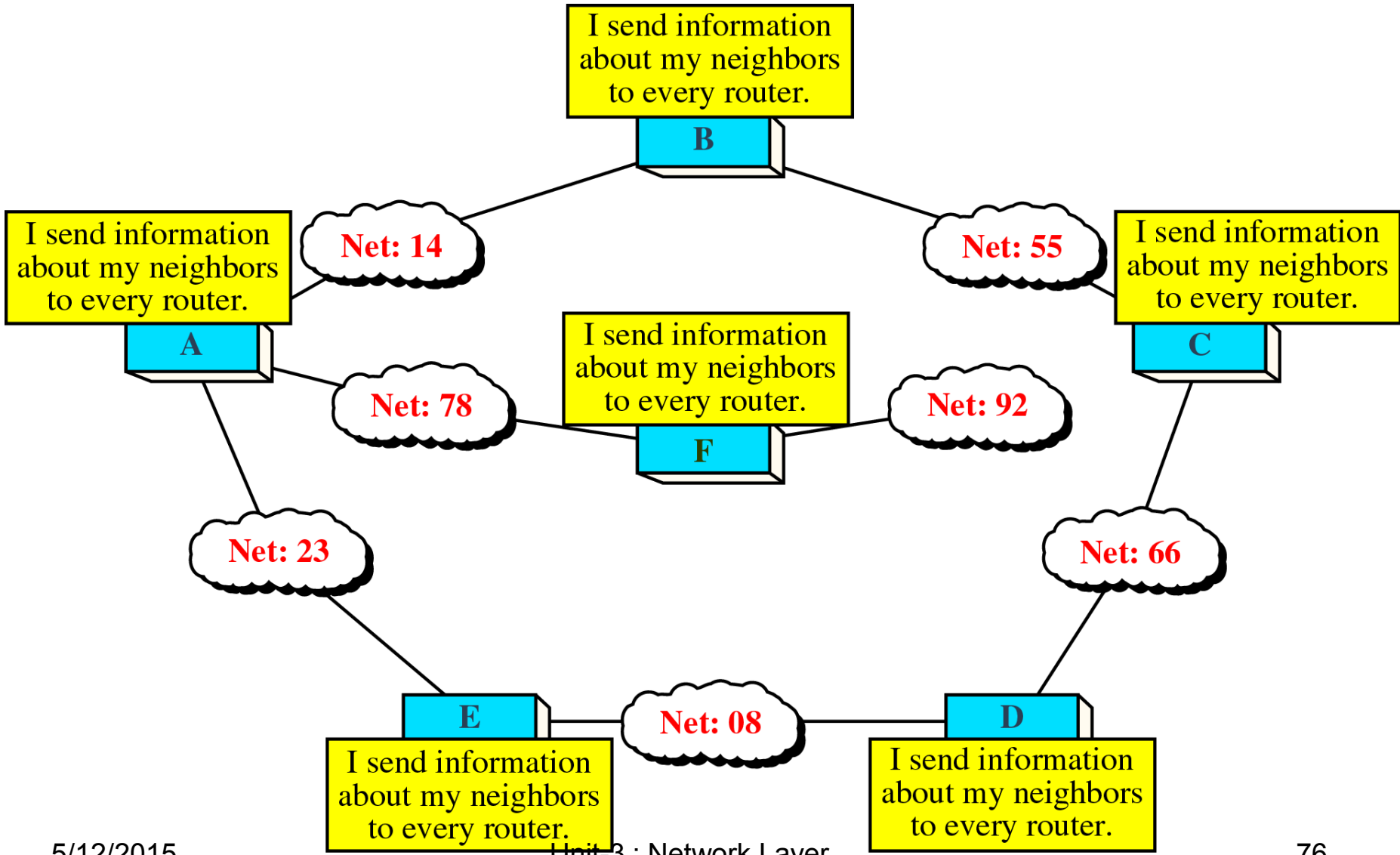
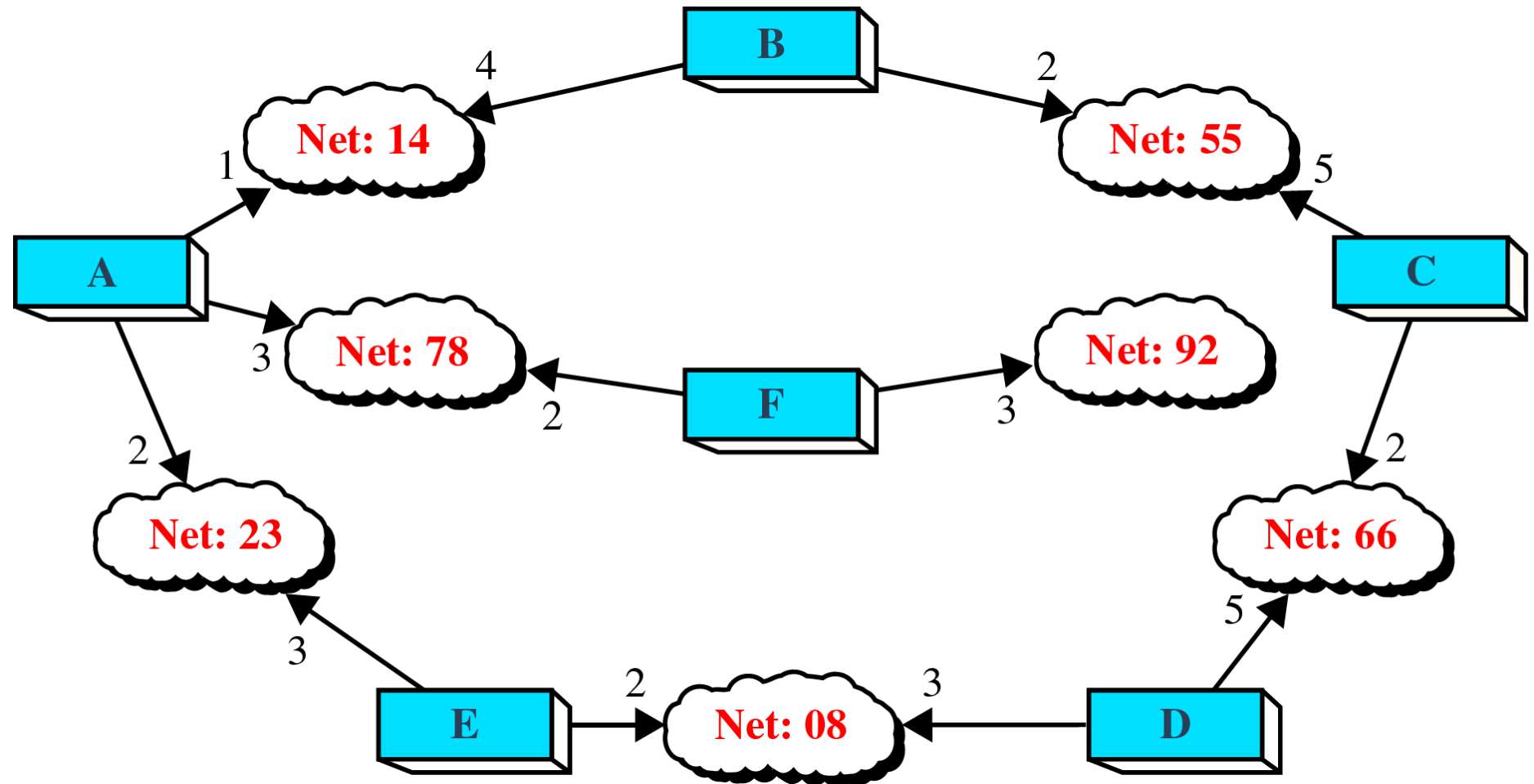


Figure 21-25

Cost in Link State Routing



Link State Packet

Advertiser	Network	Cost	Neighbor
• • • • • • •	• • • • • • •	• • • • • • • • • • • • •	• • • • • • • • • • • • •
• • • • • • •	• • • • • • •	• • • • • • • • • • • • •	• • • • • • • • • • • • •
• • • • • • •	• • • • • • •	• • • • • • • • • • • • •	• • • • • • • • • • • • •

Figure 21-27

Flooding of A's LSP

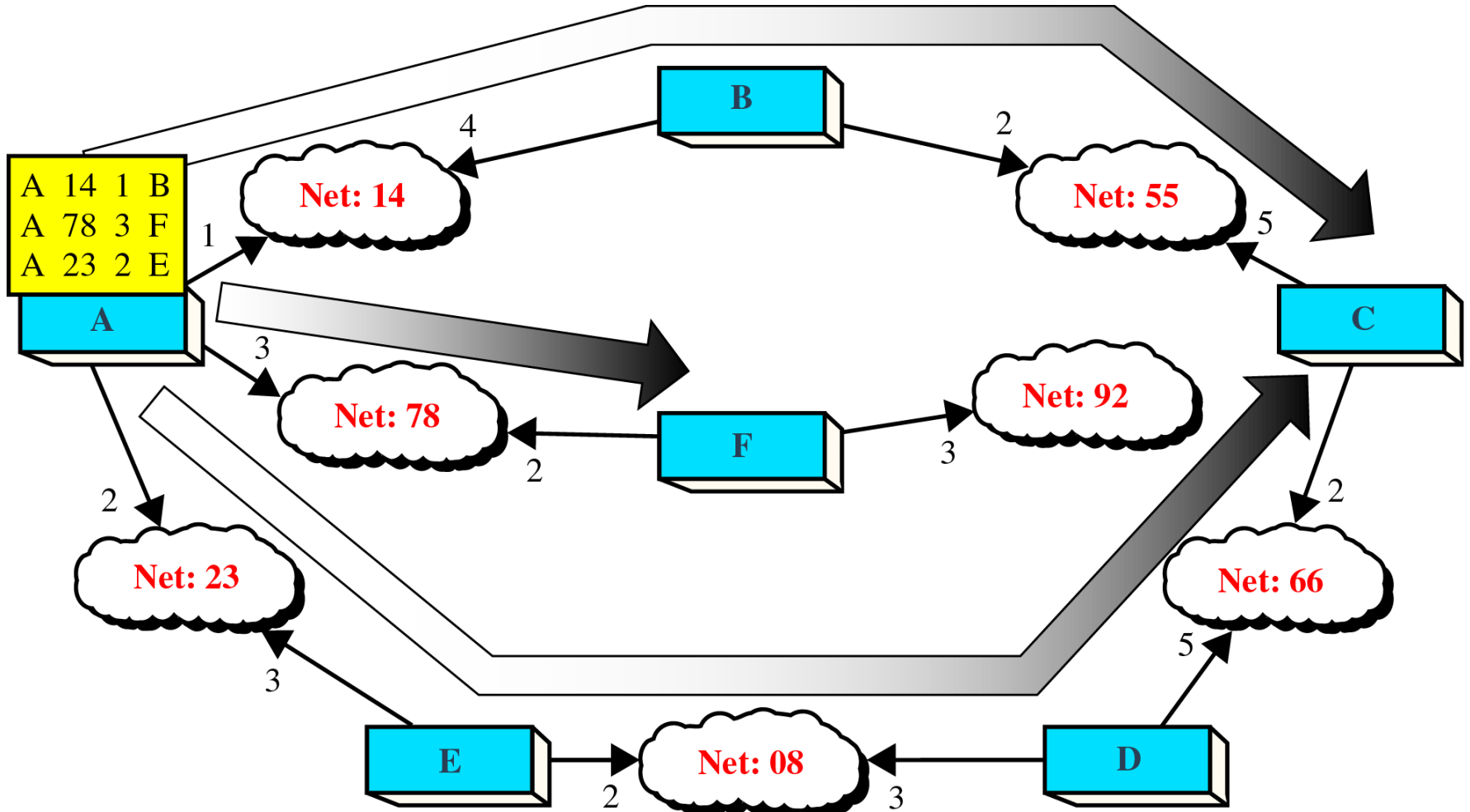


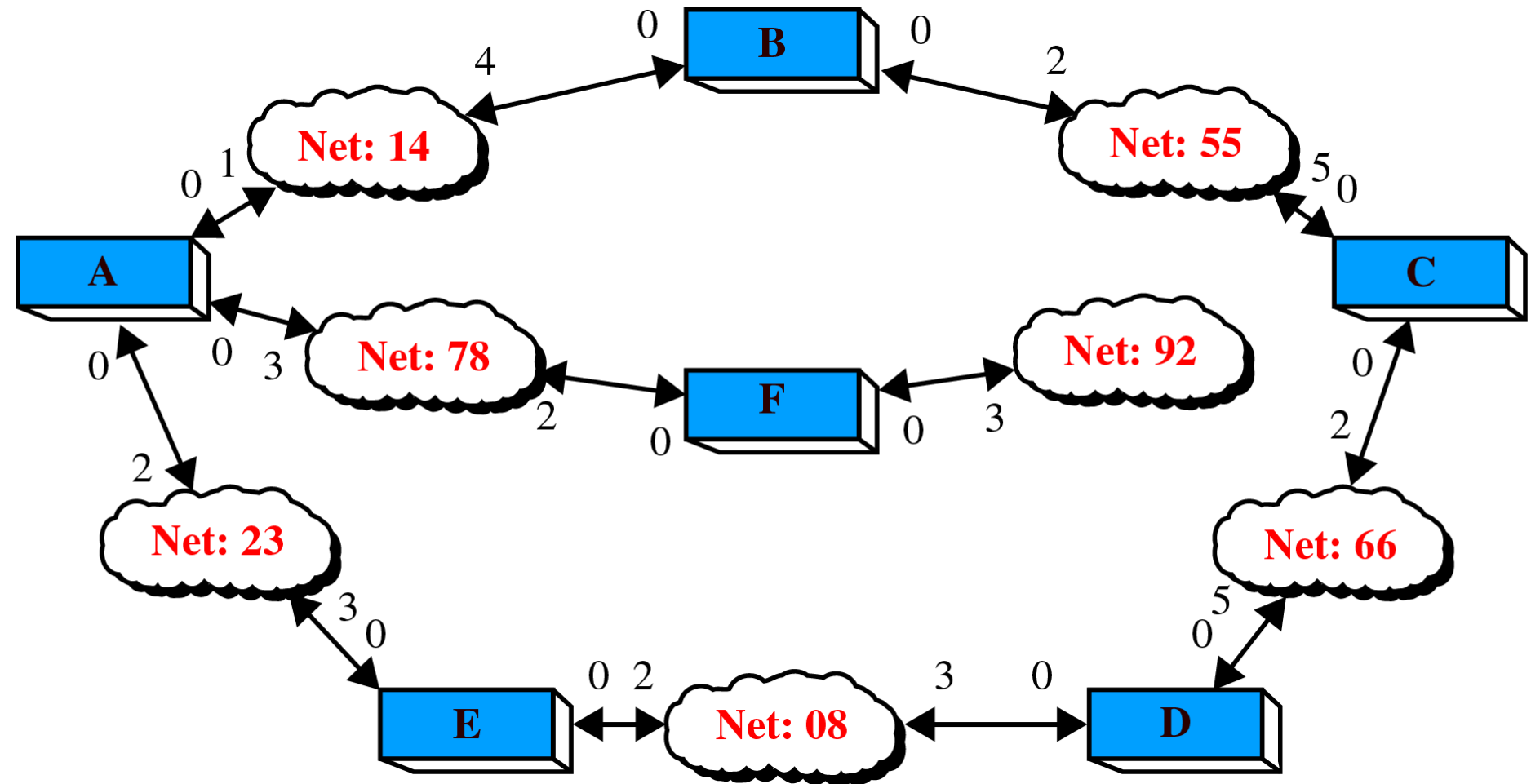
Figure 21-28

Link State Database

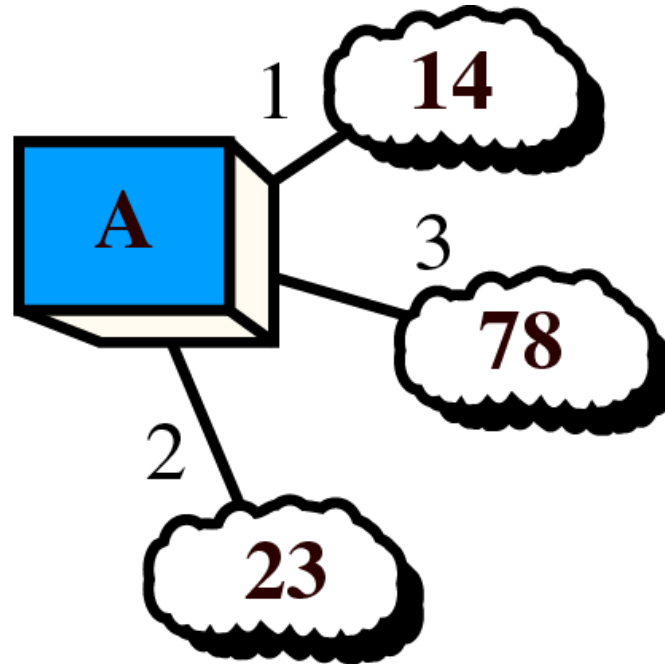
Advertiser	Network	Cost	Neighbor
A	14	1	B
A	78	3	F
A	23	2	E
B	14	4	A
B	55	2	C
C	55	5	B
C	66	2	D
D	66	5	C
D	08	3	E
E	23	3	A
E	08	2	D
F	78	2	A
F	92	3	—

Figure 21-29

Costs in the Dijkstra Algorithm

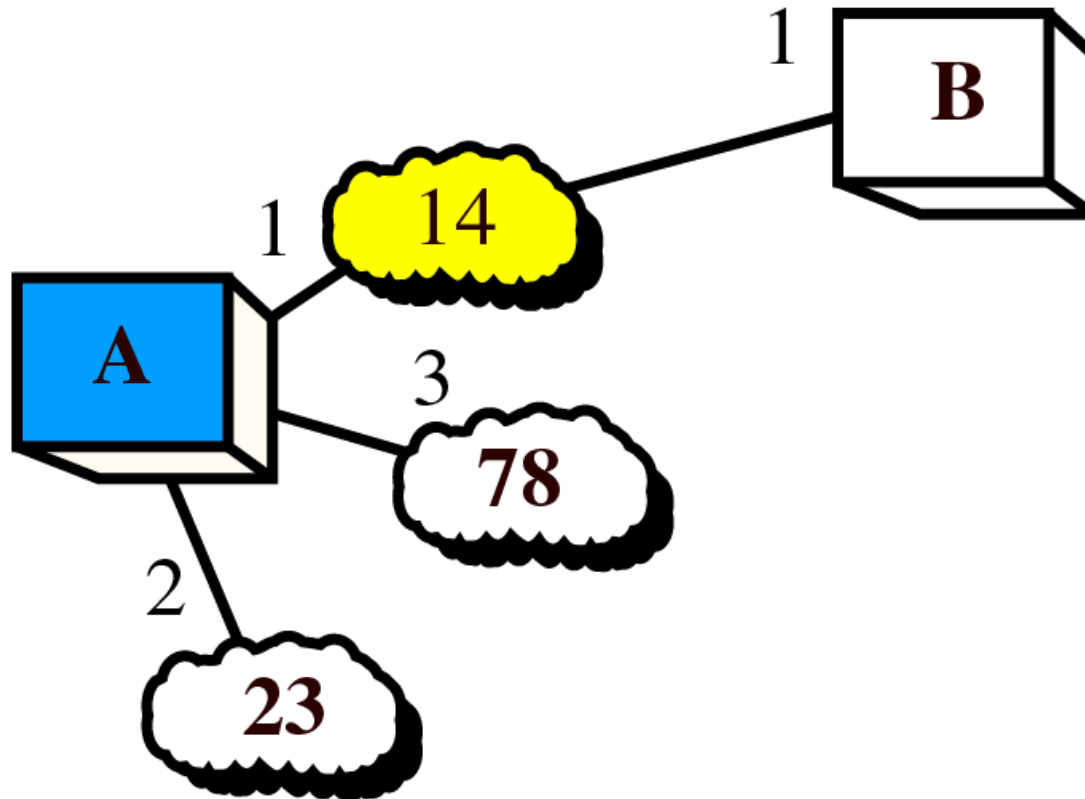


Shortest Path Calculation, Part I



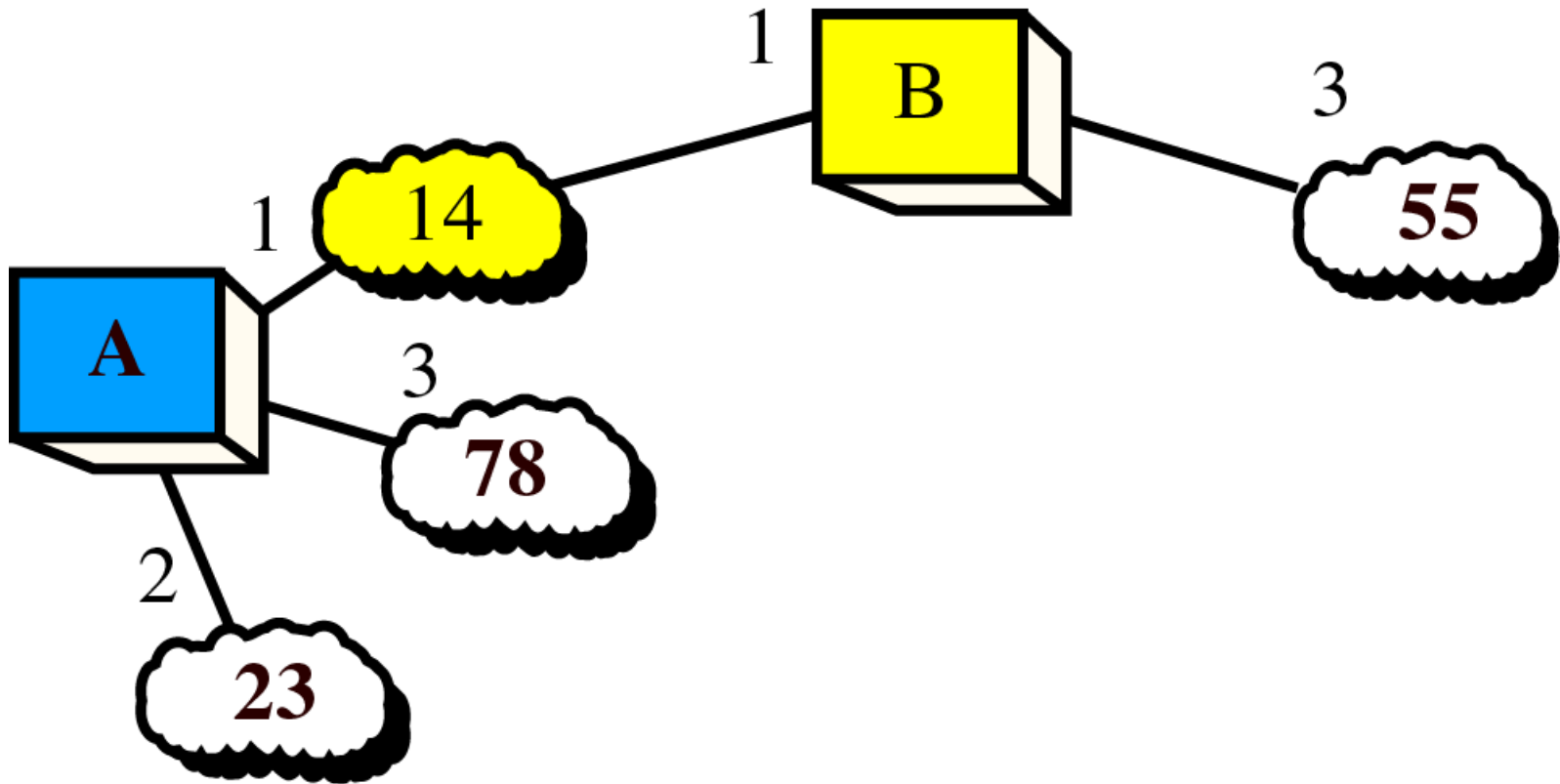
Root is A, networks
14, 78, 23 added

Shortest Path Calculation, Part II



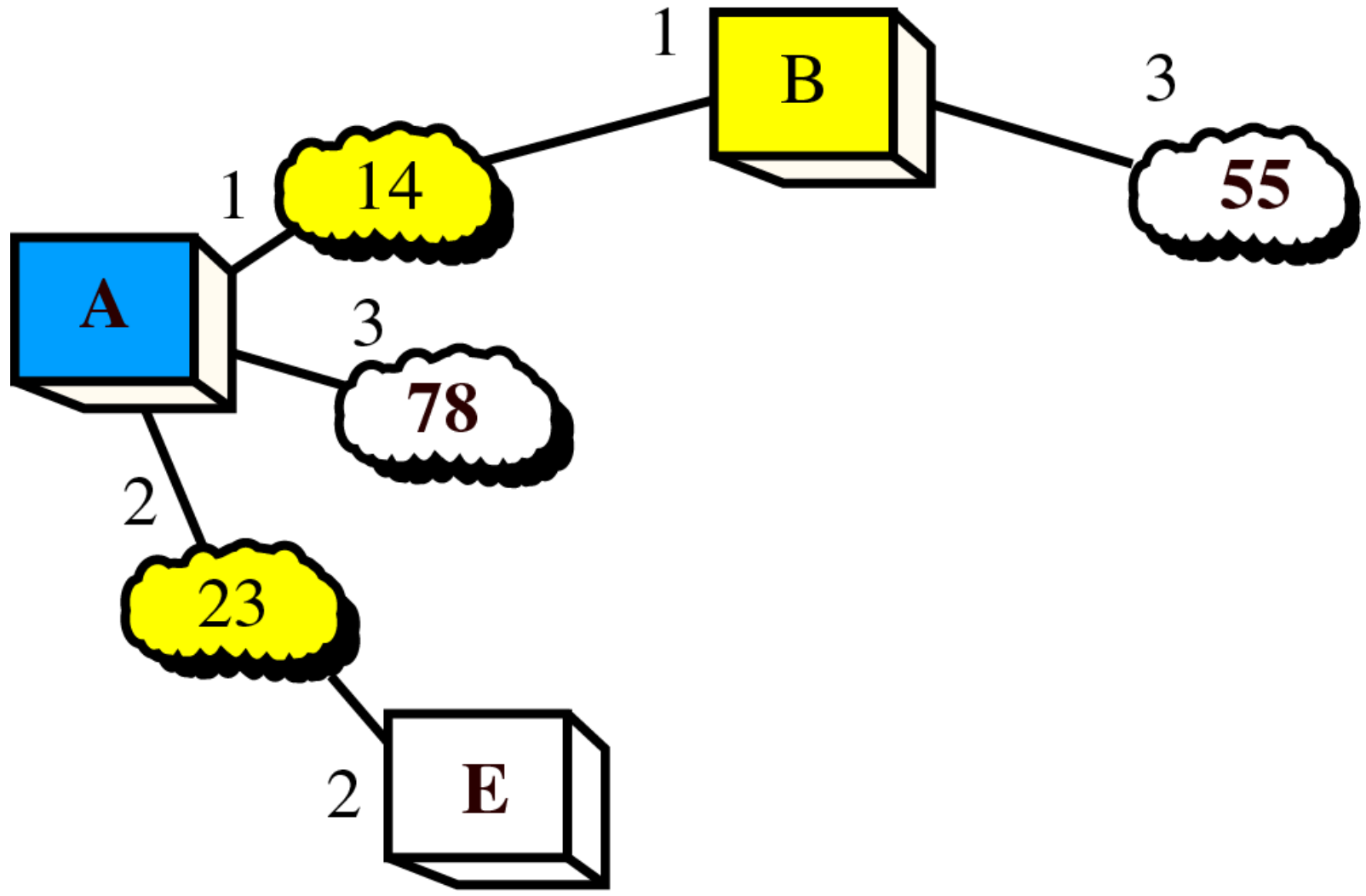
14 permanent, B added

Shortest Path Calculation, Part III



B Permanent, 55 added

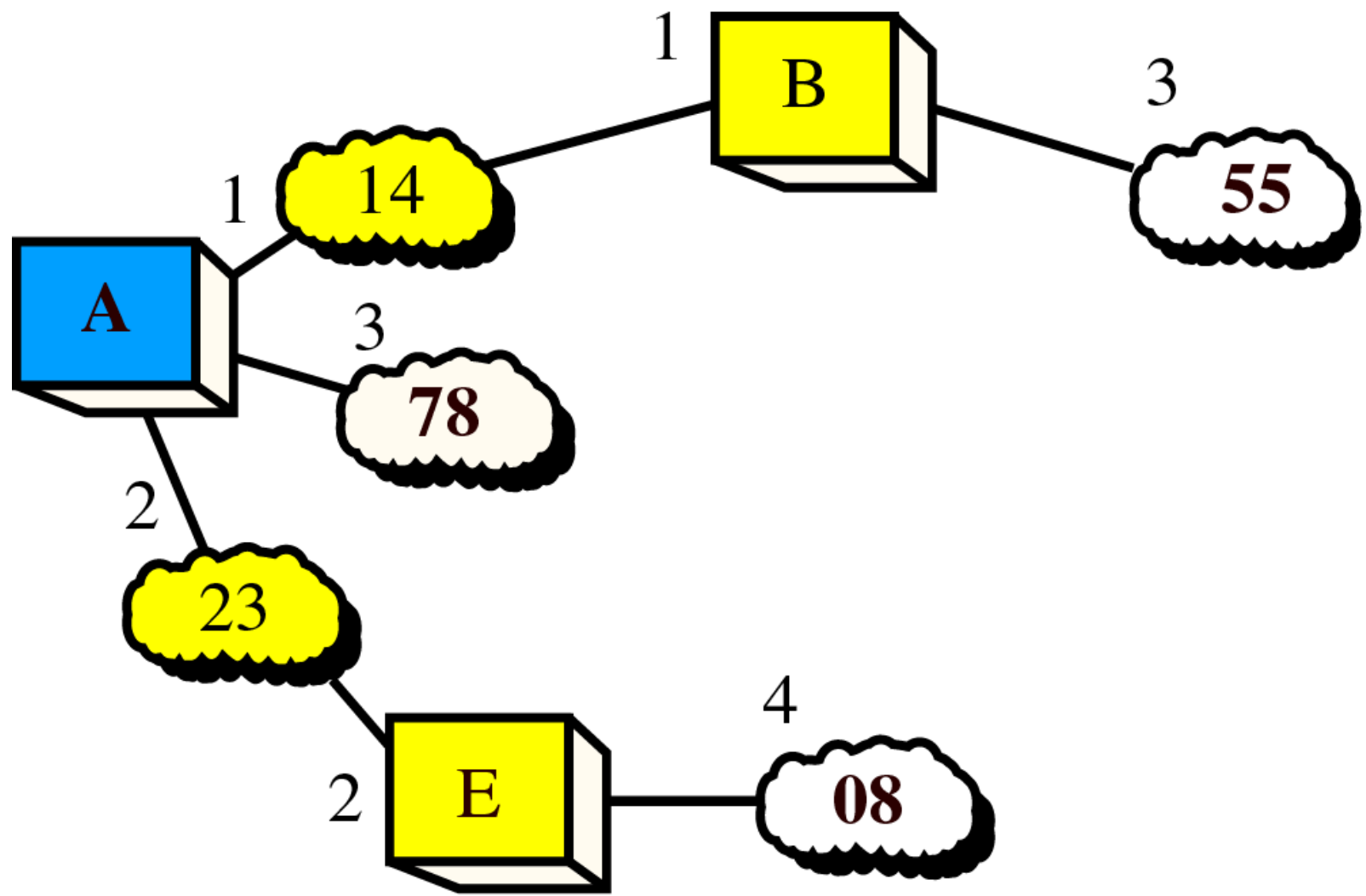
Shortest Path Calculation, Part IV



23 permanent, E added

Figure 21-30, Part V

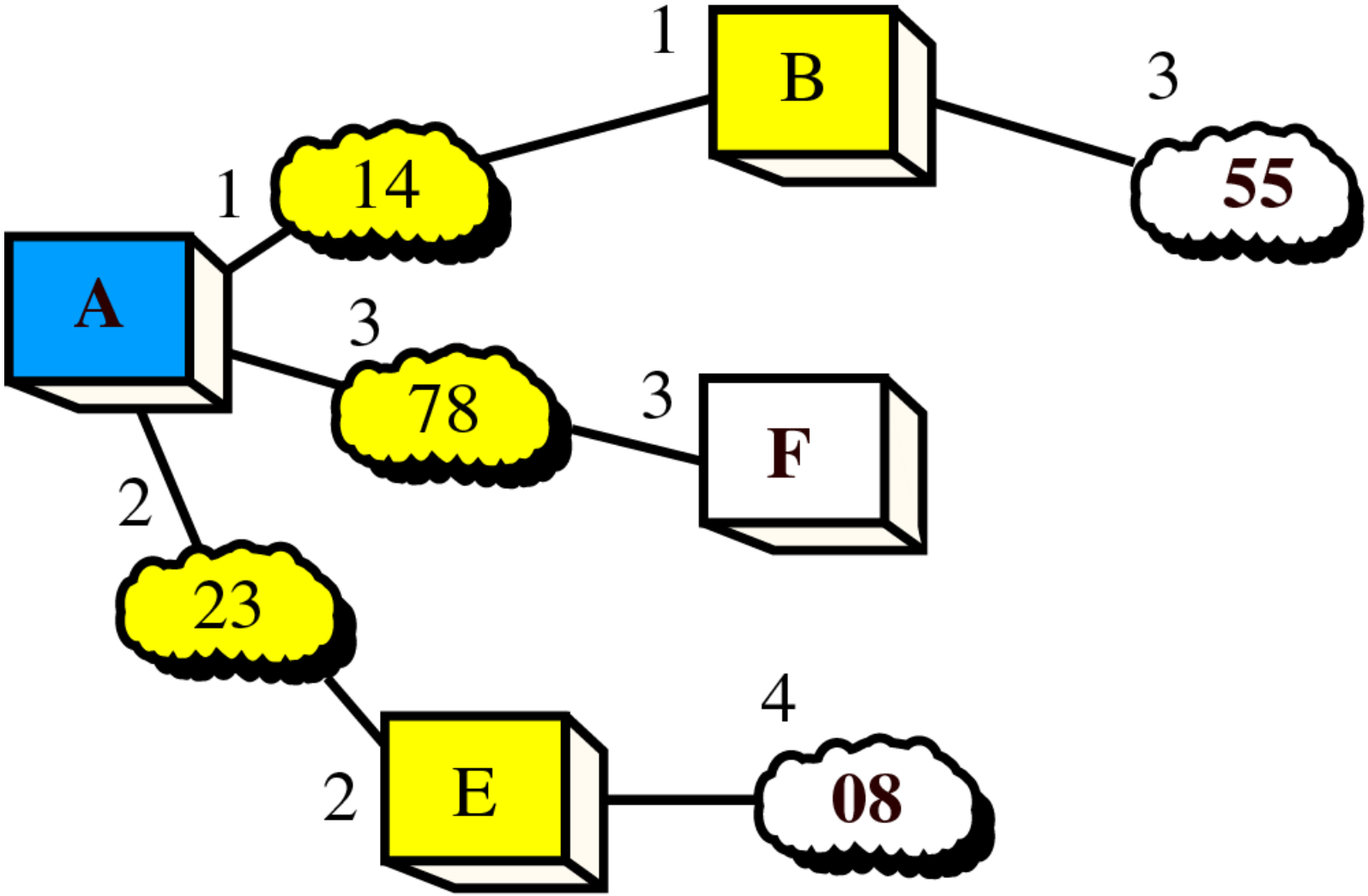
Shortest Path Calculation, Part V



E permanent, 08 added

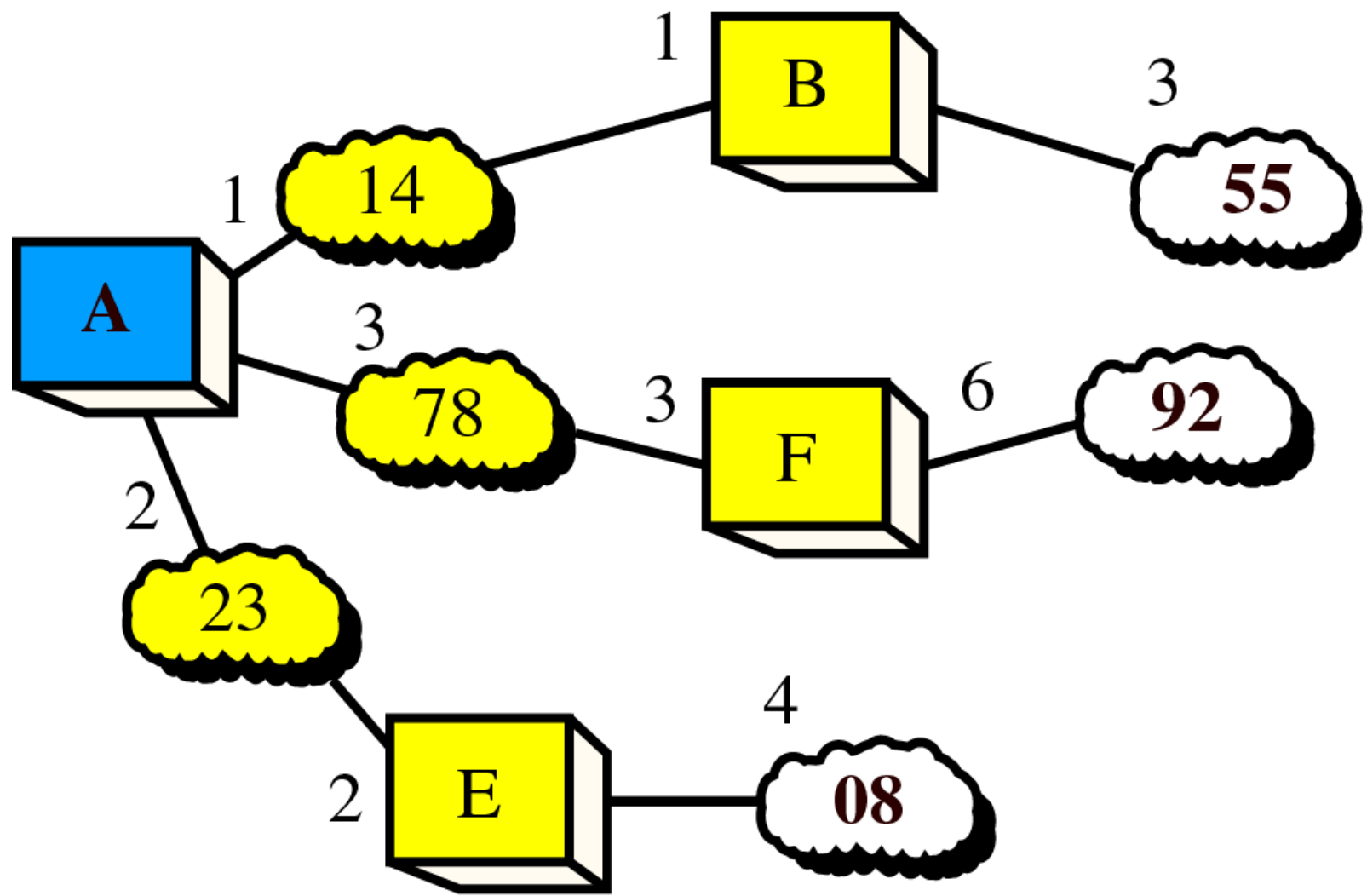
Figure 21-30, Part VI

Shortest Path Calculation, Part VI



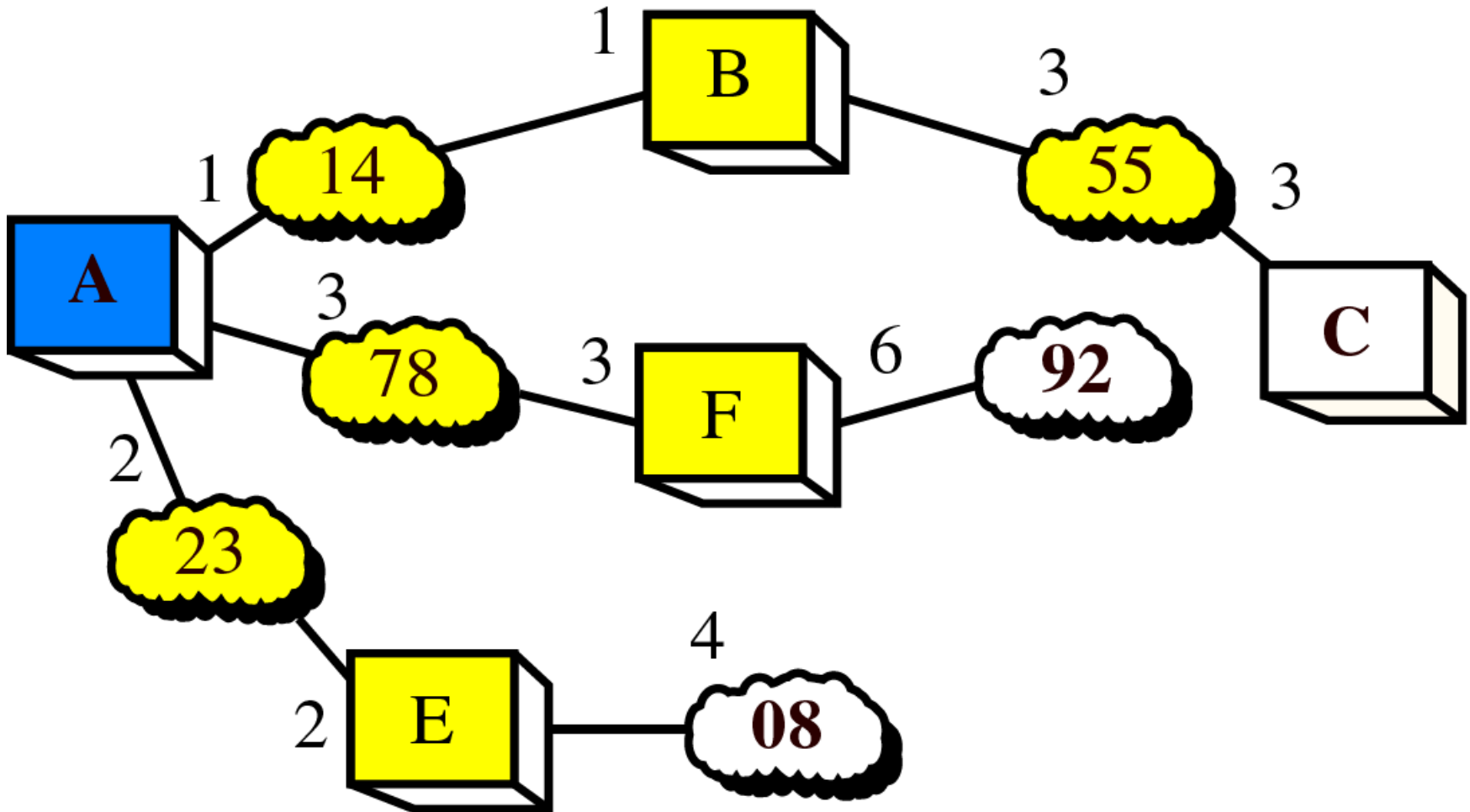
78 permanent E added

Shortest Path Calculation, Part VII



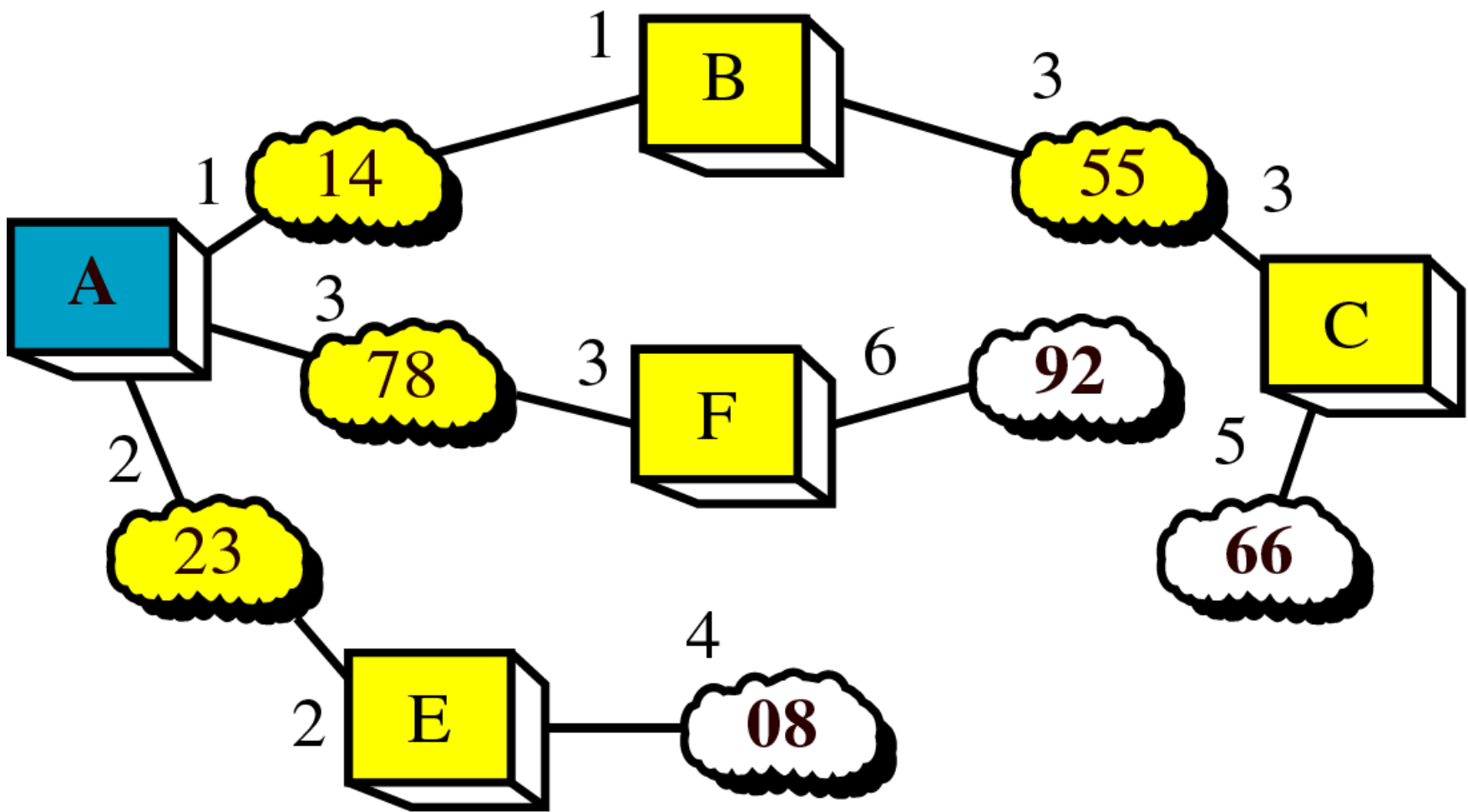
F permanent, 92 added

Shortest Path Calculation, Part VIII



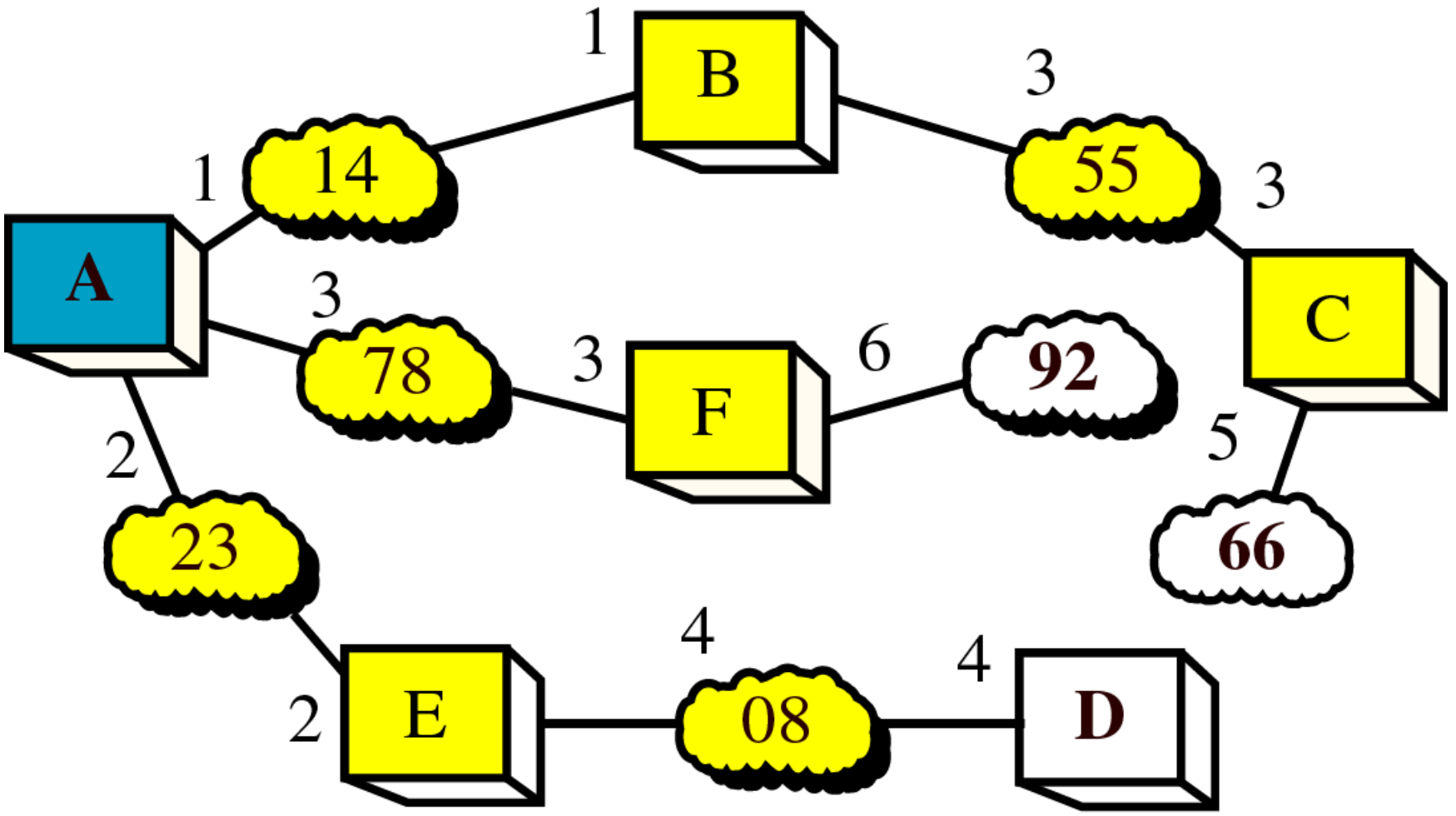
55 permanent, C added

Shortest Path Calculation, Part IX



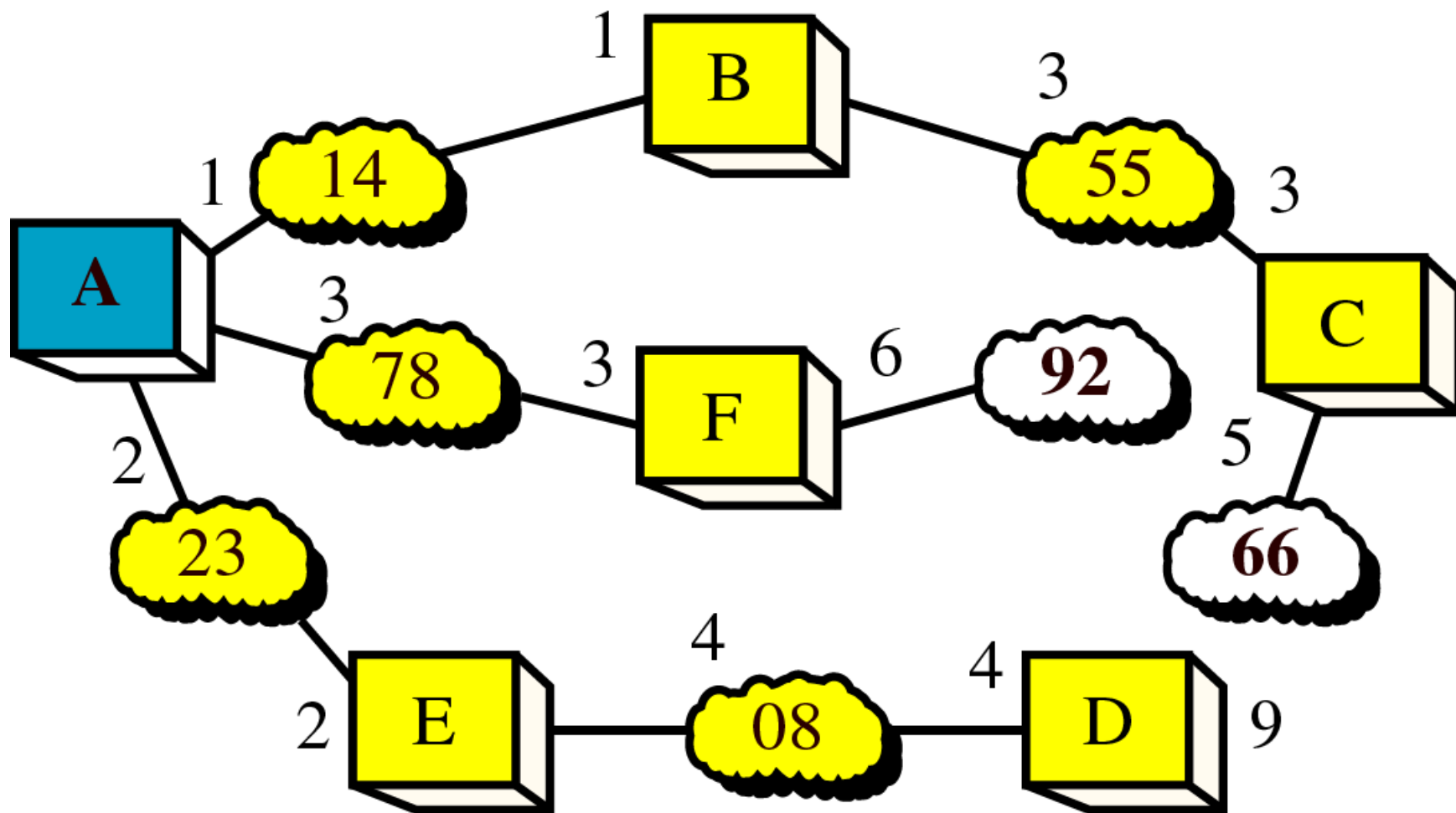
C permanent, 66 added

Shortest Path Calculation, Part X



08 permanent, D added

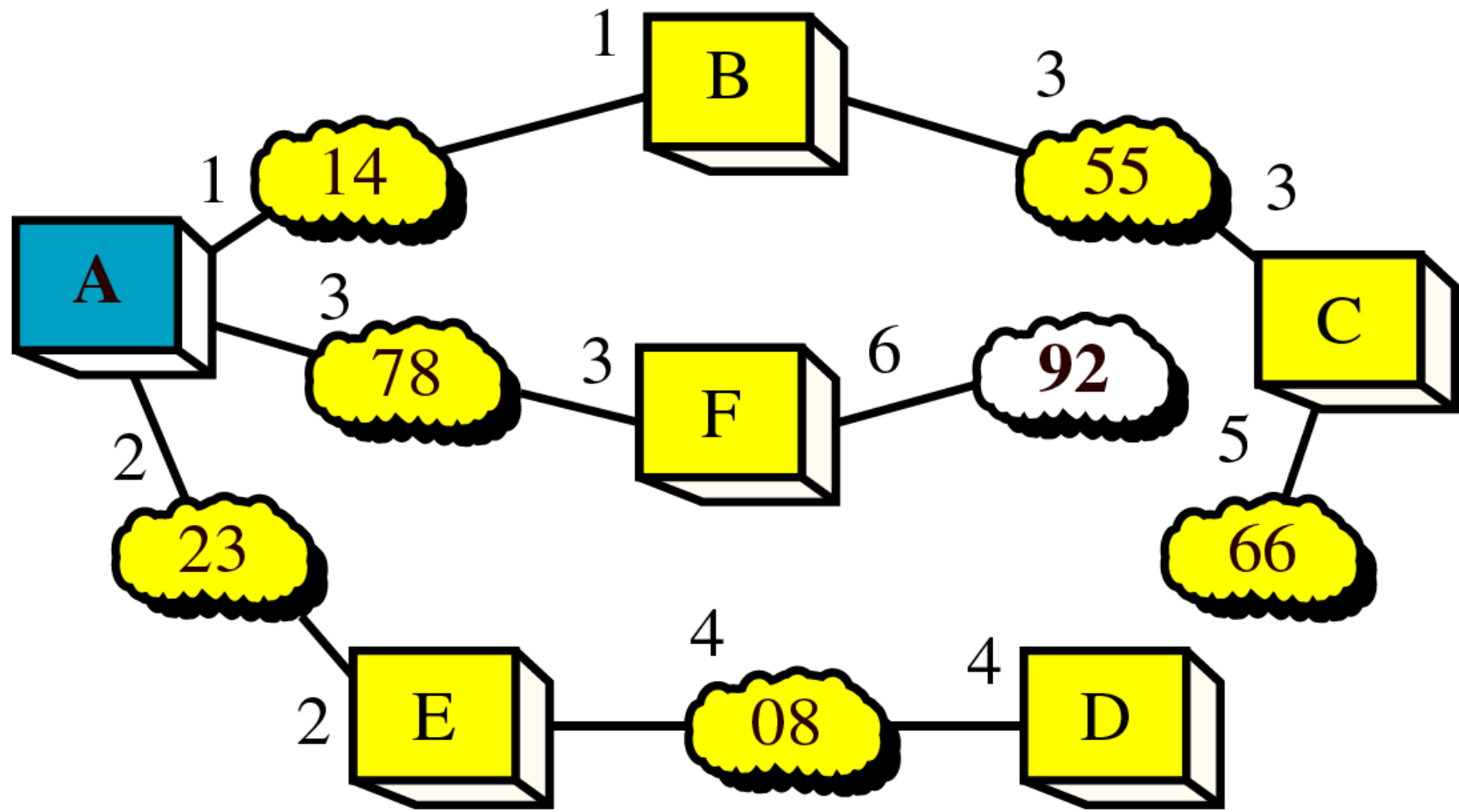
Shortest Path Calculation, Part XI



D permanent, 66 added.

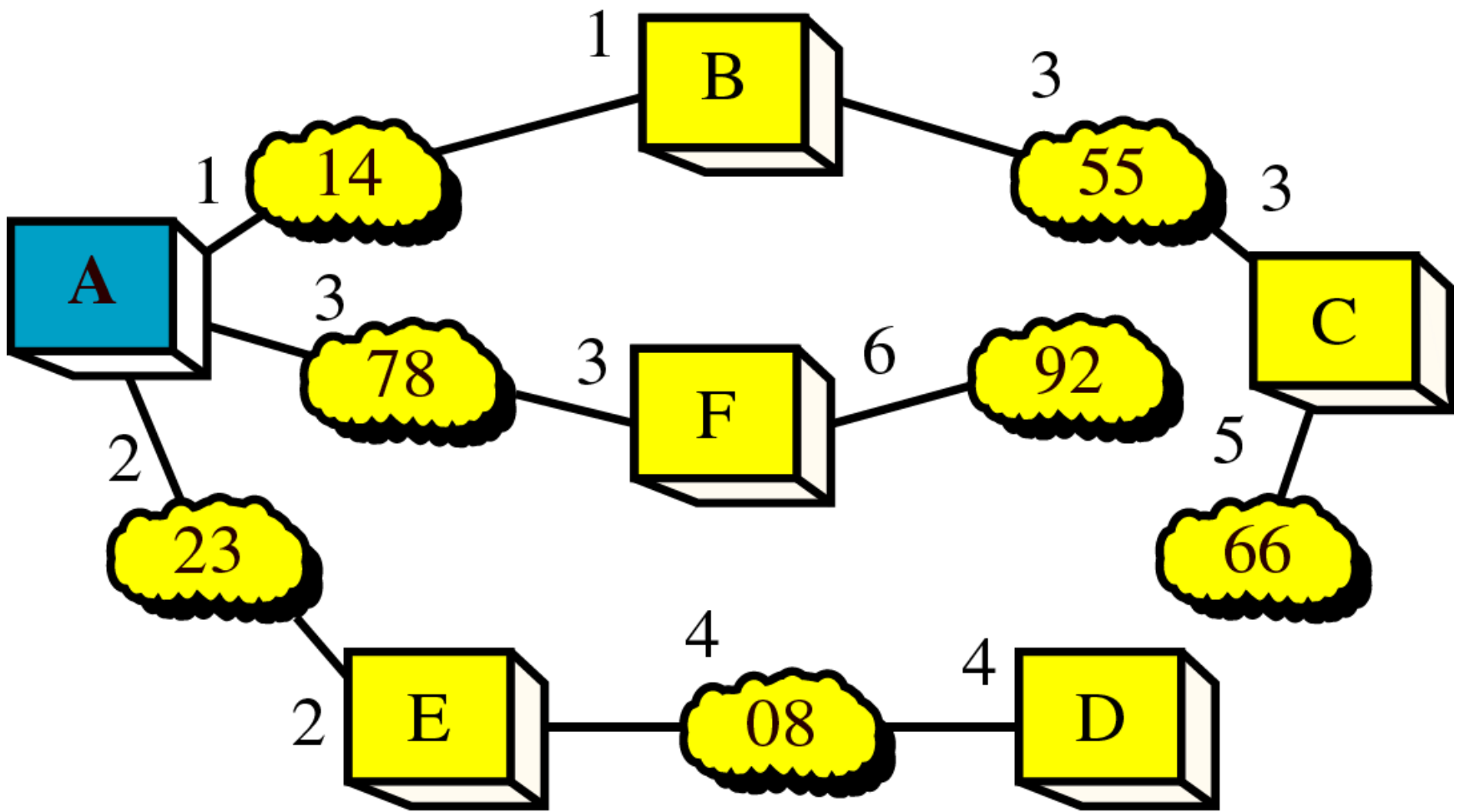
But $9 > 5$, so that link deleted

Shortest Path Calculation, Part XII



66 permanent

Shortest Path Calculation, Part XIII



Routing Table for Router A

Net	Cost	Next router
08	4	E
14	1	--
23	2	--
55	3	B
66	5	B
78	3	--
92	6	F